

# Geometry Processing

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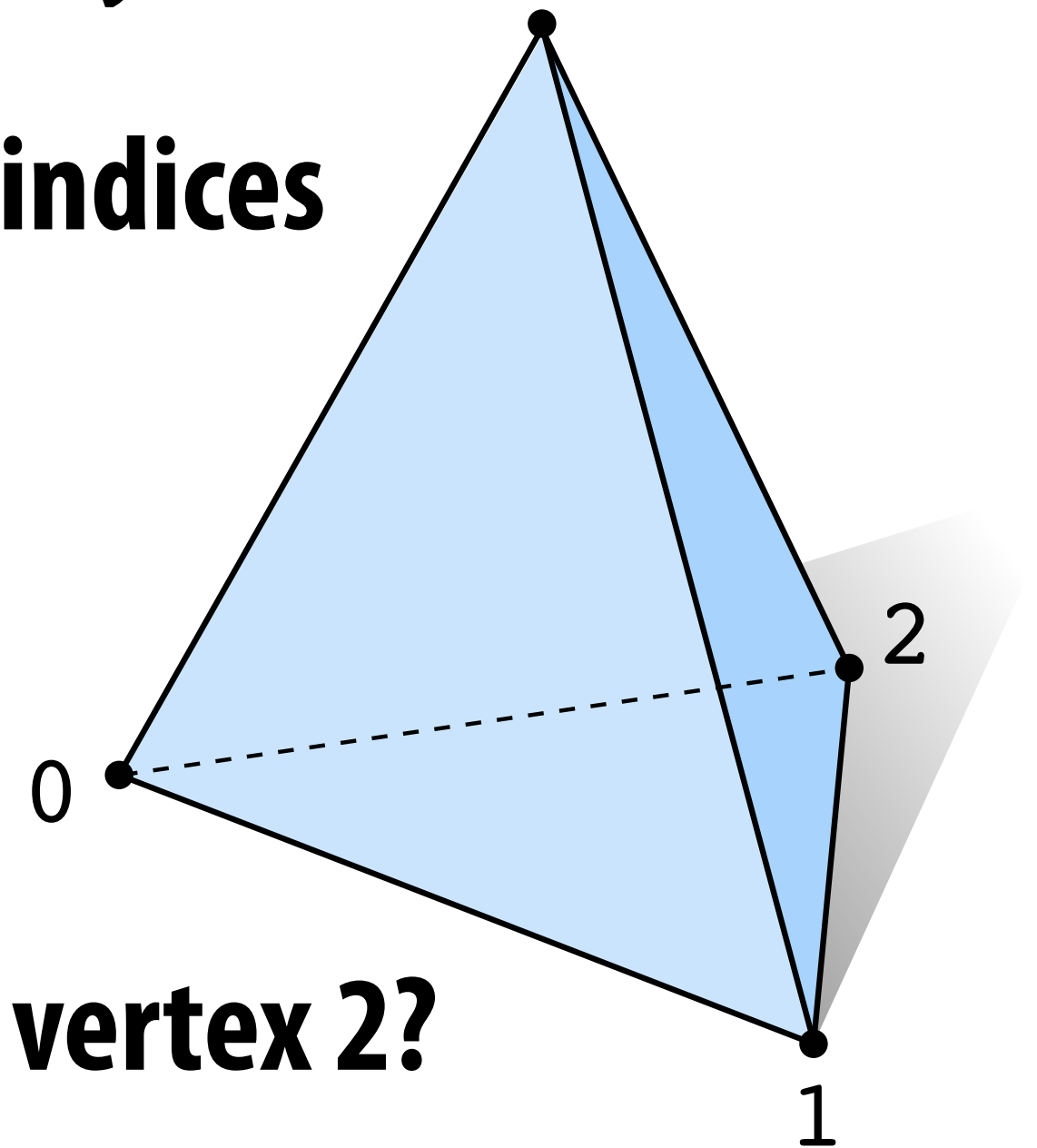
**Computer Graphics**  
**CMU 15-462/15-662**

# Last time — Adjacency List (Array-like)<sub>3</sub>

- Store triples of coordinates  $(x,y,z)$ , tuples of indices

- E.g., tetrahedron:

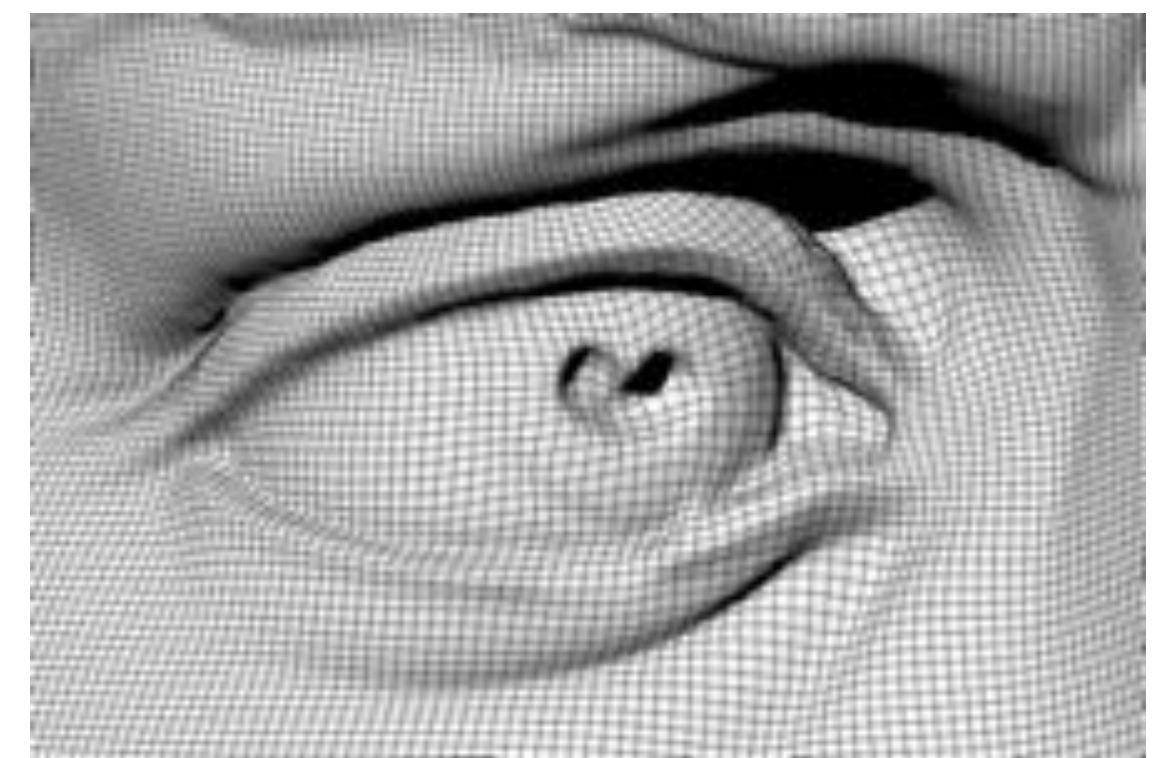
	VERTICES			POLYGONS		
	x	y	z	i	j	k
0:	-1	-1	-1	0	2	1
1:	1	-1	1	0	3	2
2:	1	1	-1	3	0	1
3:	-1	1	1	3	1	2



- Q: How do we find all the polygons touching vertex 2?

- Ok, now consider a more complicated mesh:

~1 billion polygons



**Very expensive to find the neighboring polygons! (What's the cost?)**

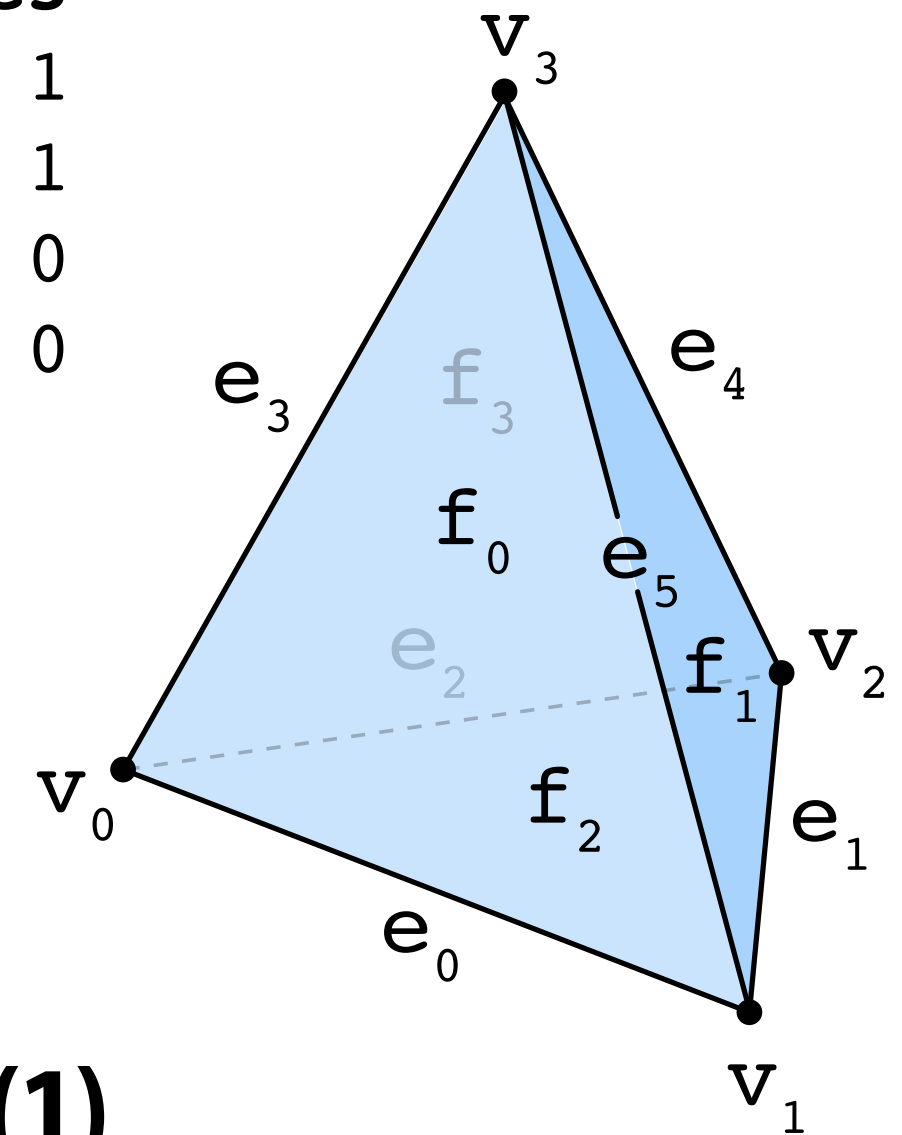
# Last time — Incidence Matrices

- If we want to know who our neighbors are, why not just store a list of neighbors?

- Can encode all neighbor information via incidence matrices

- E.g., tetrahedron:

	<u>VERTEX↔EDGE</u>				<u>EDGE↔FACE</u>						
	v0	v1	v2	v3	e0	e1	e2	e3	e4	e5	
e0	1	1	0	0	f0	1	0	0	1	0	1
e1	0	1	1	0	f1	0	1	0	0	1	1
e2	1	0	1	0	f2	1	1	1	0	0	0
e3	1	0	0	1	f3	0	0	1	1	1	0
e4	0	0	1	1							
e5	0	1	0	1							

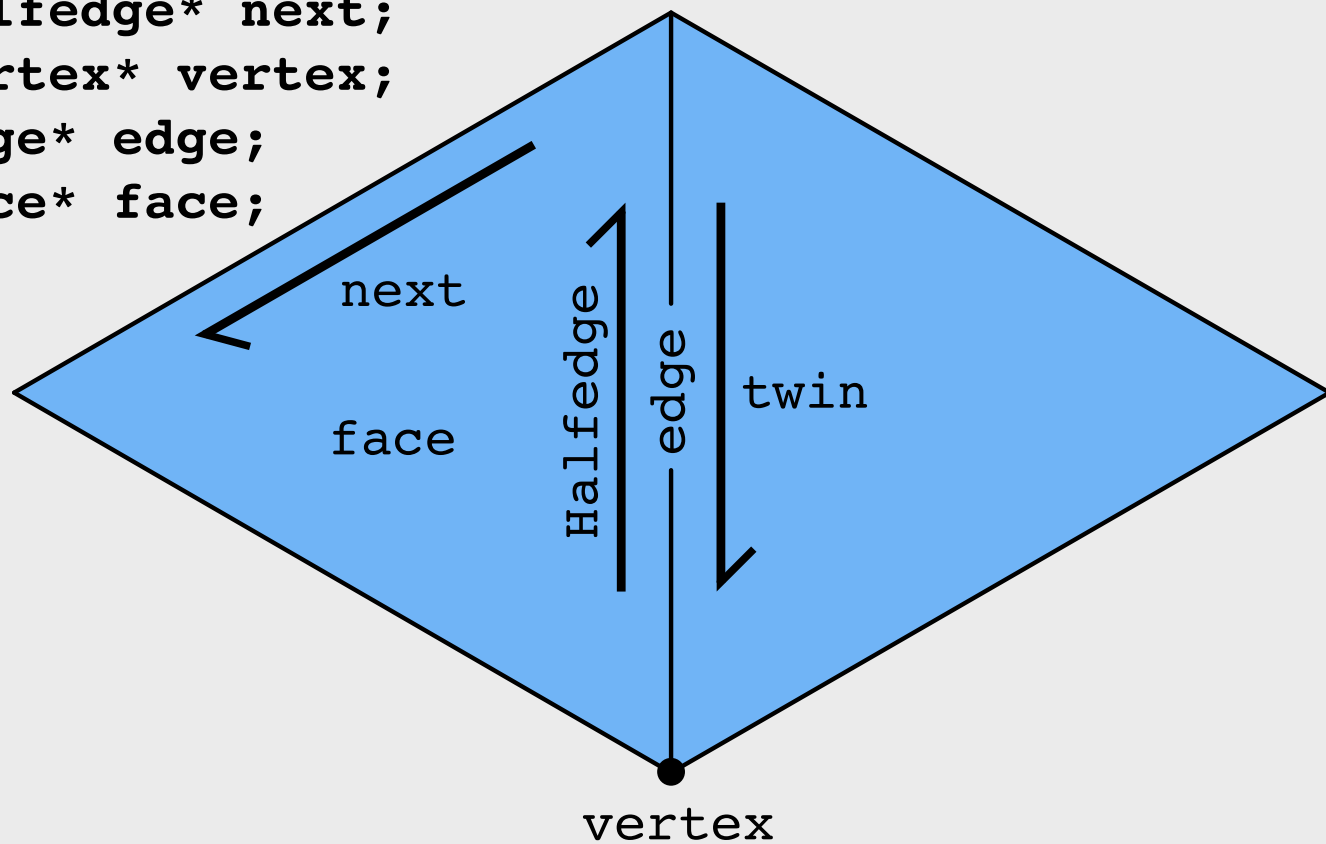


- 1 means “touches”; 0 means “does not touch”
- Instead of storing lots of 0’s, use sparse matrices
- Still large storage cost, but finding neighbors is now  $O(1)$
- Hard to change connectivity, since we used fixed indices
- Bonus feature: mesh does not have to be manifold

# Last time — Halfedge Data Structure

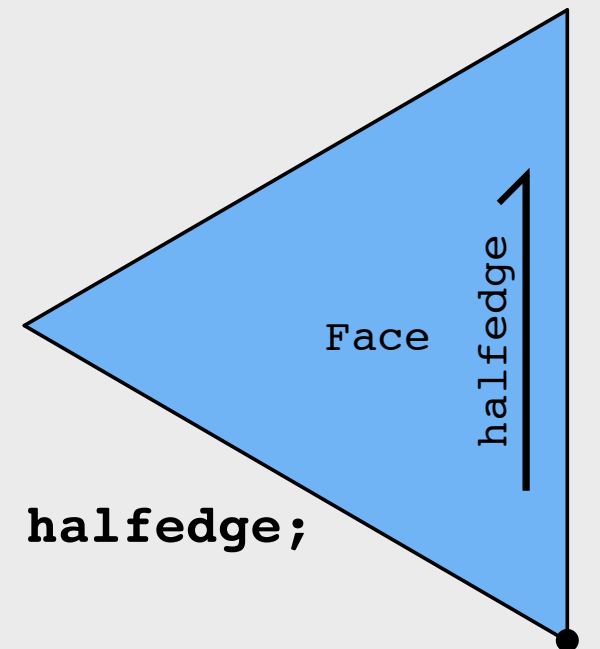
- Store some information about neighbors
- Don't need an exhaustive list; just a few key pointers
- Key idea: two halfedges act as “glue” between mesh elements:

```
struct Halfedge
{
  Halfedge* twin;
  Halfedge* next;
  Vertex* vertex;
  Edge* edge;
  Face* face;
};
```

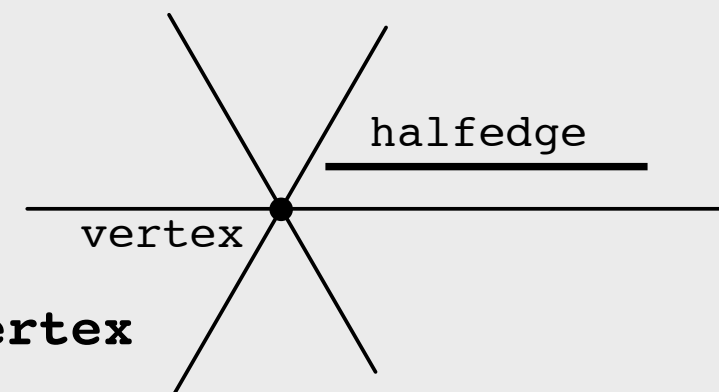


```
struct Edge
{
  Halfedge* halfedge;
};
```

```
struct Face
{
  Halfedge* halfedge;
};
```



```
struct Vertex
{
  Halfedge* halfedge;
};
```



- Each vertex, edge face points to just one of its halfedges.

# Comparison of Polygon Mesh Data Structures

	Adjacency List	Incidence Matrices	Halfedge Mesh
constant-time neighborhood access?	NO	YES	YES
easy to add/remove mesh elements?	NO	NO	YES
nonmanifold geometry?	YES	YES	NO

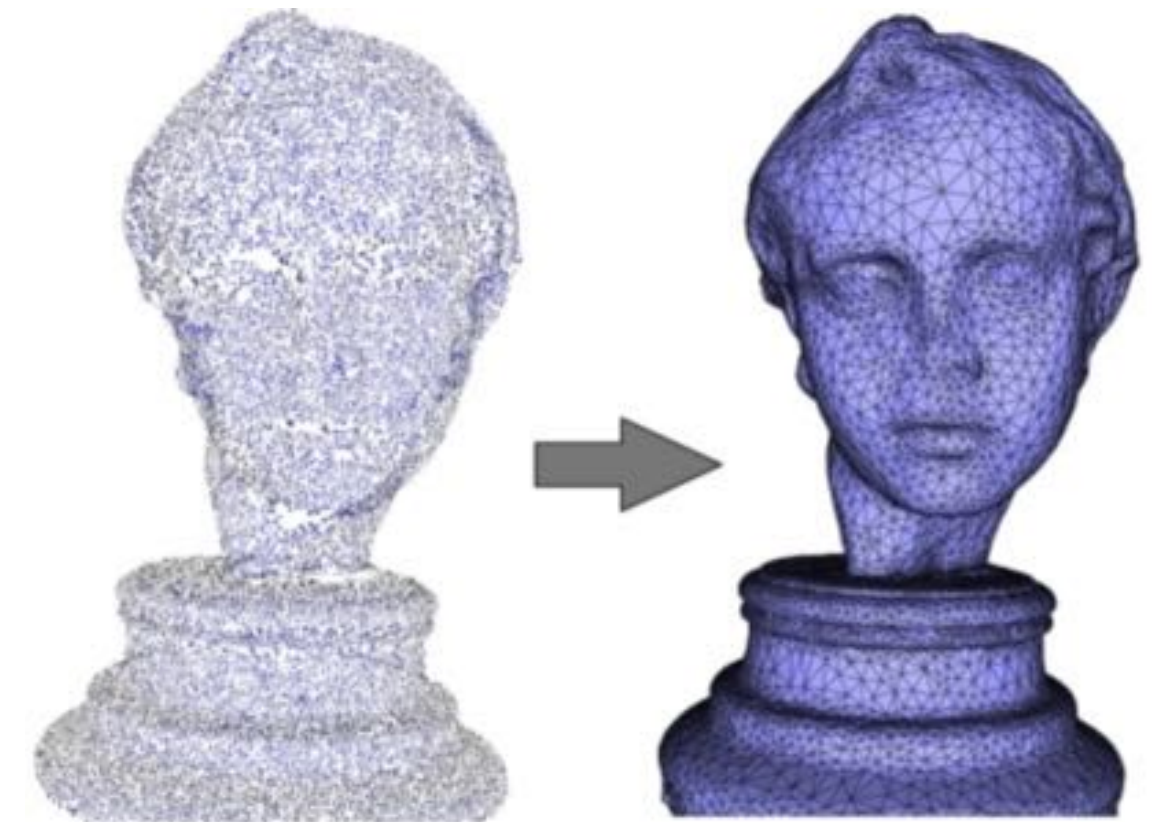
**Conclusion: pick the right data structure for the job!**

**Ok, but what can we actually do with our  
fancy new data structures?**



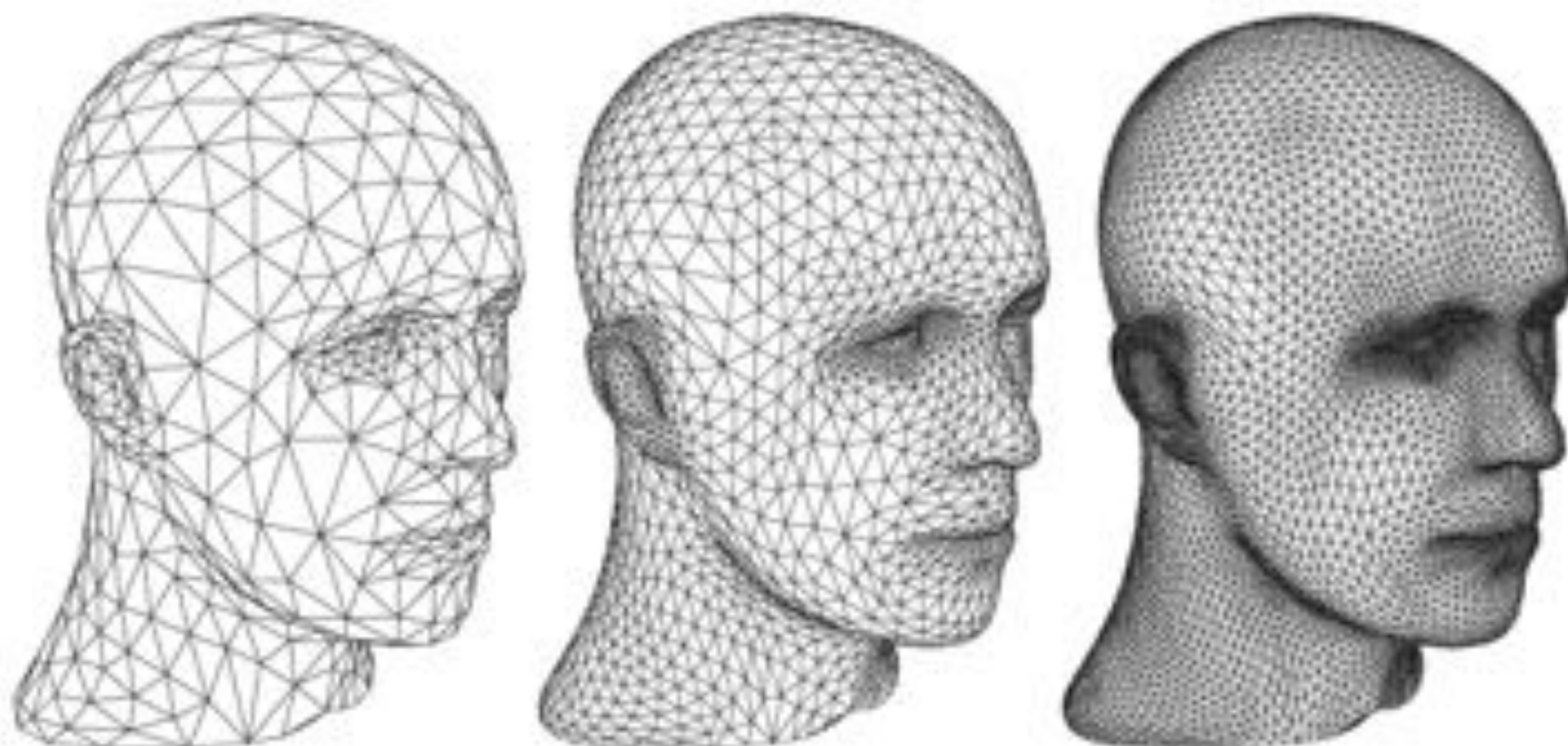
# Geometry Processing: Reconstruction

- **Given samples of geometry, reconstruct surface**
- **What are “samples”? Many possibilities:**
  - **points, points & normals, ...**
  - **image pairs / sets (multi-view stereo)**
  - **line density integrals (MRI/CT scans)**
- **How do you get a surface? Many techniques:**
  - **silhouette-based (visual hull)**
  - **Voronoi-based (e.g., power crust)**
  - **PDE-based (e.g., Poisson reconstruction)**
  - **Radon transform / isosurfacing (marching cubes)**



# Geometry Processing: Upsampling

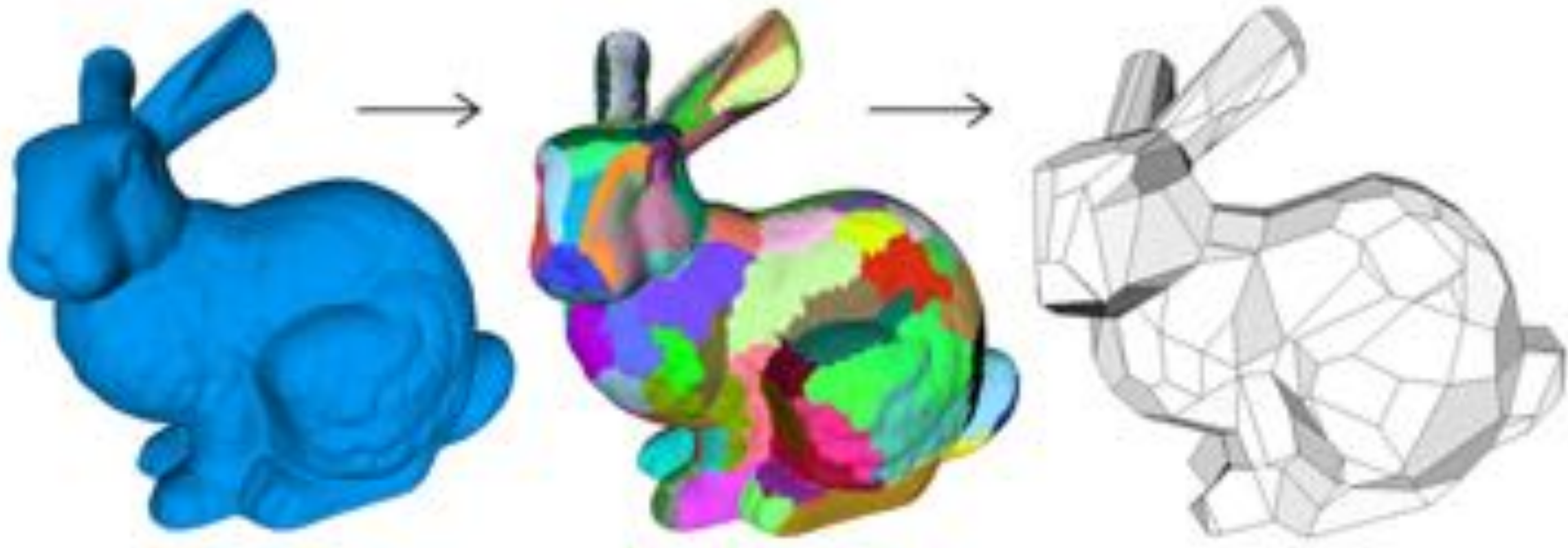
- Increase resolution via interpolation
- Images: e.g., bilinear, bicubic interpolation
- Polygon meshes:
  - subdivision
  - bilateral upsampling
  - ...





# Geometry Processing: Downsampling

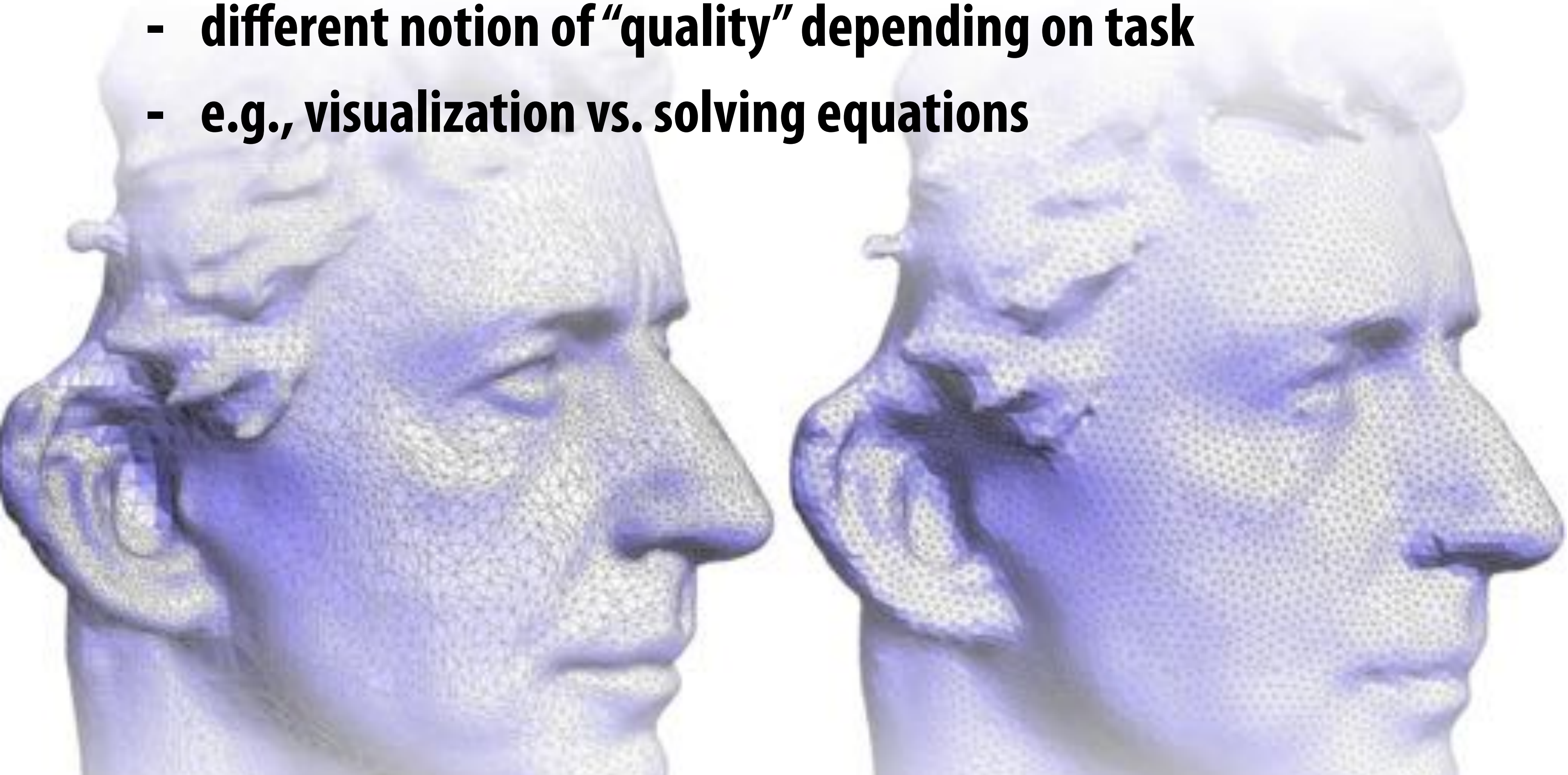
- **Decrease resolution; try to preserve shape/appearance**
- **Images: nearest-neighbor, bilinear, bicubic interpolation**
- **Point clouds: subsampling (just take fewer points!)**
- **Polygon meshes:**
  - **iterative decimation, variational shape approximation, ...**





# Geometry Processing: Resampling

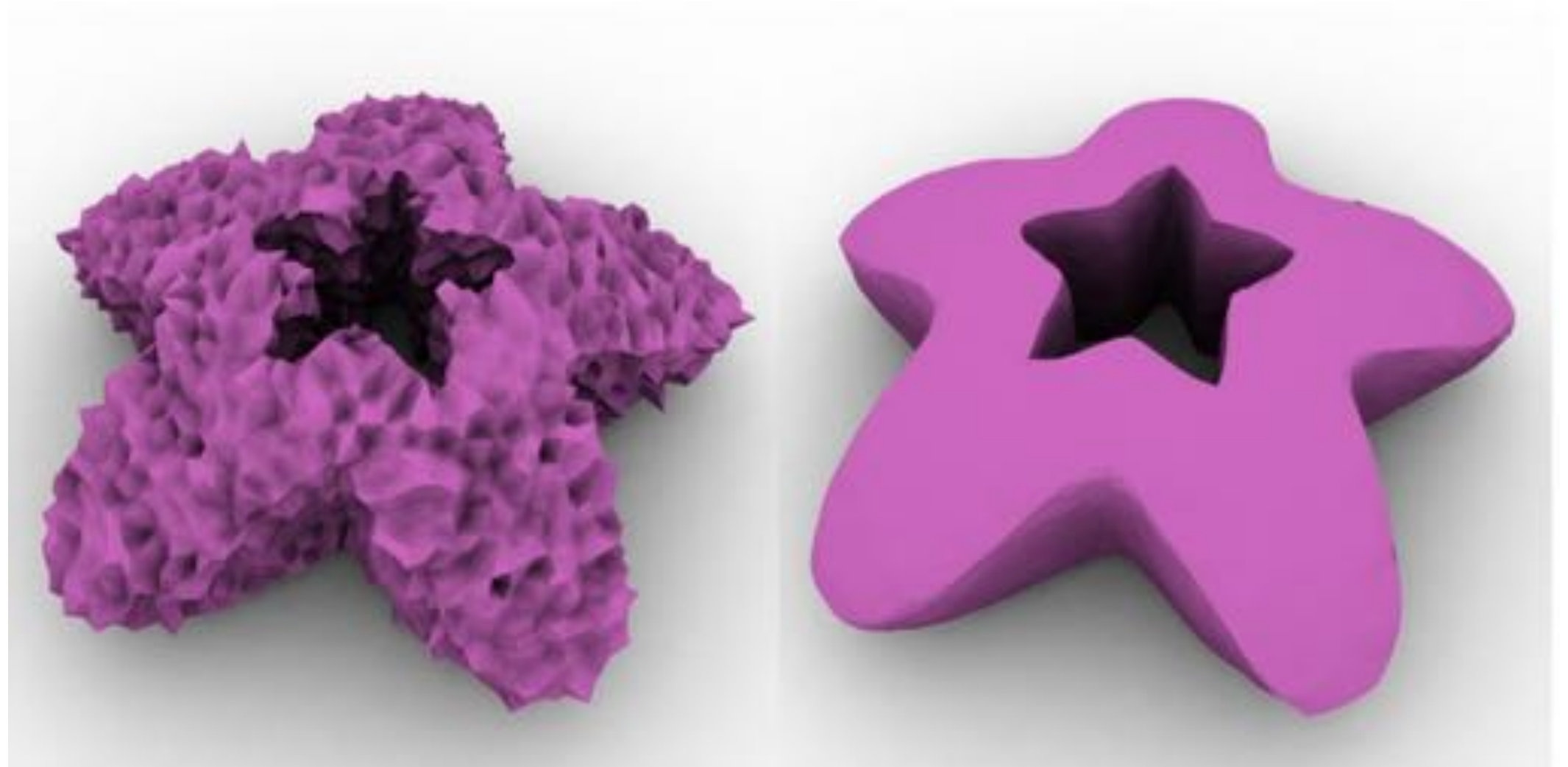
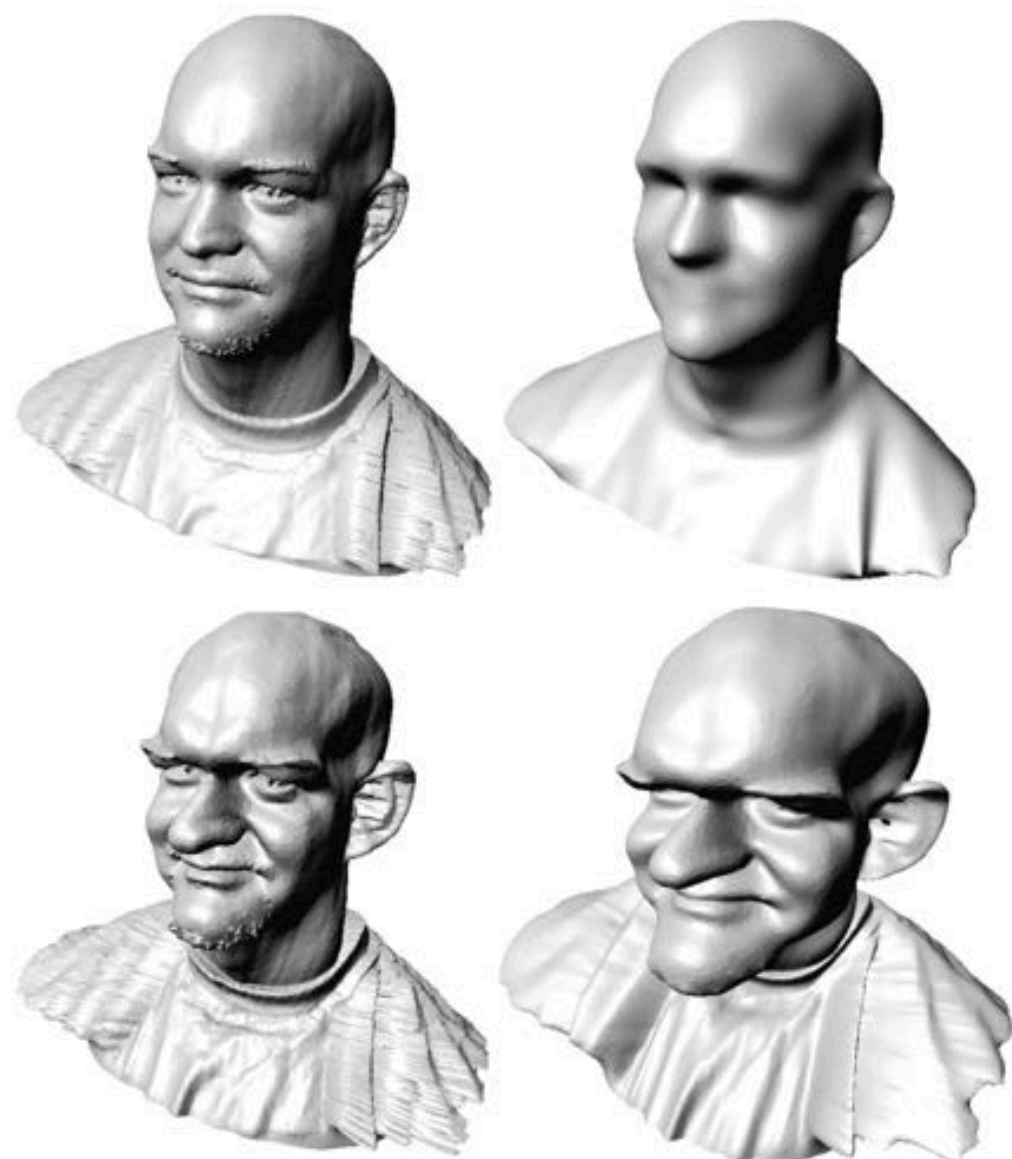
- **Modify sample distribution to improve quality**
- **Images: not an issue! (Pixels always stored on a regular grid)**
- **Meshes: shape of polygons is extremely important!**
  - **different notion of “quality” depending on task**
  - **e.g., visualization vs. solving equations**





# Geometry Processing: Filtering

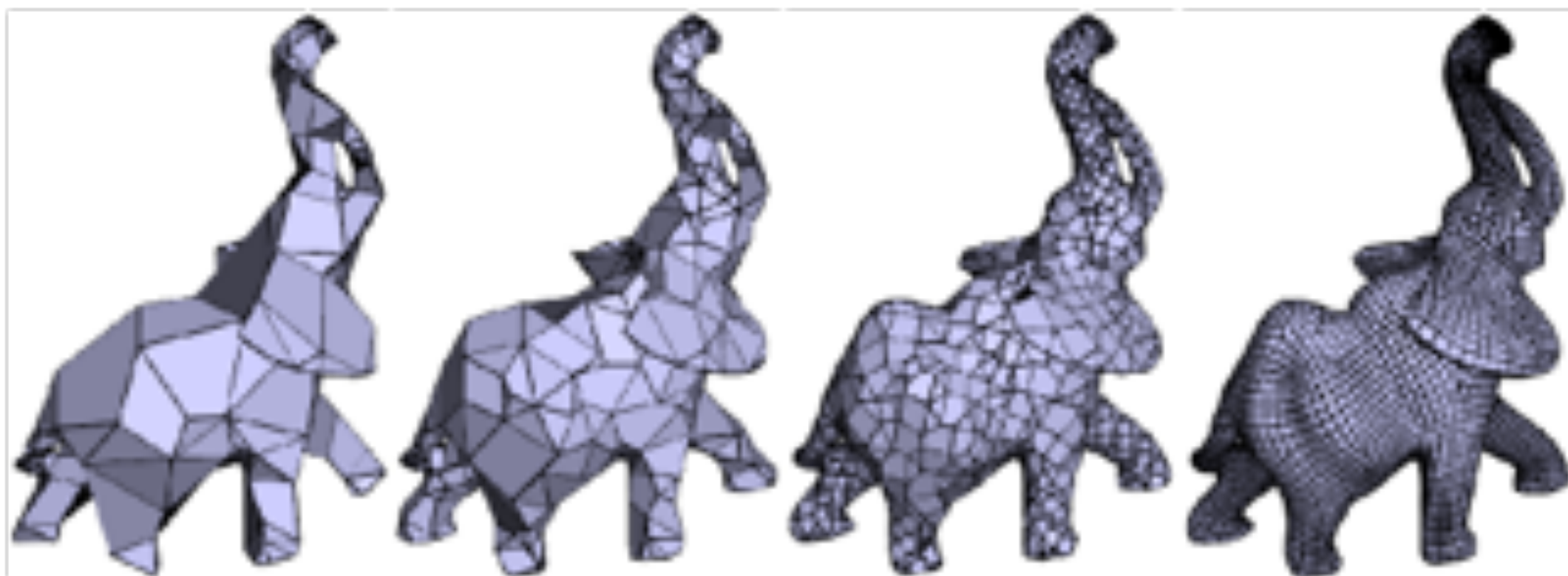
- Remove noise, or emphasize important features (e.g., edges)
- Images: blurring, bilateral filter, edge detection, ...
- Polygon meshes:
  - curvature flow
  - bilateral filter
  - spectral filter





# Geometry Processing: Compression

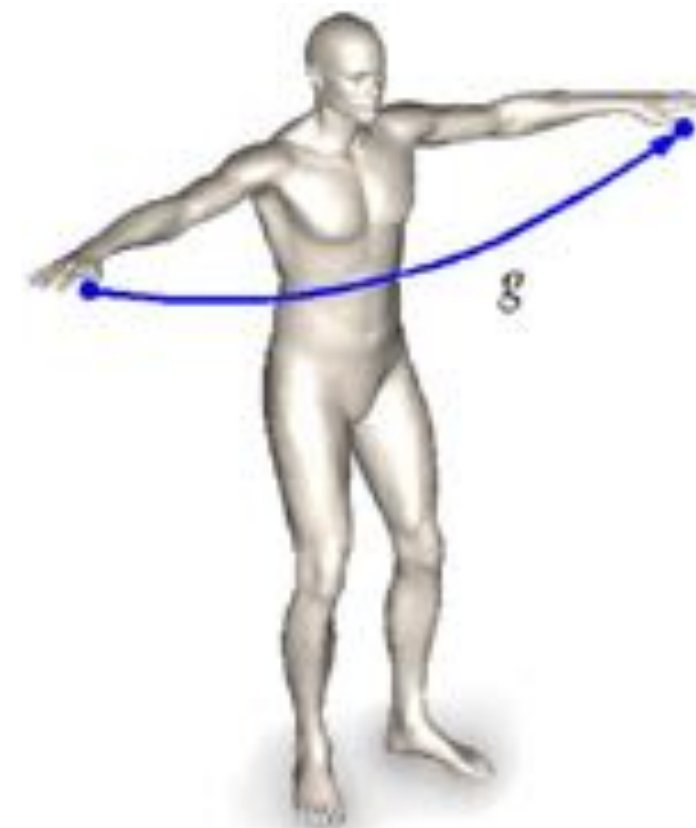
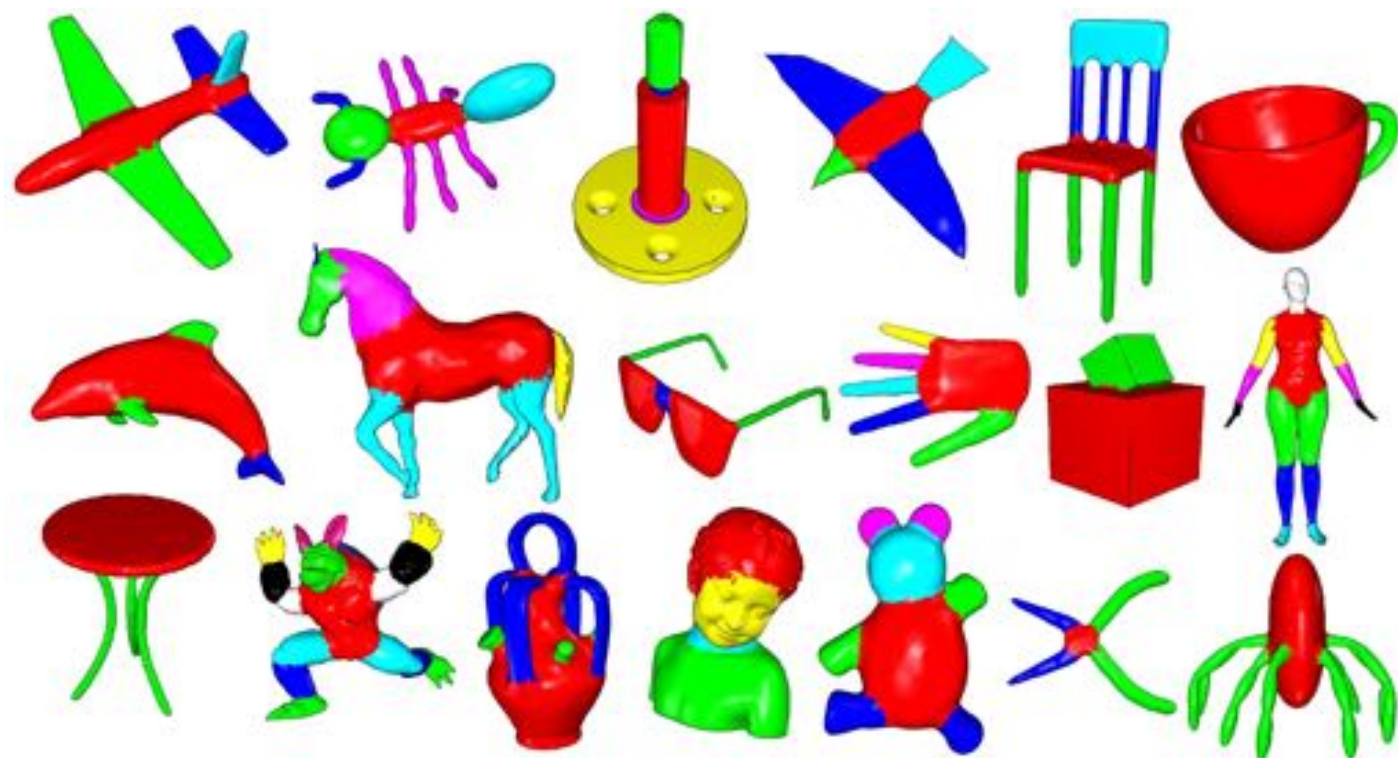
- Reduce storage size by eliminating redundant data/  
approximating unimportant data
- Images:
  - run-length, Huffman coding - lossless
  - cosine/wavelet (JPEG/MPEG) - lossy
- Polygon meshes:
  - compress geometry and connectivity
  - many techniques (lossy & lossless)





# Geometry Processing: Shape Analysis

- Identify/understand important semantic features
- Images: computer vision, segmentation, face detection, ...
- Polygon meshes:
  - segmentation, correspondence, symmetry detection, ...



Extrinsic symmetry



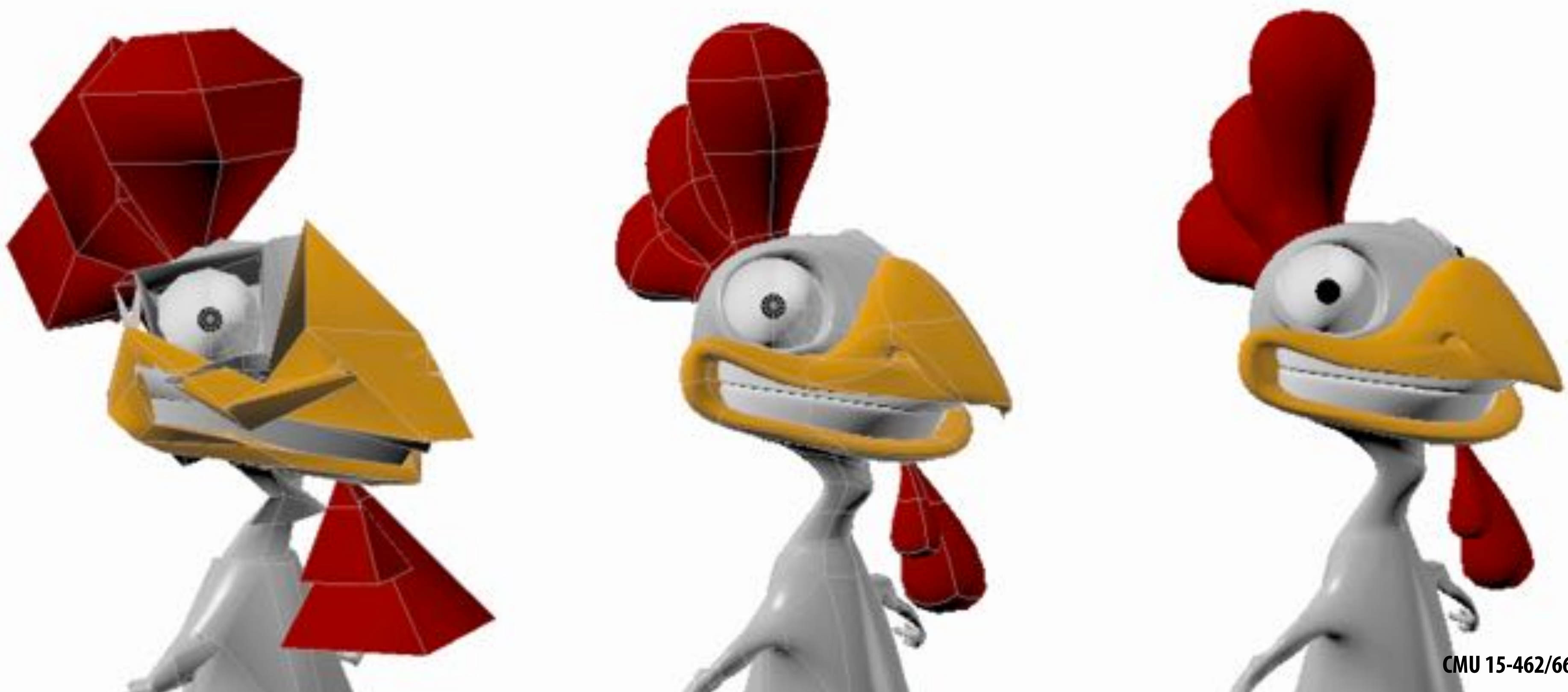
Intrinsic symmetry





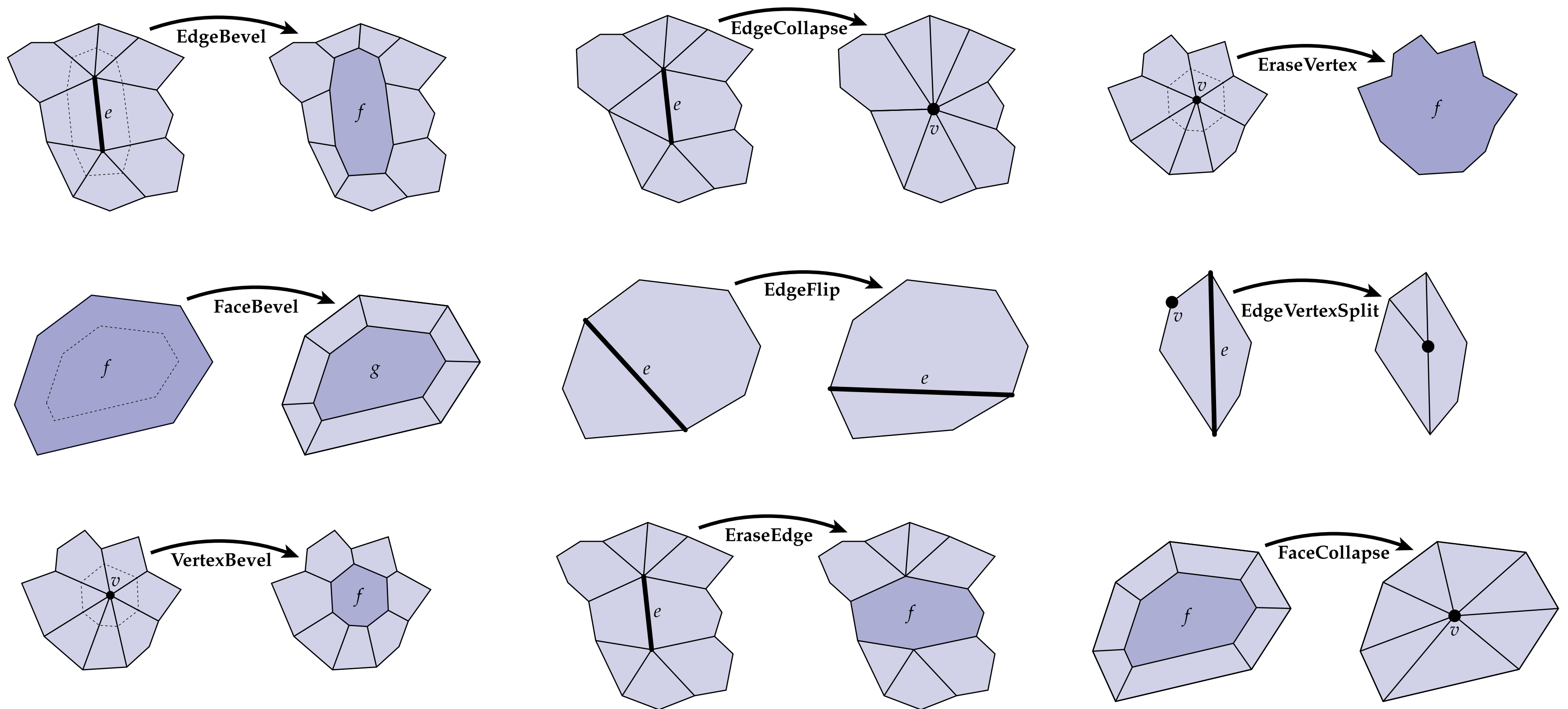
# Subdivision Modeling

- **Common modeling paradigm in modern 3D tools:**
  - **Coarse “control cage”**
  - **Perform local operations to control/edit shape**
  - **Global subdivision process determines final surface**



# Subdivision Modeling—Local Operations

- For general polygon meshes, we can dream up lots of local mesh operations that might be useful for modeling:



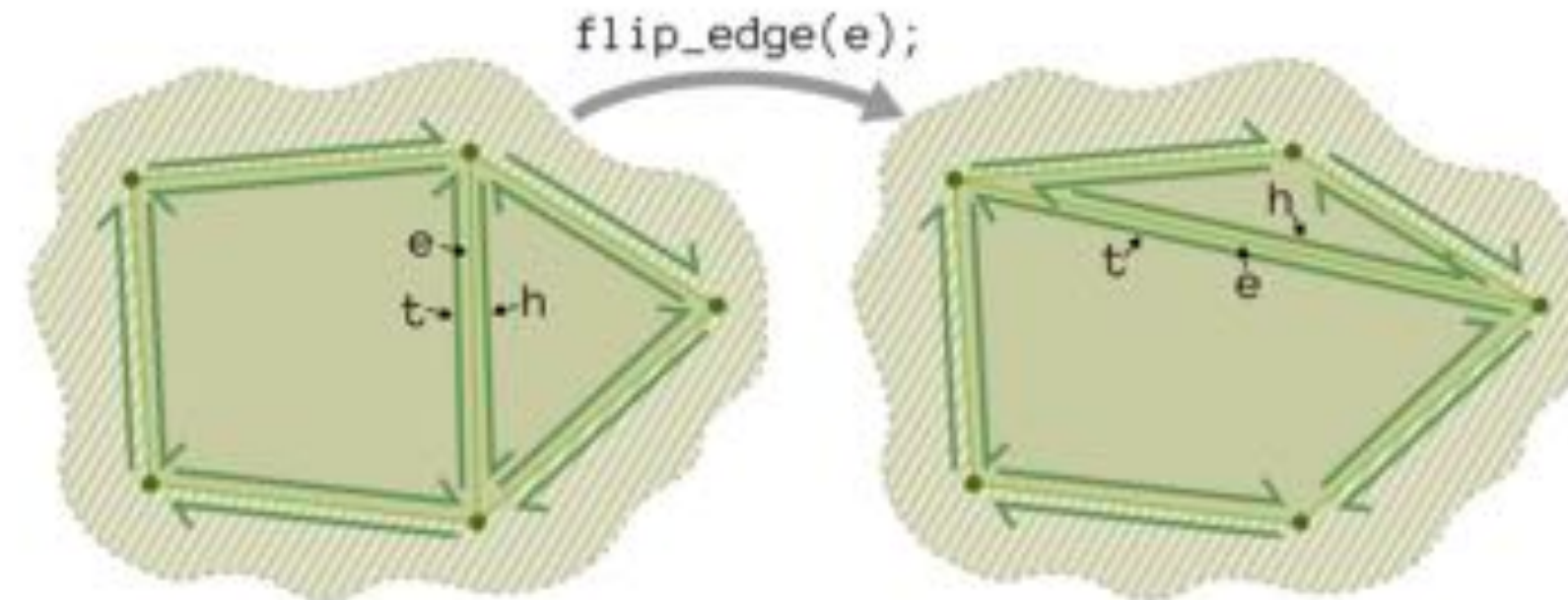
...and many, many more!



# A2.0 Diagrams

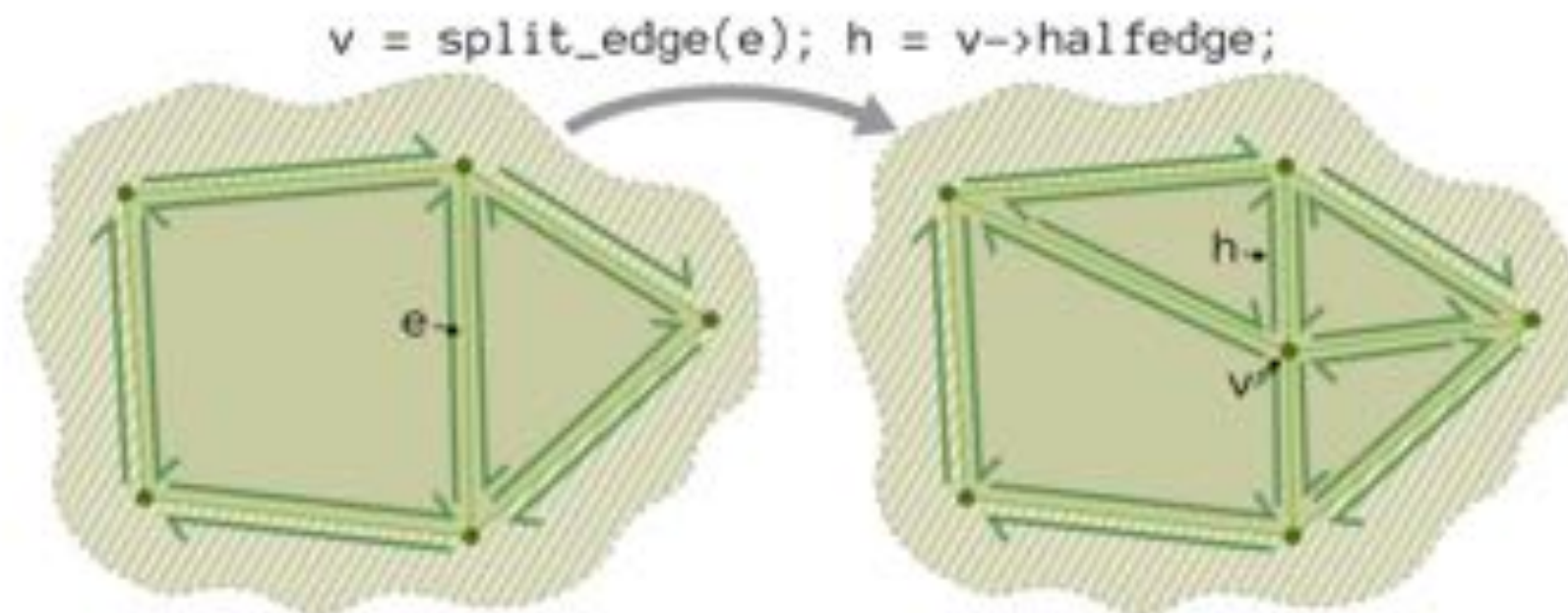
## A2L1 Flip Edge

Given a non-boundary edge, rotate it to the vertices one step counterclockwise from its current endpoints.

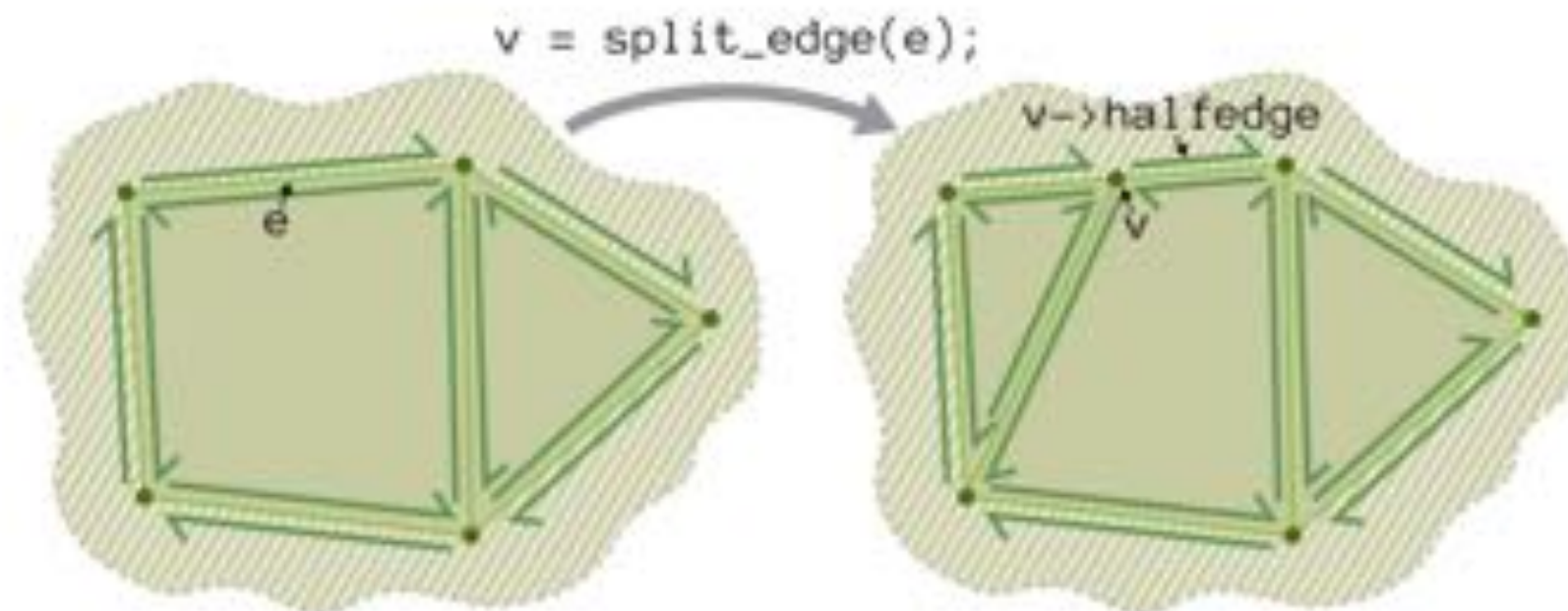


## A2L2 Split Edge

Given an edge, add a vertex in the middle and an edge from the newly-added vertex to the next-in-ccw-order-from-the-edge vertex in the face.



Split edge does not subdivide boundary faces.

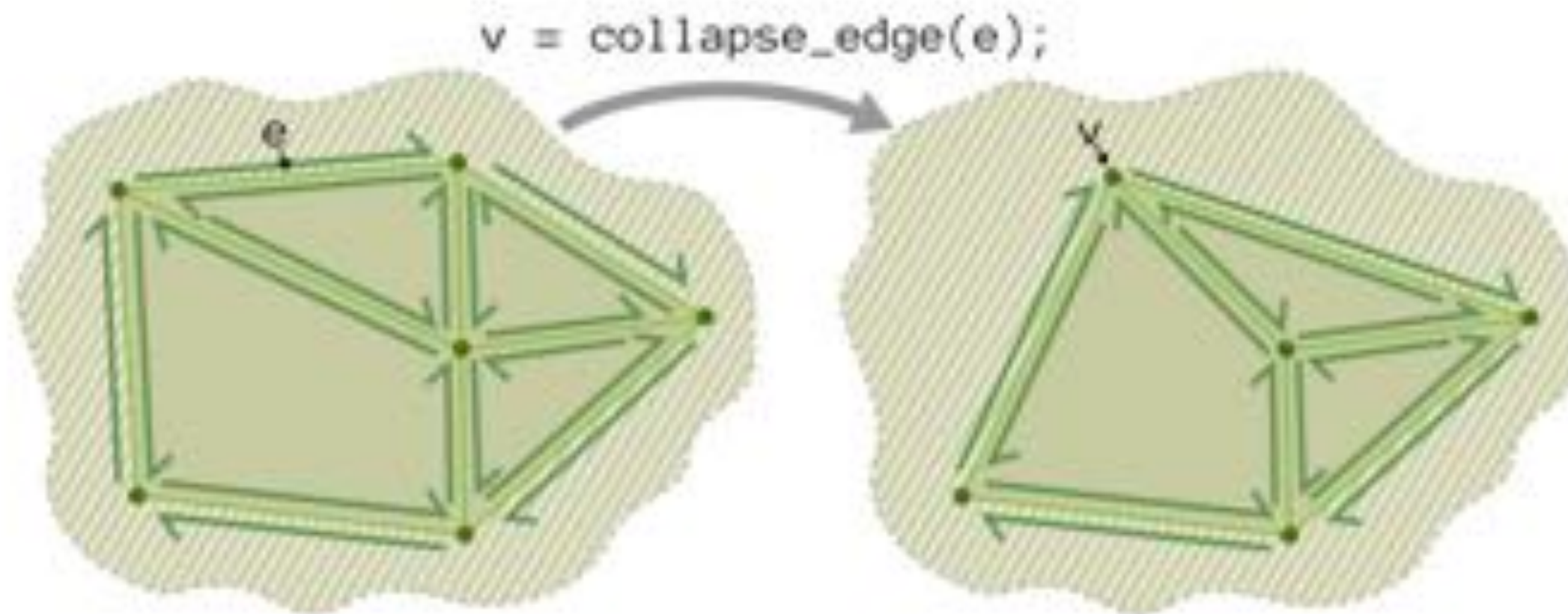




# A2.0 Diagrams

## A2L3 Collapse Edge

Merge the two endpoints of an edge at the midpoint of the edge. This will combine all the edges to both endpoints into a new vertex (be careful to not have two edges overlapping).



If collapsing the edge would result in an invalid mesh, you should instead do nothing and return `std::nullopt`. Otherwise returns the newly collapsed vertex.

Examples of edge cases that we will be testing would be



(figure 8, and any generalizations of this)





# What is a valid mesh?

A `Halfedge_Mesh` structure is a sea of pointers. It is surprisingly easy to tangle these pointers into a mess that fails to represent a manifold mesh (or, really, a consistent mesh at all). In order to provide guidance on what constitutes a reasonable setting for a mesh, we define a valid mesh as one for which the following properties hold:

1. The mesh is self-contained. All pointers are to elements in the vertices, edges, faces, or halfedges lists.
2. Edges correspond to twinned halfedges. That is, for every edge `e`, `e->halfedge(->twin)^n` is a cycle of exactly two halfedges, and these are exactly the halfedges with `h->edge` equal to `e`.
3. Faces correspond to cycles of halfedges. That is, for every face `f`, `f->halfedge(->next)^n` is a cycle of at least three halfedges, and these are exactly the halfedges with `h->face` equal to `f`.
4. Vertices correspond to stars of halfedges. That is, for every vertex `v`, `v->halfedge(->twin->next)^n` is a cycle of at least two halfedges, and these are exactly the halfedges with `h->vertex` equal to `v`.
5. Vertices are not orphaned, nor is the surface adjacent to them hourglass-shaped. That is, vertices are adjacent to at least one non-boundary face and at most one boundary face.
6. Edges not orphaned. That is, edges are adjacent to at least one non-boundary face.
7. Faces are simple. That is, faces touch each vertex and edge at most once.

These properties are checked by the `Halfedge_Mesh::validate()` function, which returns `nullopt` if the mesh is valid or an element reference and an explanatory message if something is wrong with the mesh.

# Mesh Processing:

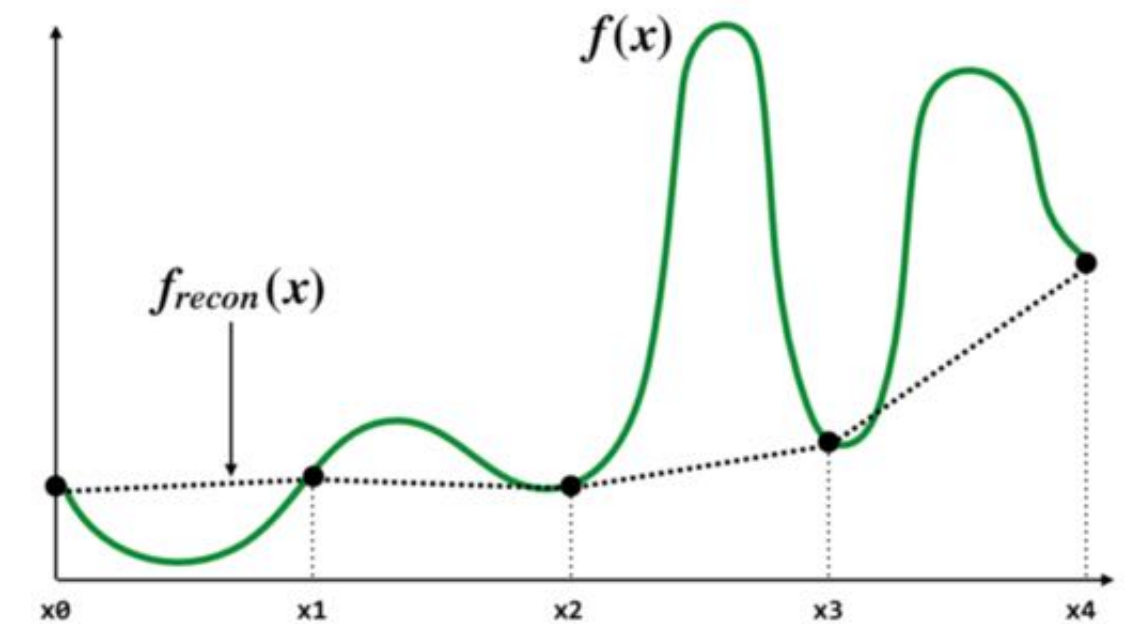
“If the surface resulting from an operation can be represented by *some* valid mesh, then run the operation and produce a valid mesh representing the result.”

**A few things to consider —**



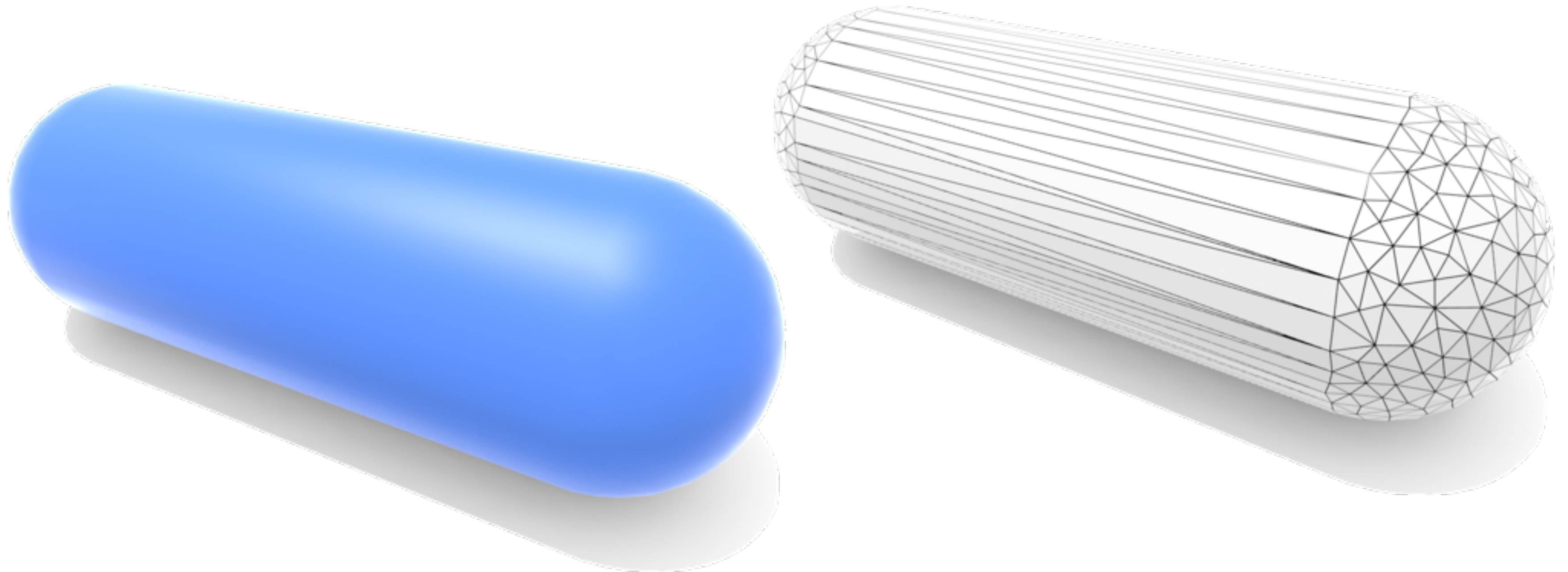
# Remeshing as resampling

- Remember our discussion of aliasing
- Bad sampling makes signal appear different than it really is
- E.g., undersampled curve looks flat
- Geometry is no different!
  - undersampling destroys features
  - oversampling bad for performance



# What makes a “good” mesh?

- **One idea: good approximation of original shape!**
- **Keep only elements that contribute information about shape**
- **Add additional information where, e.g., curvature is large**

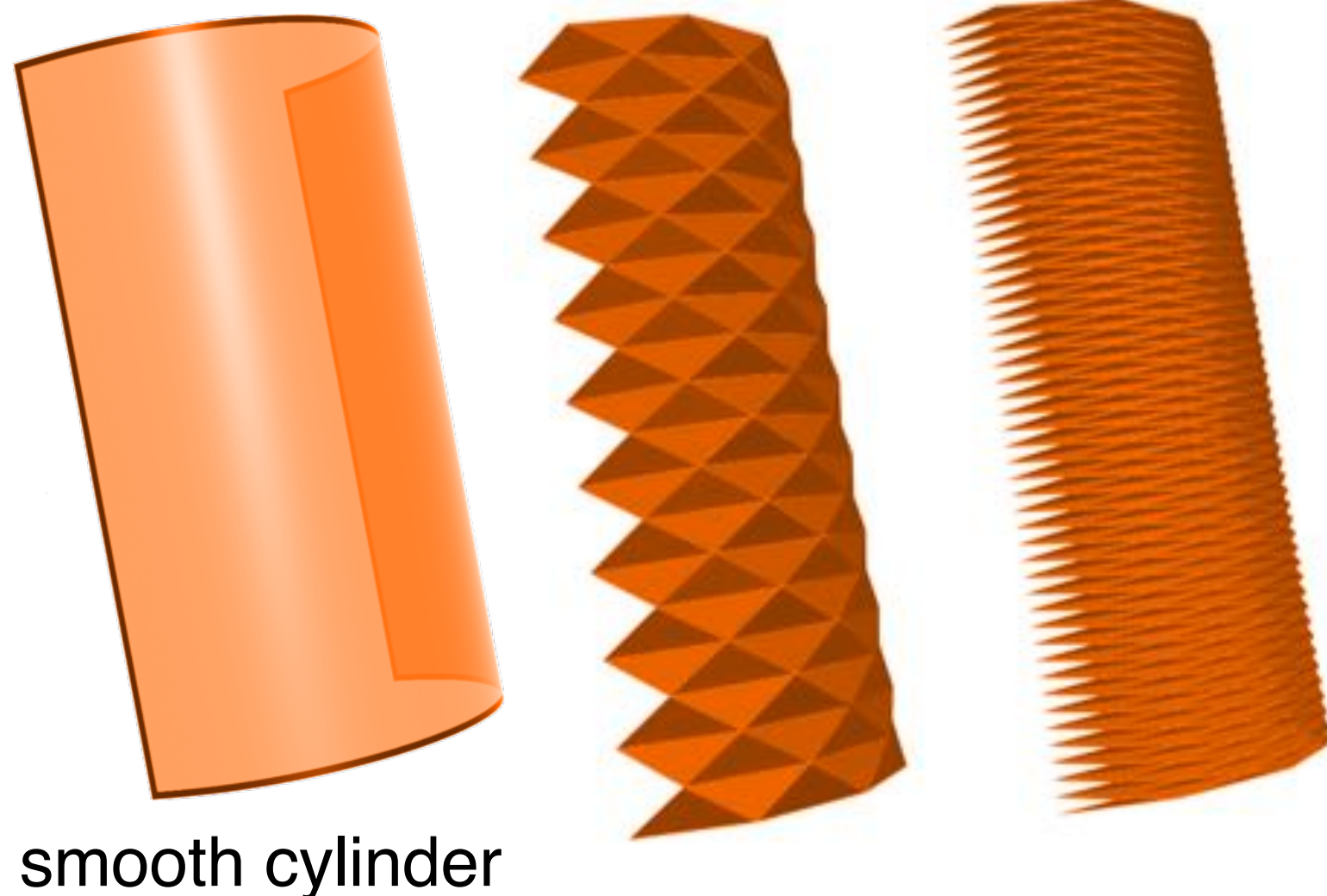




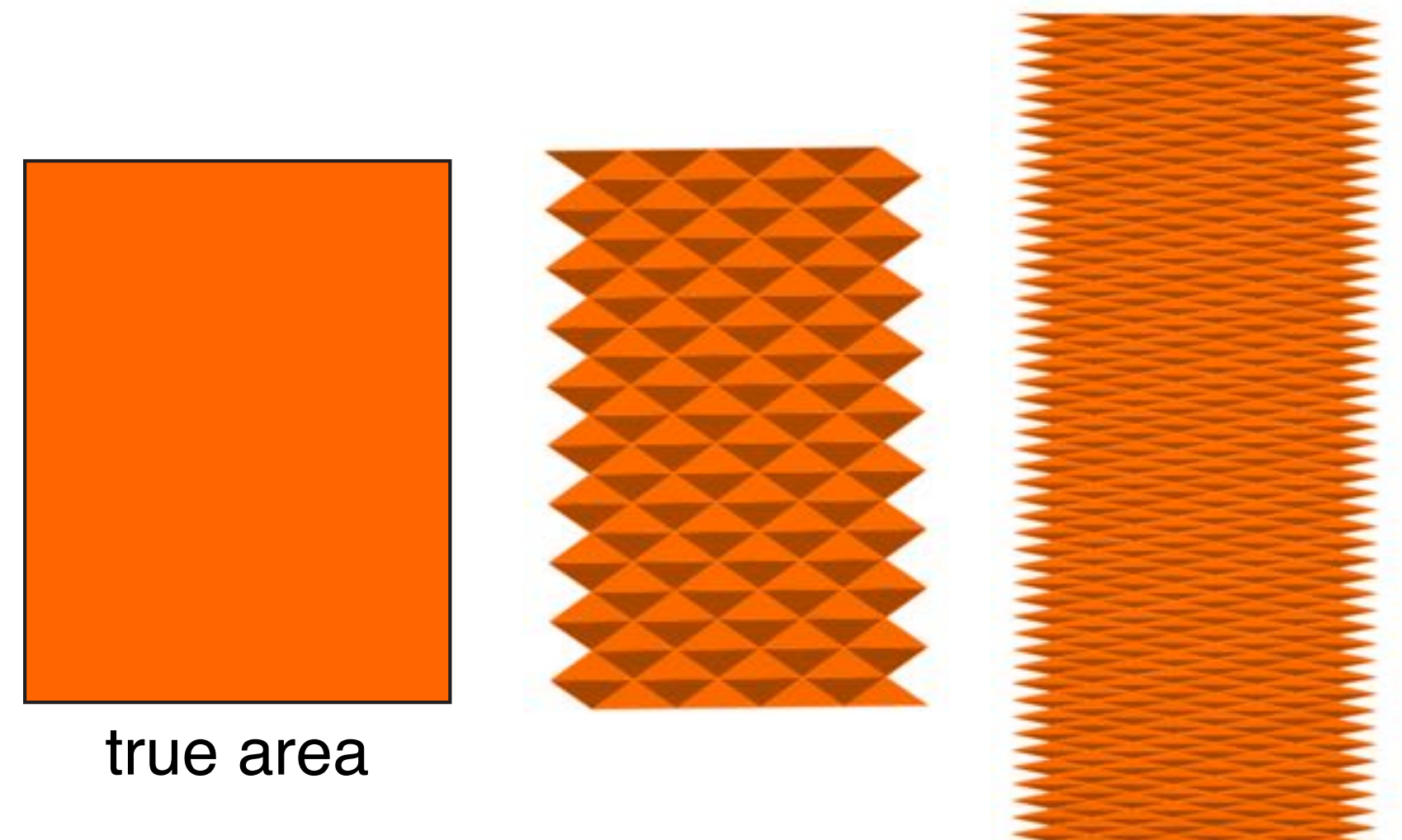
# Approximation of position is not enough!

- Just because the vertices of a mesh are close to the surface it approximates does not mean it's a good approximation!
- Can still have wrong appearance, wrong area, wrong...
- Need to consider other factors\*, e.g., close approximation of surface normals

vertices exactly on smooth cylinder



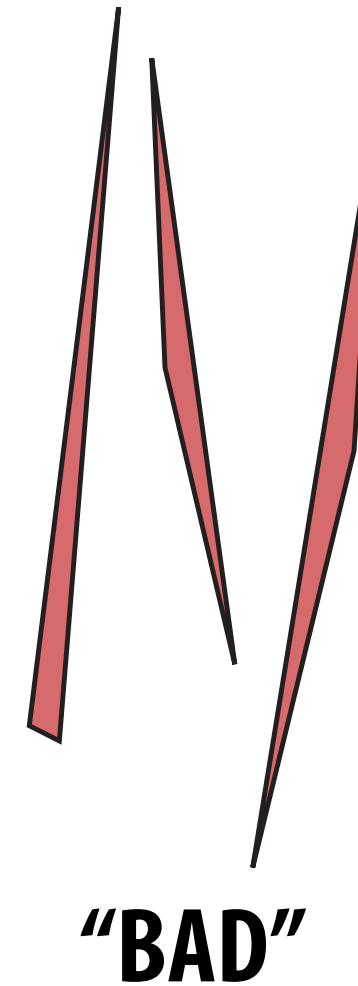
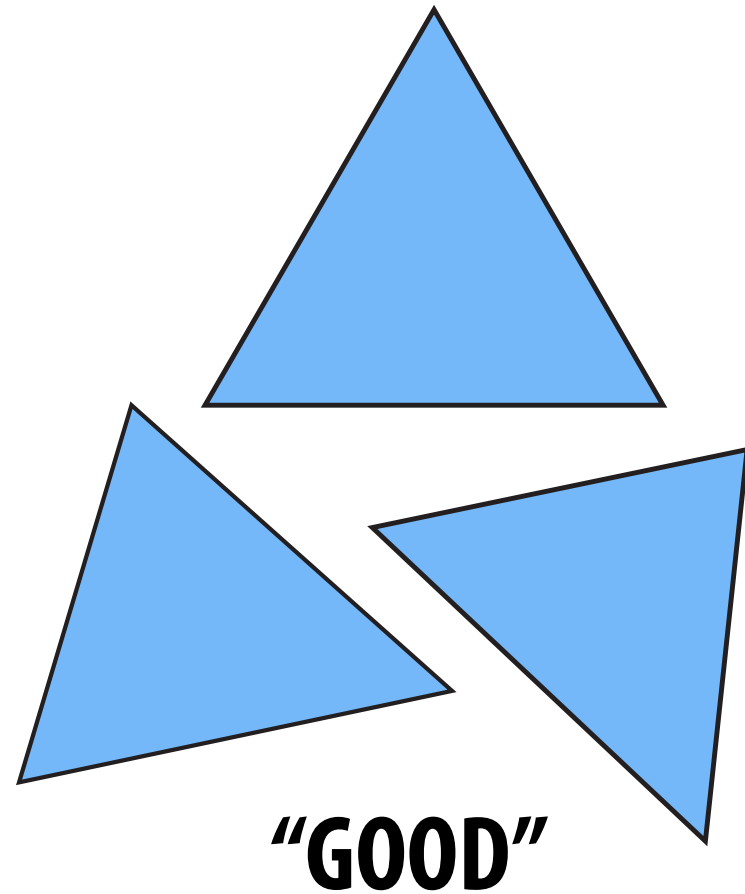
flattening of smooth cylinder & meshes



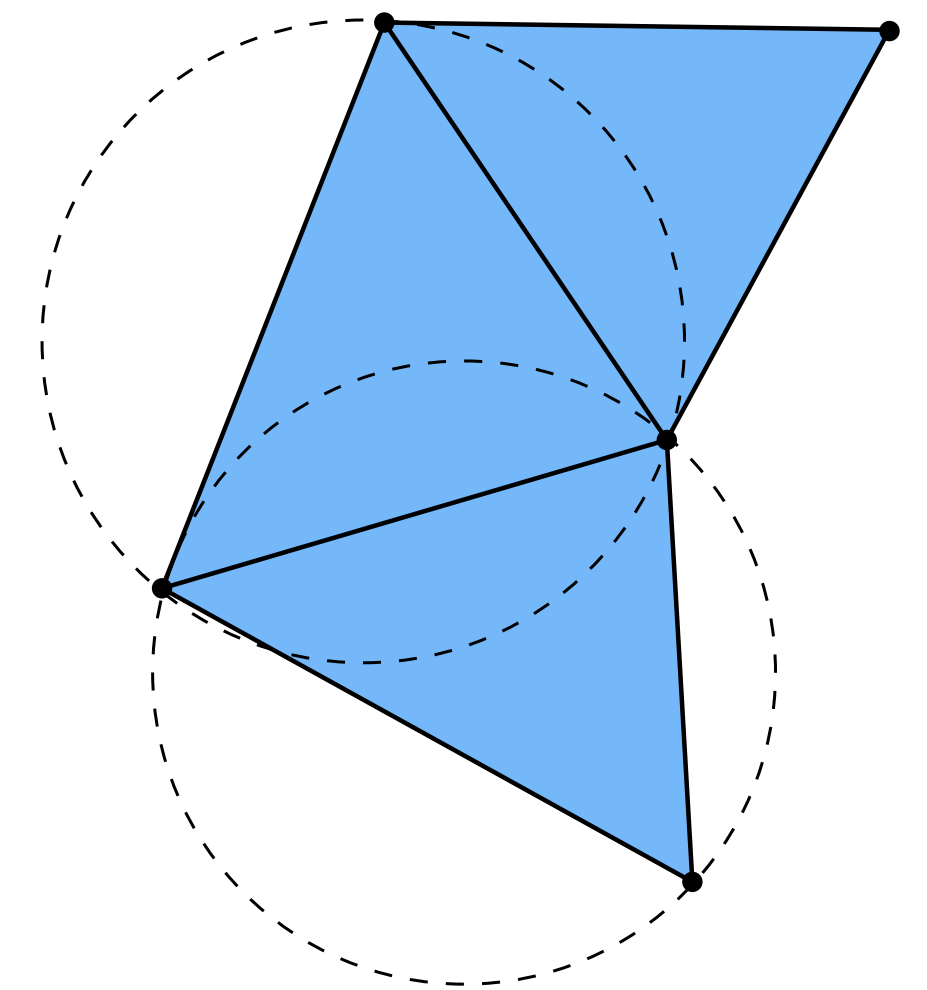
\*See Hildebrandt et al (2007), "On the convergence of metric and geometric properties of polyhedral surfaces"

# What else makes a “good” triangle mesh?

## ■ Another rule of thumb: triangle



DELAUNAY



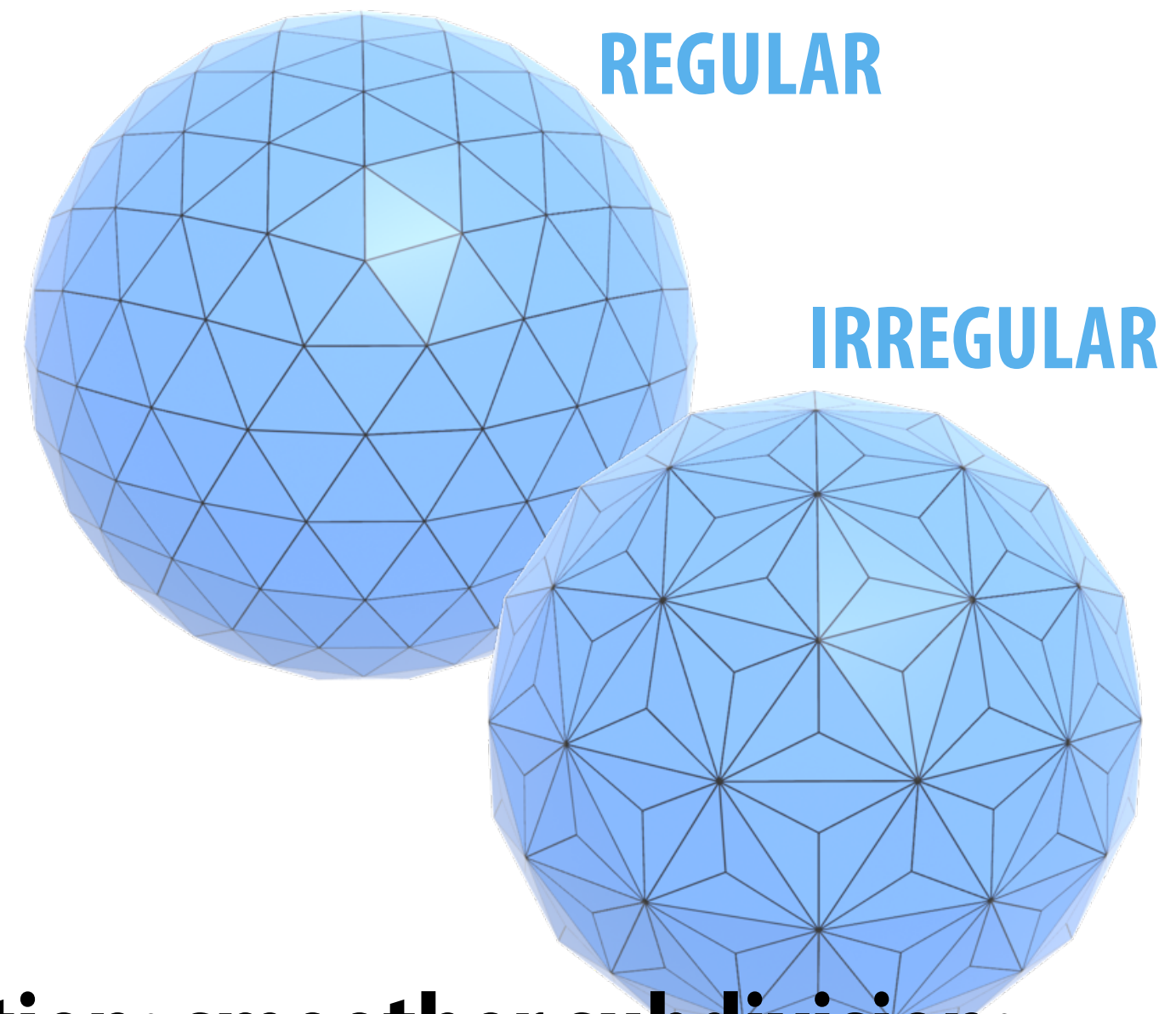
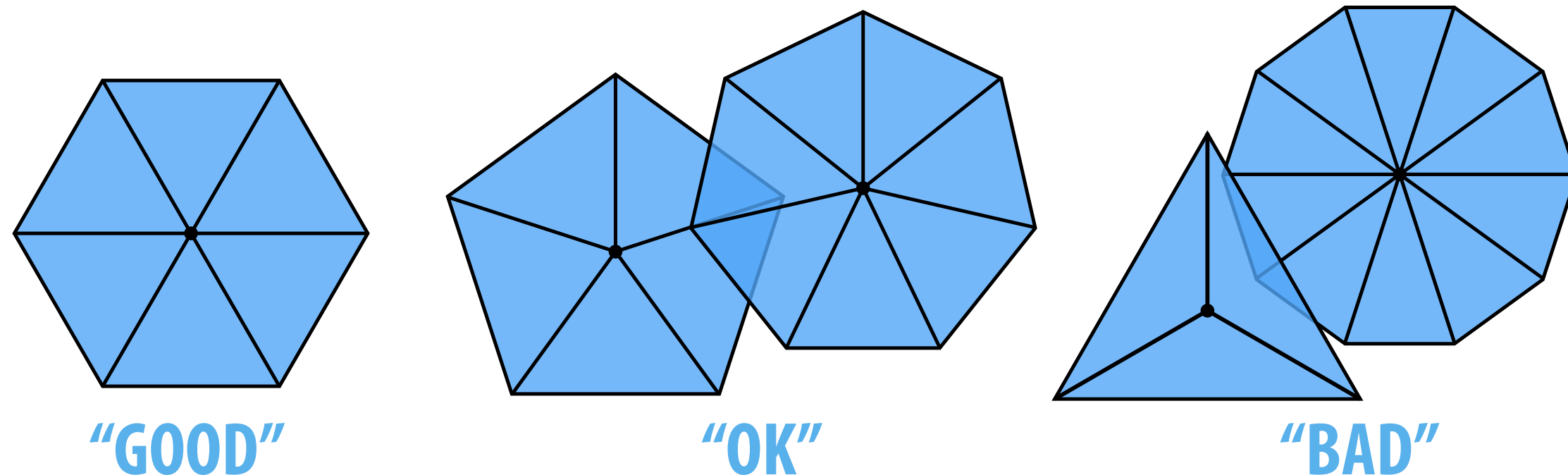
- E.g., all angles close to 60 degrees
- More sophisticated condition: Delaunay (empty circumcircles)
  - often helps with numerical accuracy/stability
  - coincides with shockingly many other desirable properties (maximizes minimum angle, provides smoothest interpolation, guarantees maximum principle...)
- Tradeoffs w/ good geometric approximation\*
  - e.g., long & skinny might be “more efficient”

\*see Shewchuk, “What is a Good Linear Element”

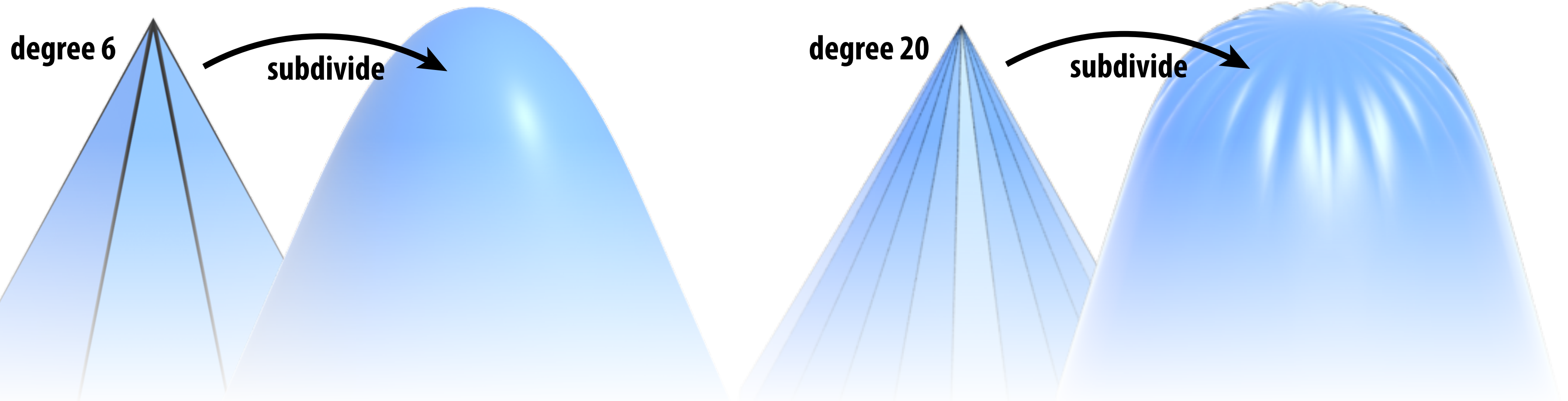


# What else constitutes a “good” mesh?

- Another rule of thumb: regular vertex degree
- Degree 6 for triangle mesh, 4 for quad mesh



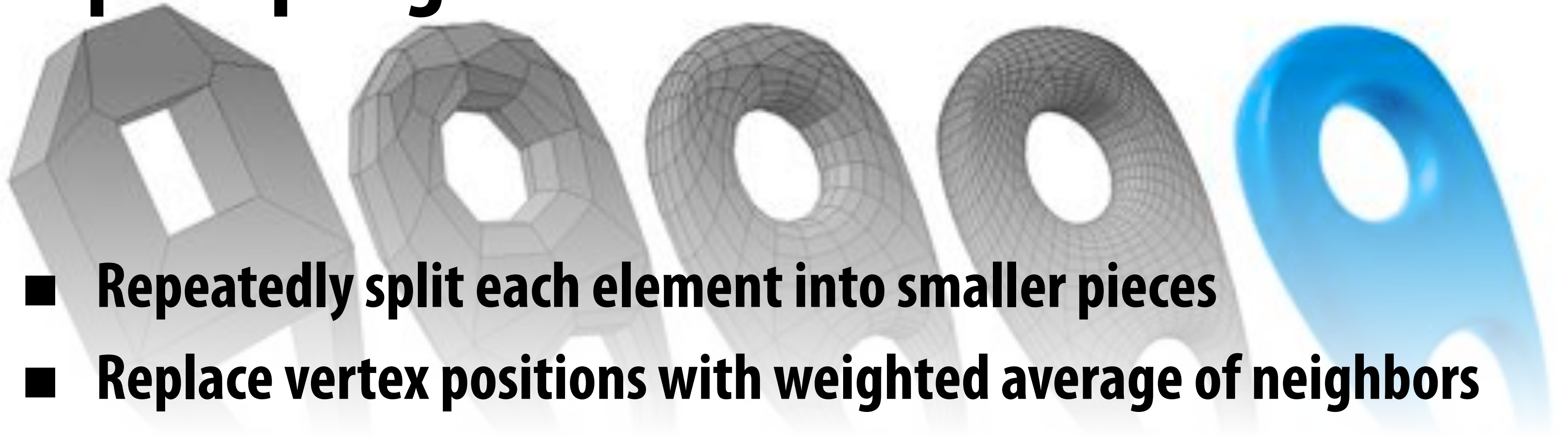
**Why? Better polygon shape; more regular computation; smoother subdivision:**



**Fact: in general, can't have regular vertex degree everywhere!**

# Upsampling via Subdivision

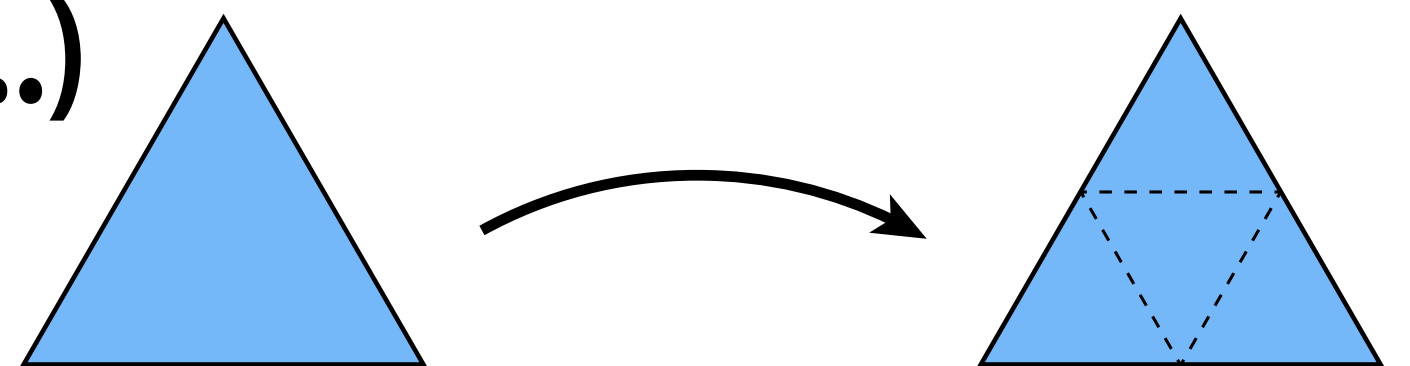
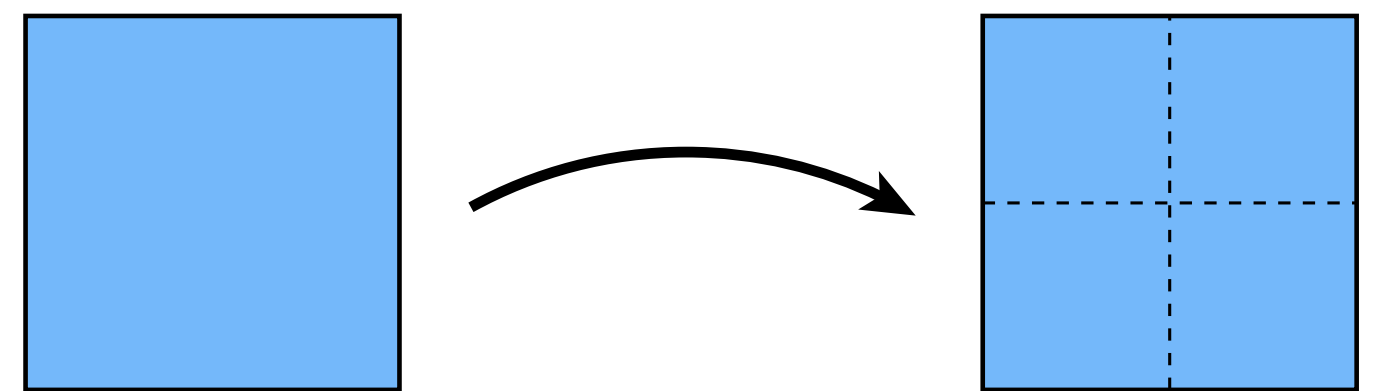
# Upsampling via Subdivision



- Repeatedly split each element into smaller pieces
- Replace vertex positions with weighted average of neighbors

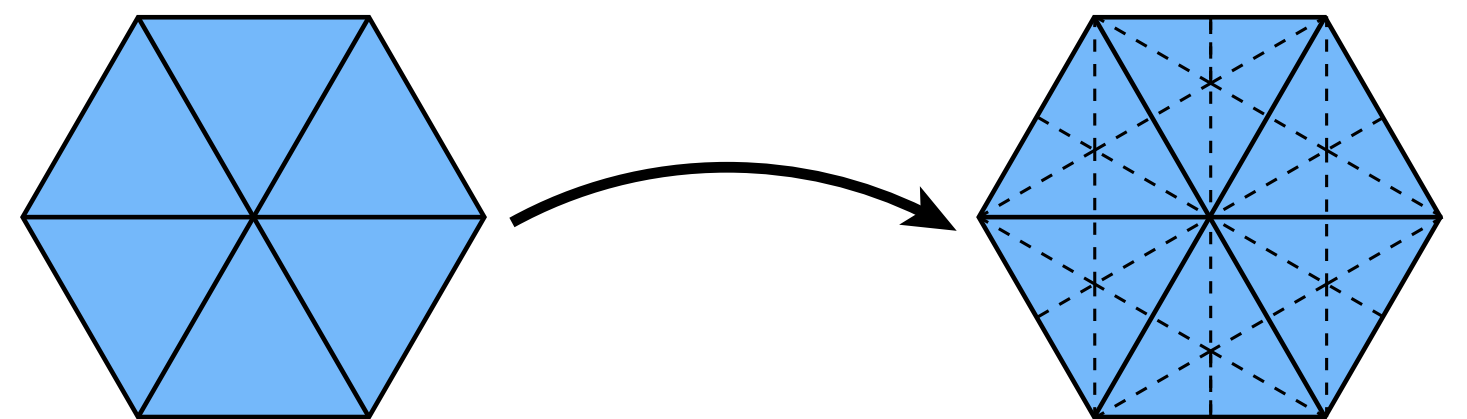
- Main considerations:

- interpolating vs. approximating
- limit surface continuity ( $C^1, C^2, \dots$ )
- behavior at irregular vertices



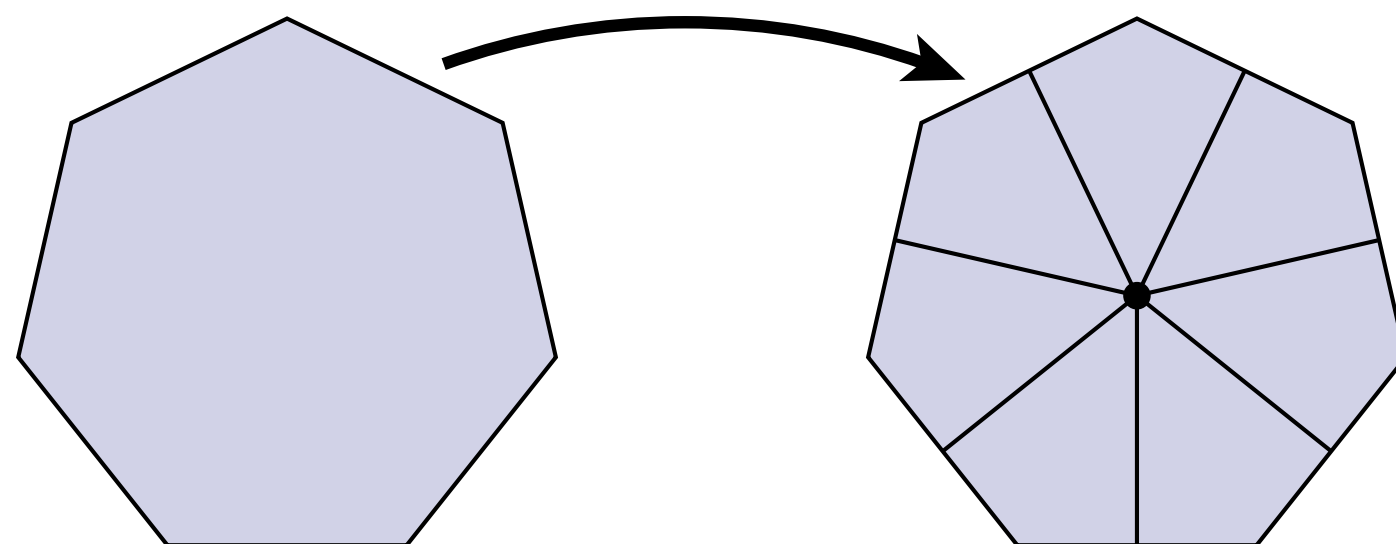
- Many options:

- Quad: Catmull-Clark
- Triangle: Loop, Butterfly, Sqrt(3)

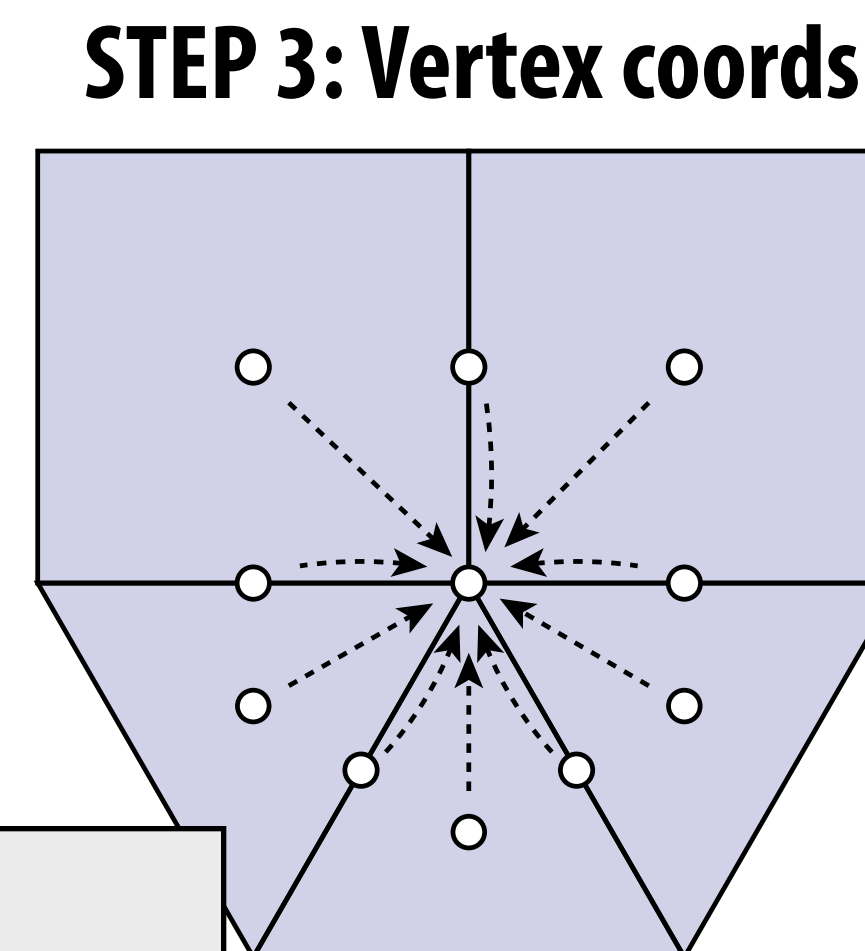
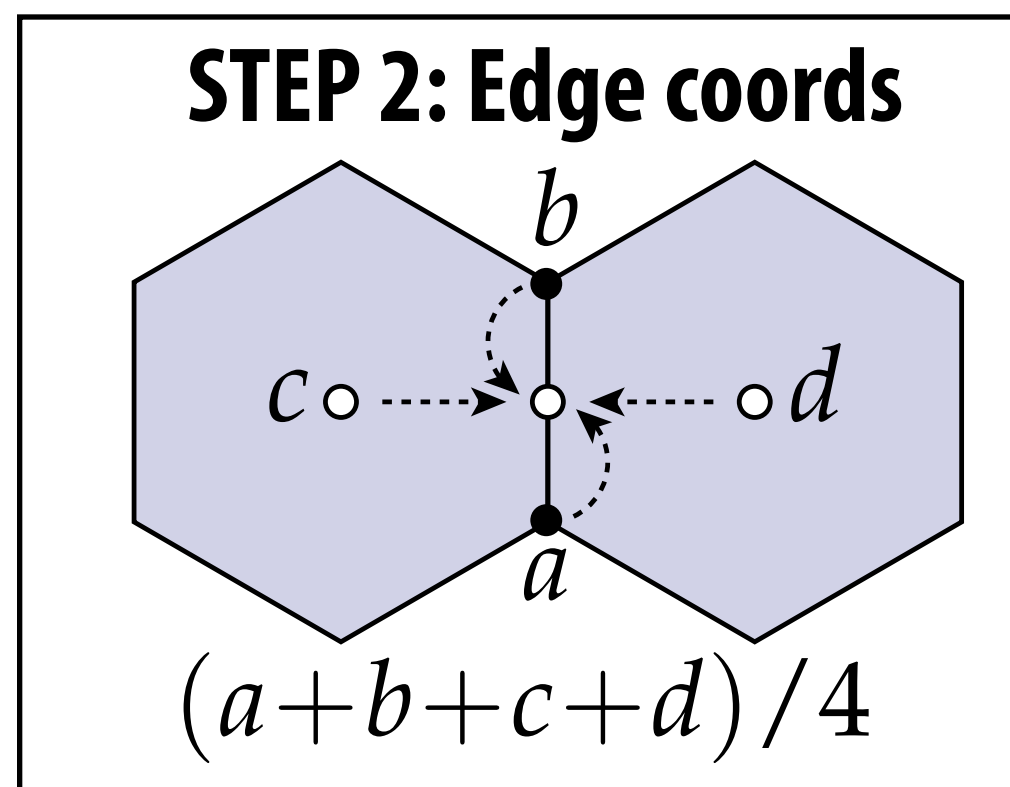
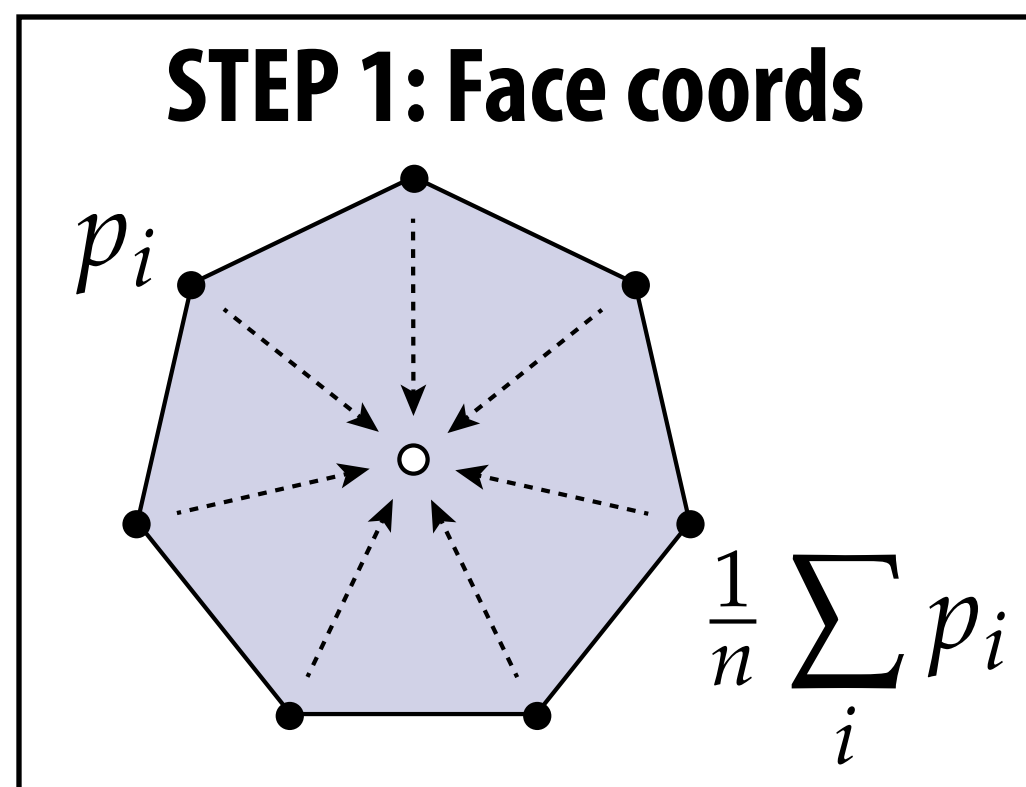


# Catmull-Clark Subdivision

- Step 0: split every polygon (any # of sides) into quadrilaterals:



- New vertex positions are weighted combination of old ones:



**New vertex coords:**

$$\frac{Q + 2R + (n - 3)S}{n}$$

$n$  - vertex degree

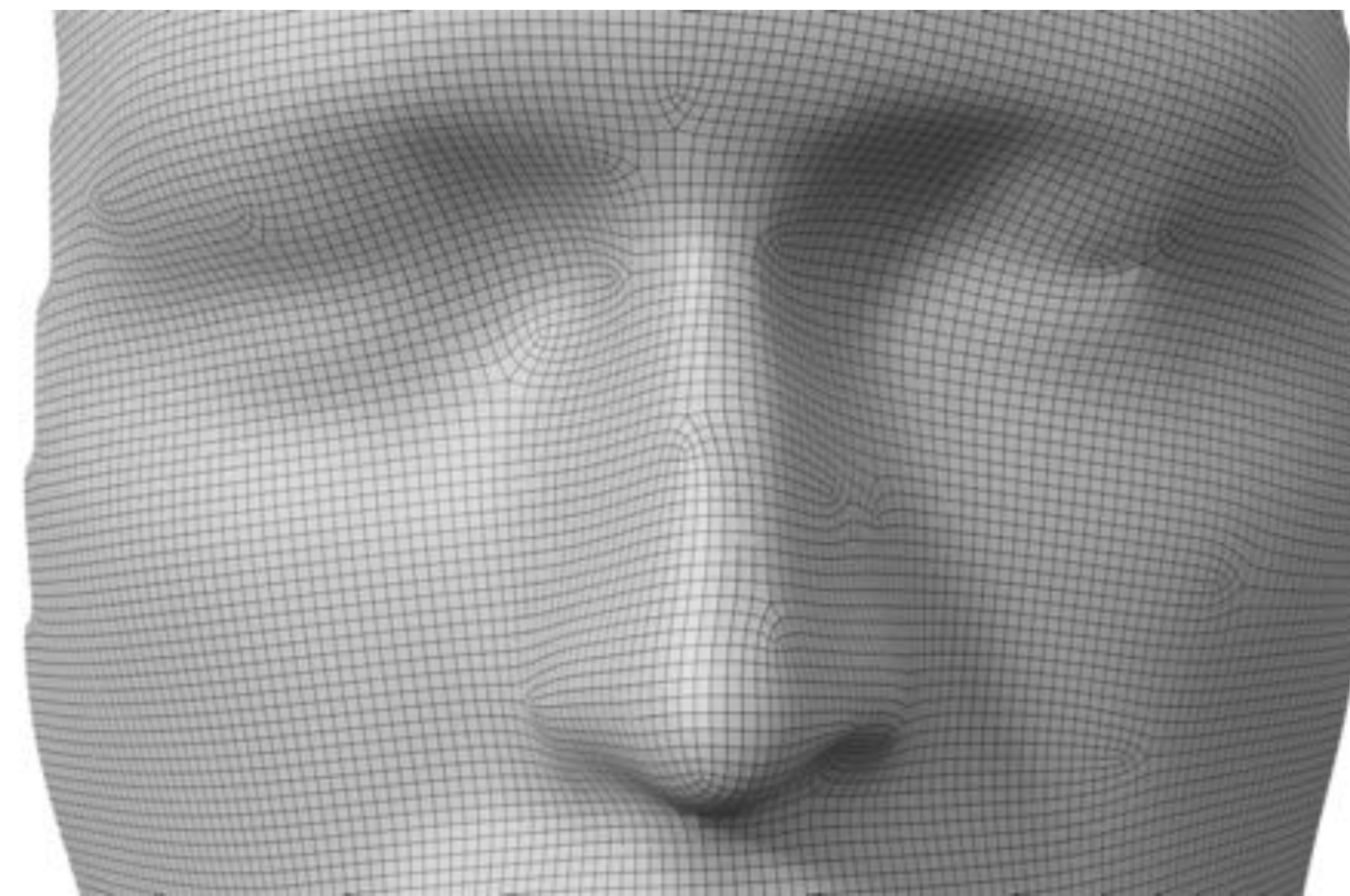
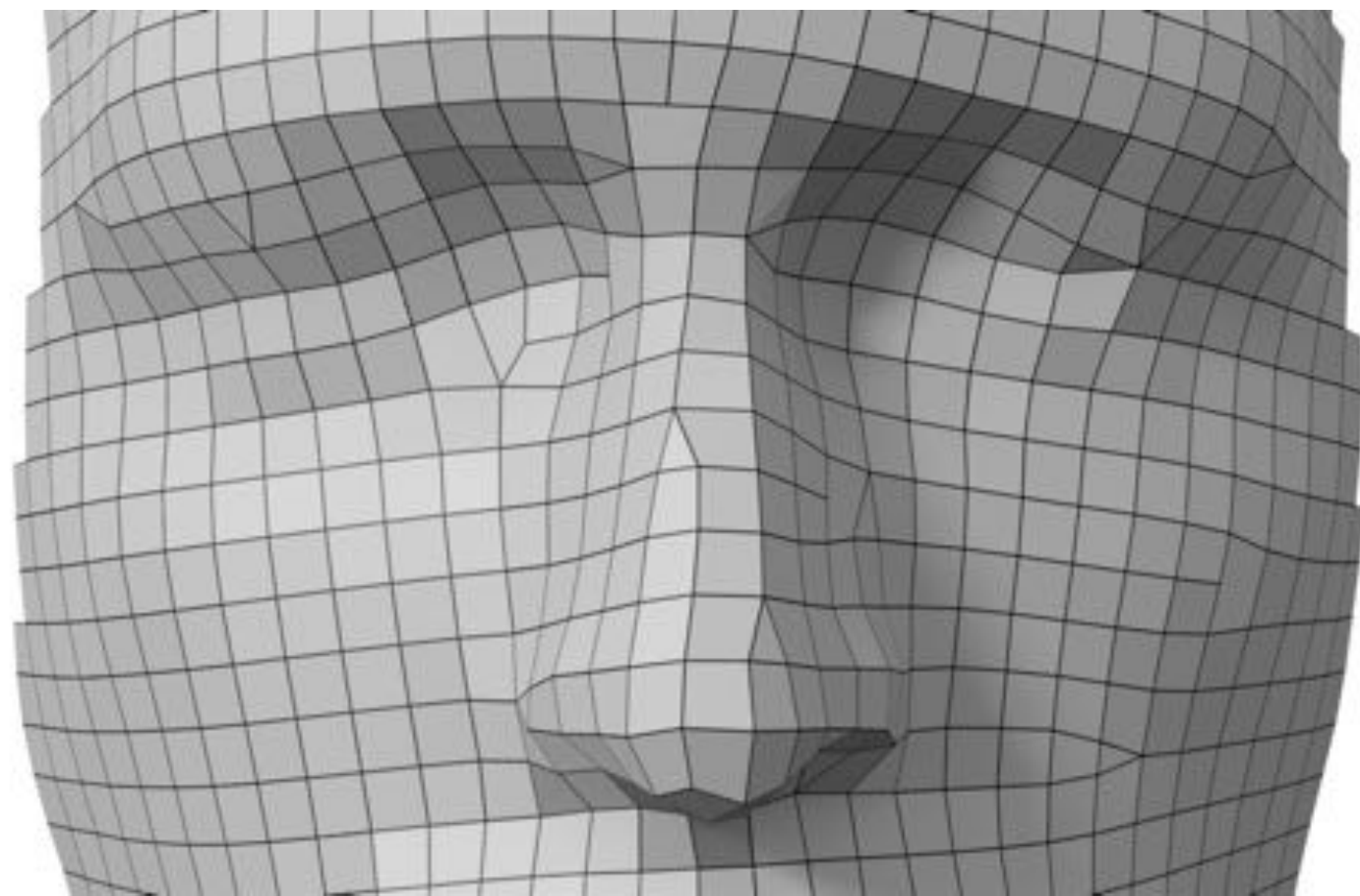
$Q$  - average of face coords around vertex

$R$  - average of edge coords around vertex

$S$  - original vertex position

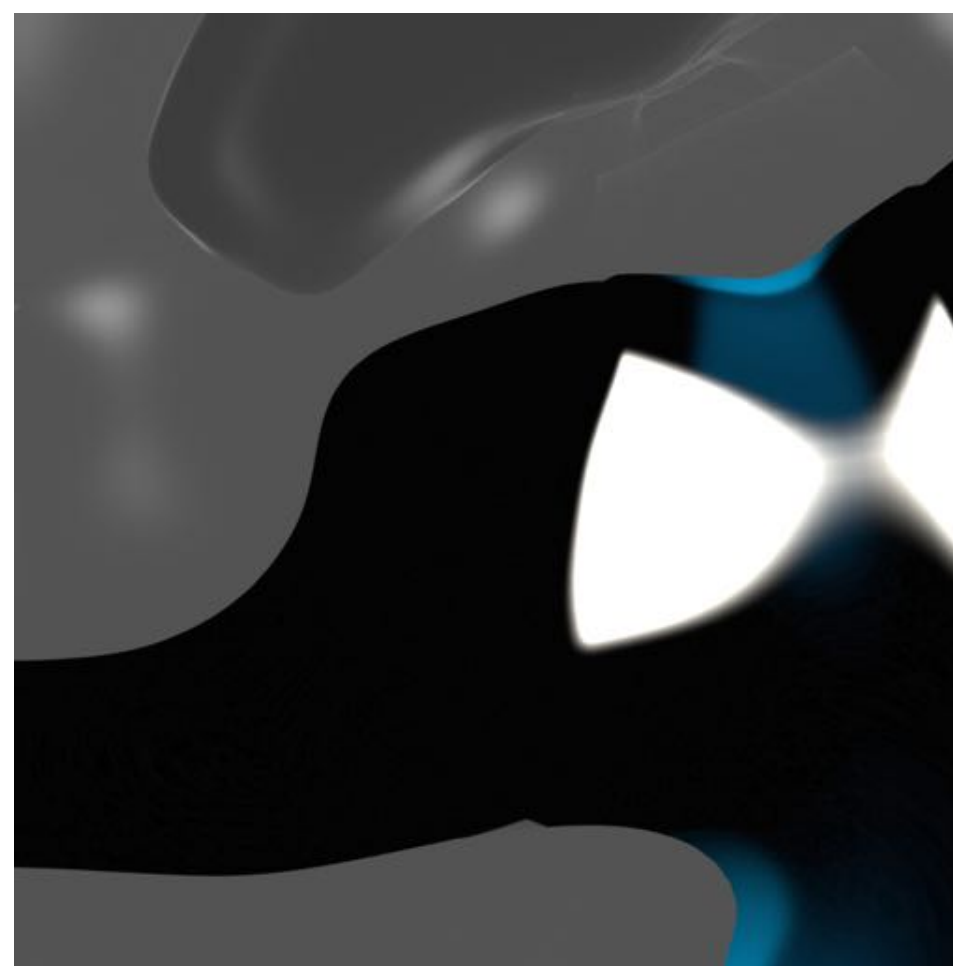


# Catmull-Clark on quad mesh

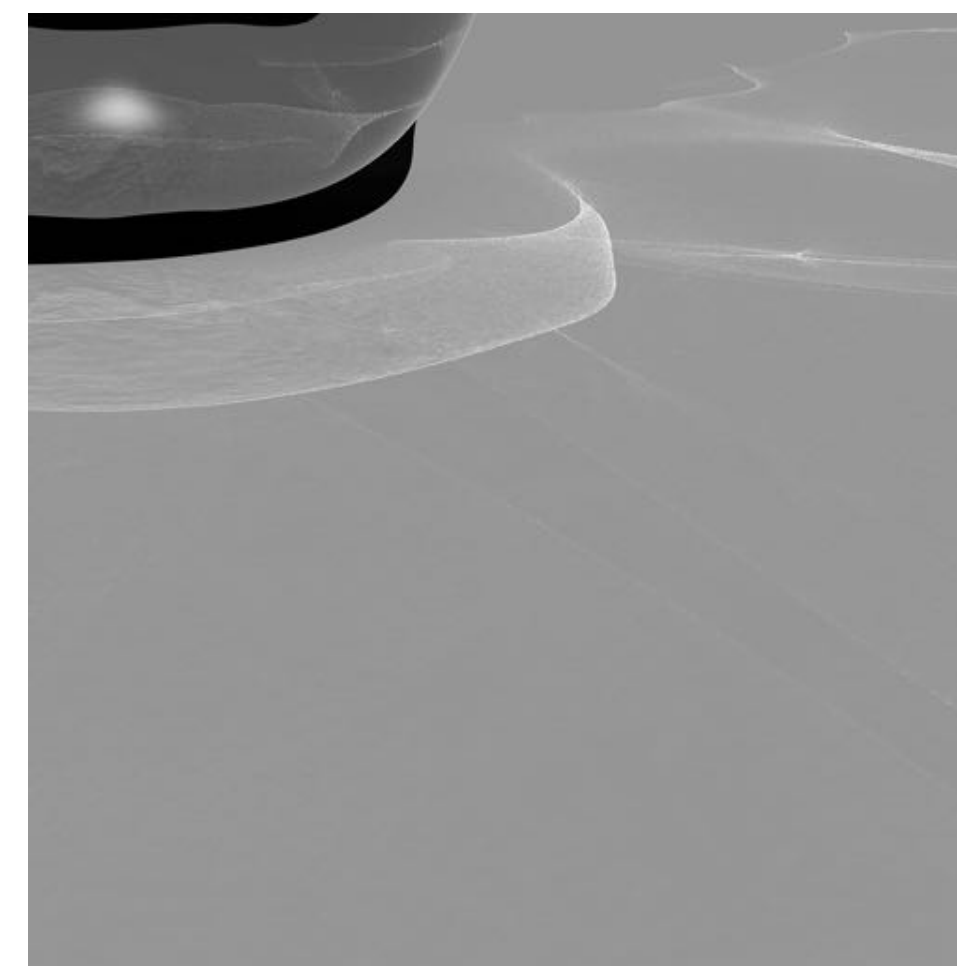


**few irregular vertices**

⇒ **smoothly-varying surface normals**



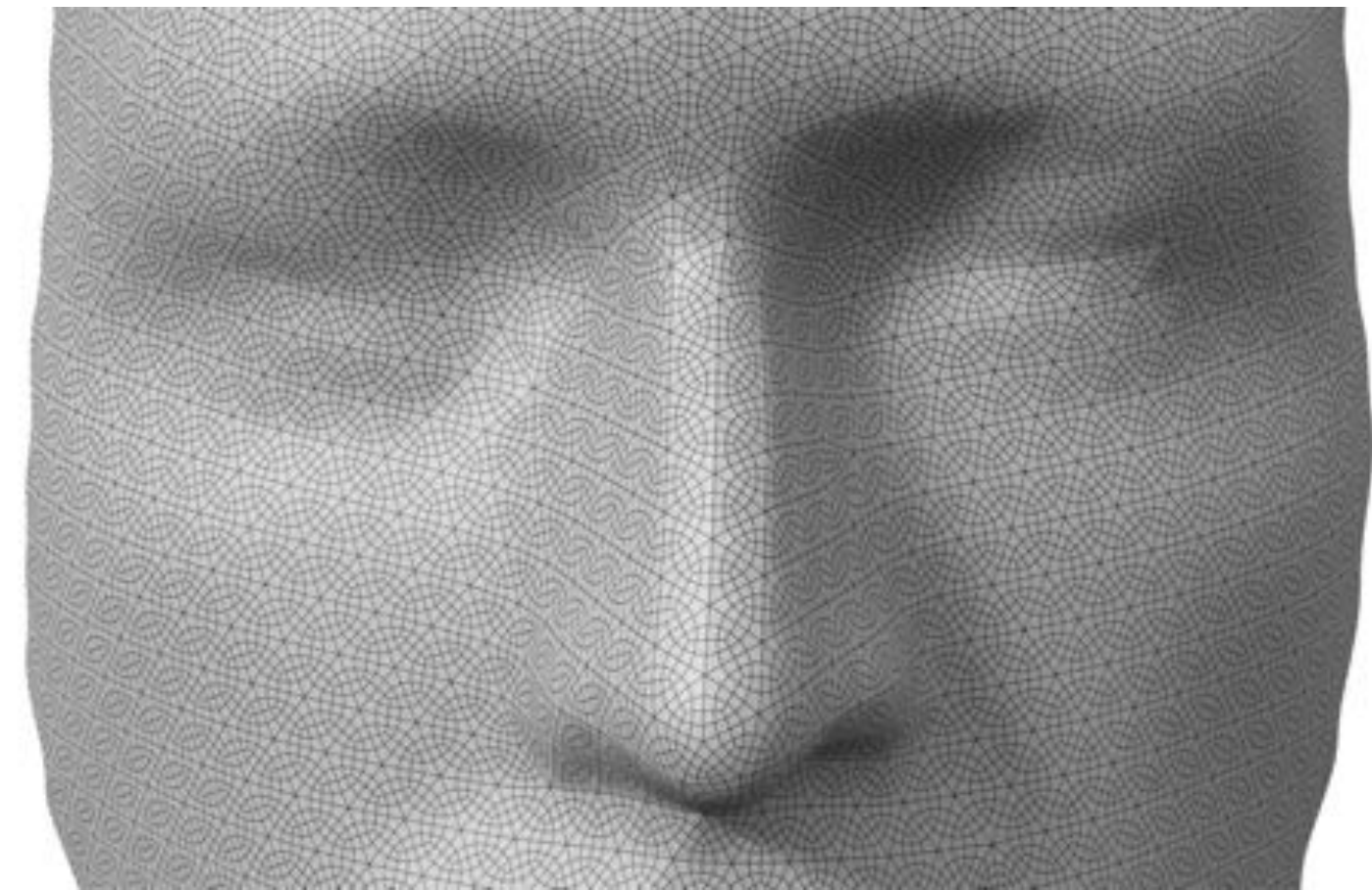
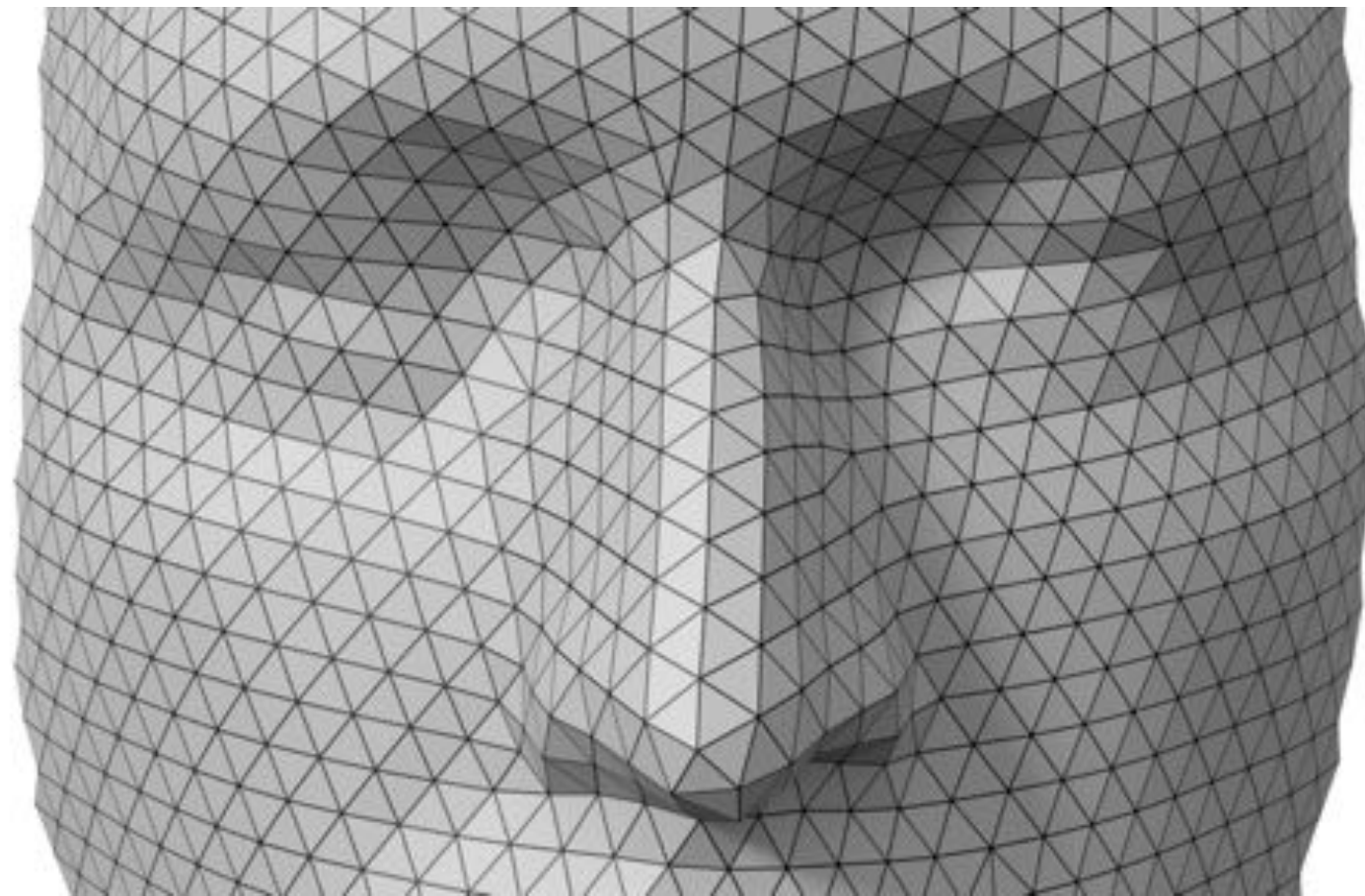
**smooth  
reflection lines**



**smooth  
caustics**

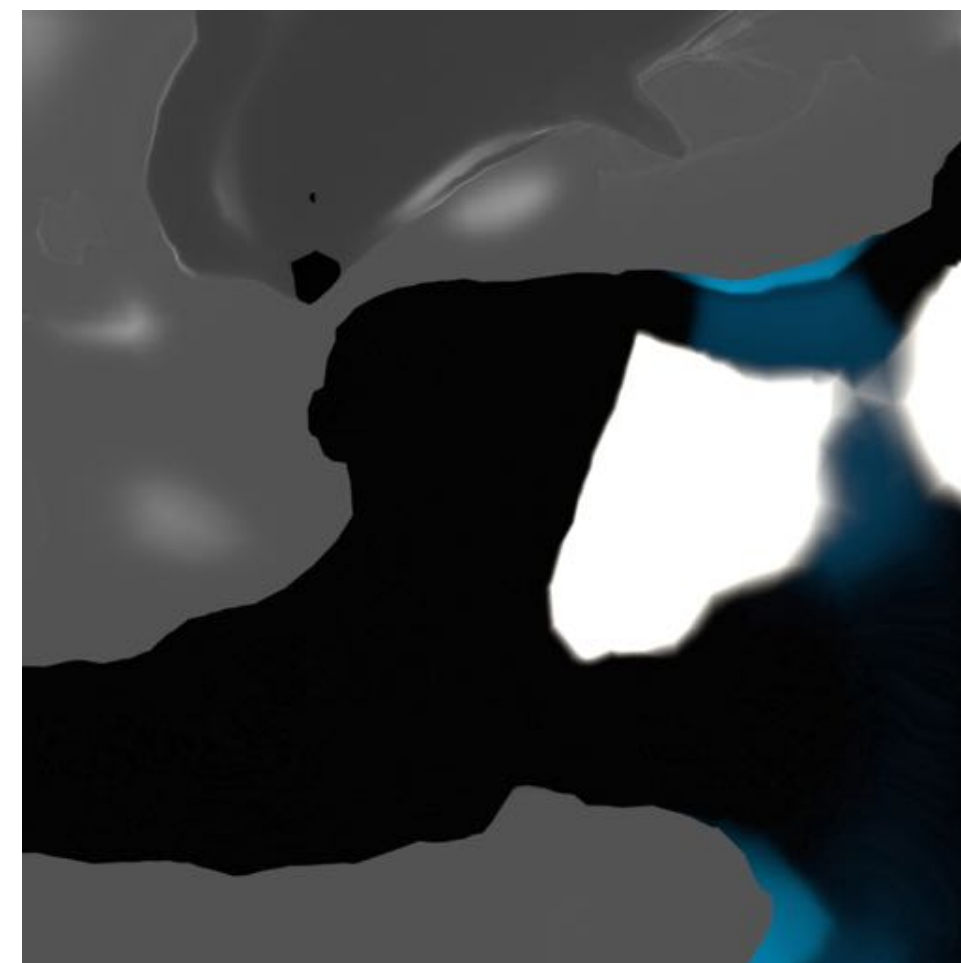


# Catmull-Clark on triangle mesh

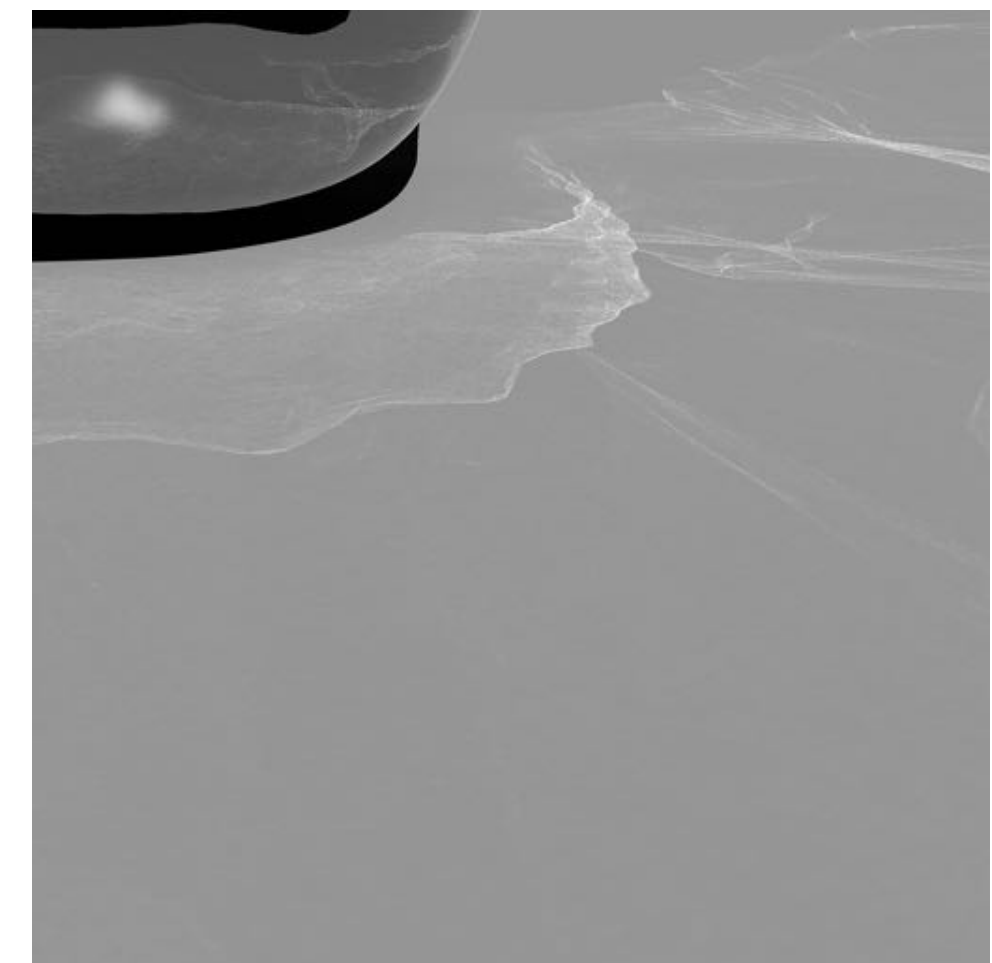


**many irregular vertices**

**⇒ erratic surface normals**



**jagged  
reflection lines**



**jagged  
caustics**

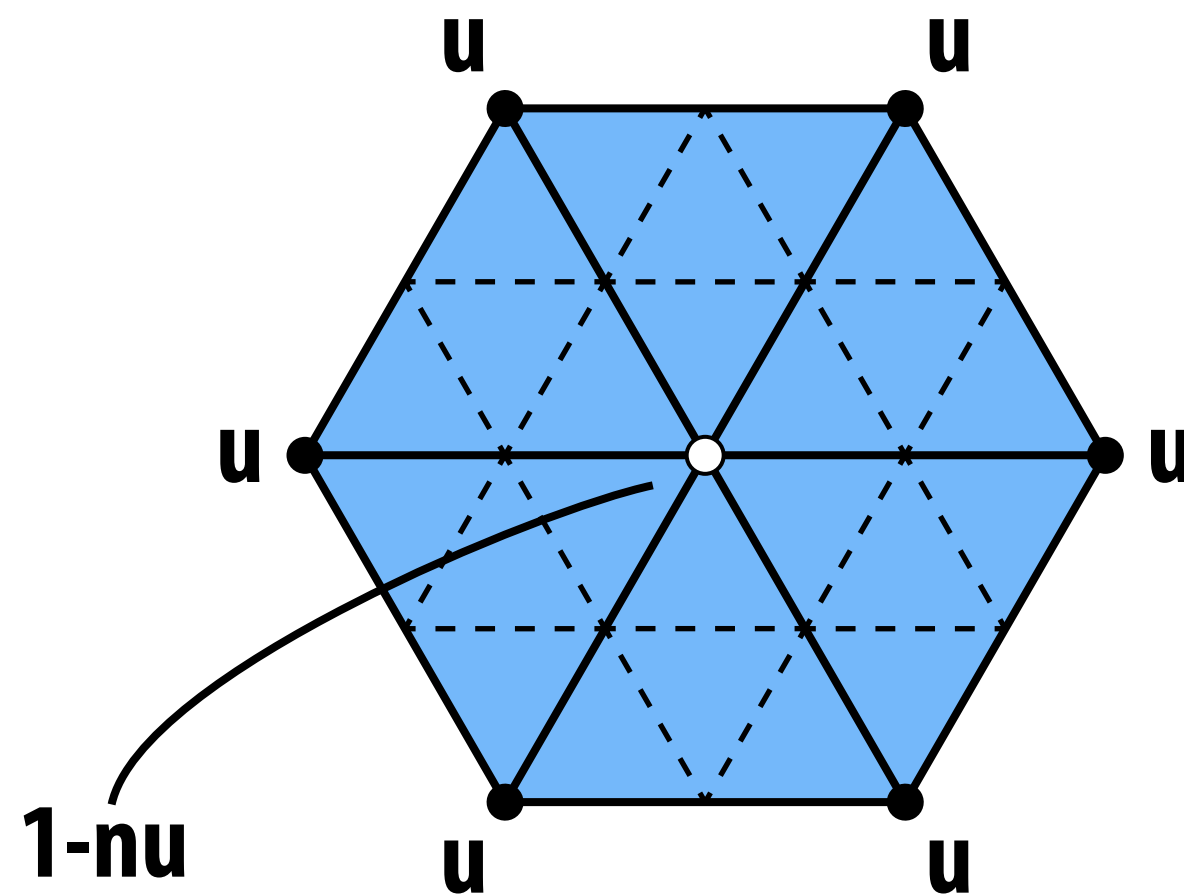
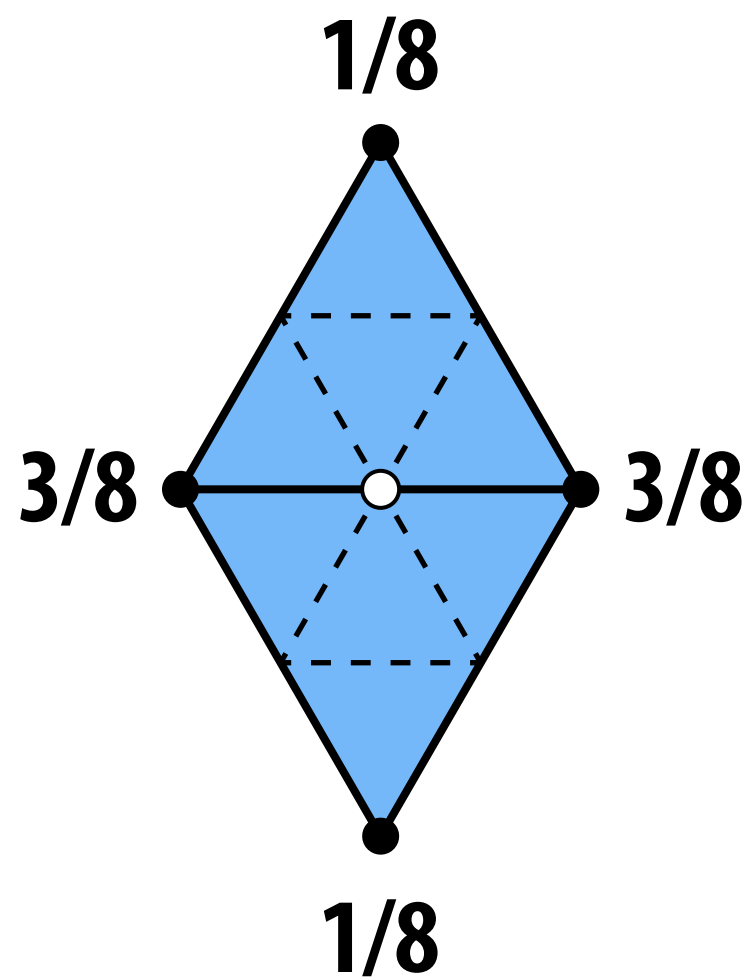
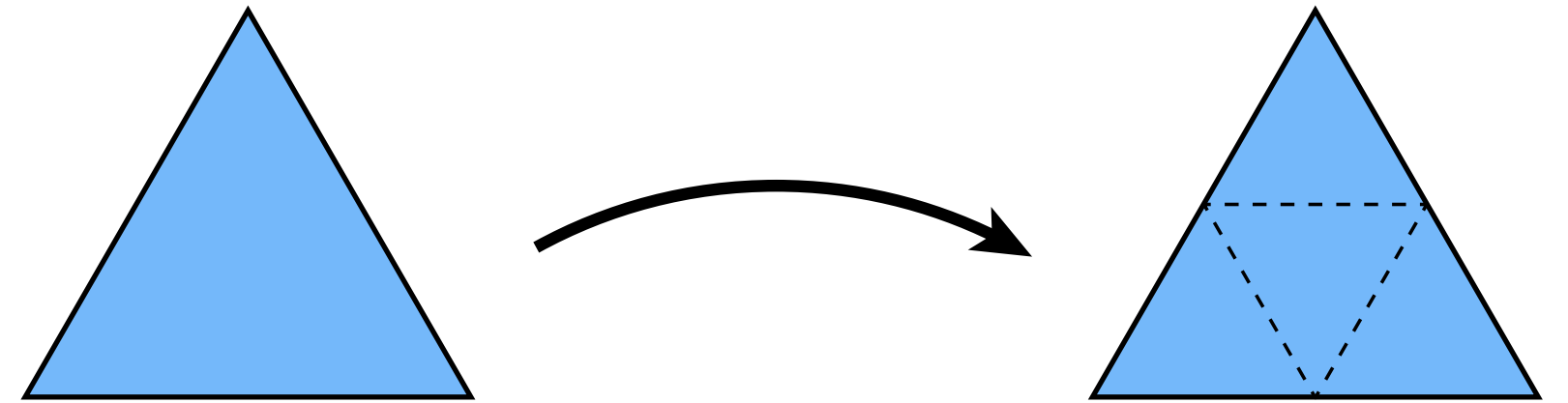


# Loop Subdivision

- Alternative subdivision scheme for triangle meshes
- Curvature is continuous away from irregular vertices (" $C^2$ ")

- **Algorithm:**

- Split each triangle into four
- Assign new vertex positions according to weights:

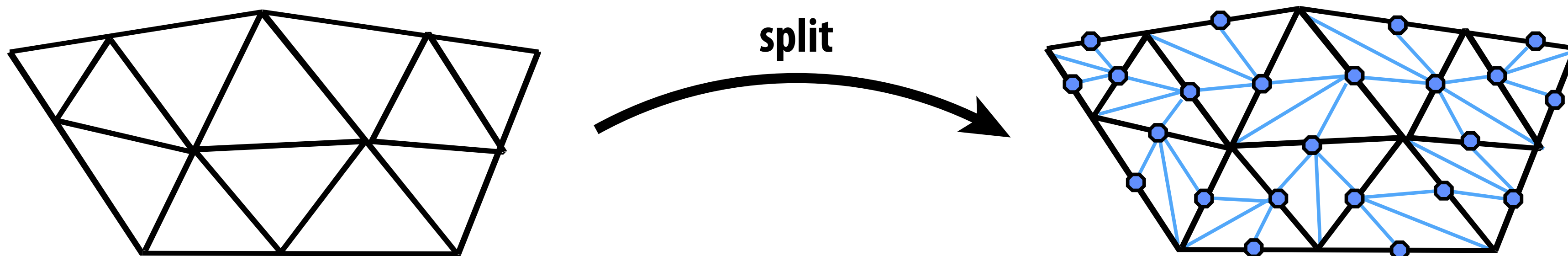


$n$ : vertex degree

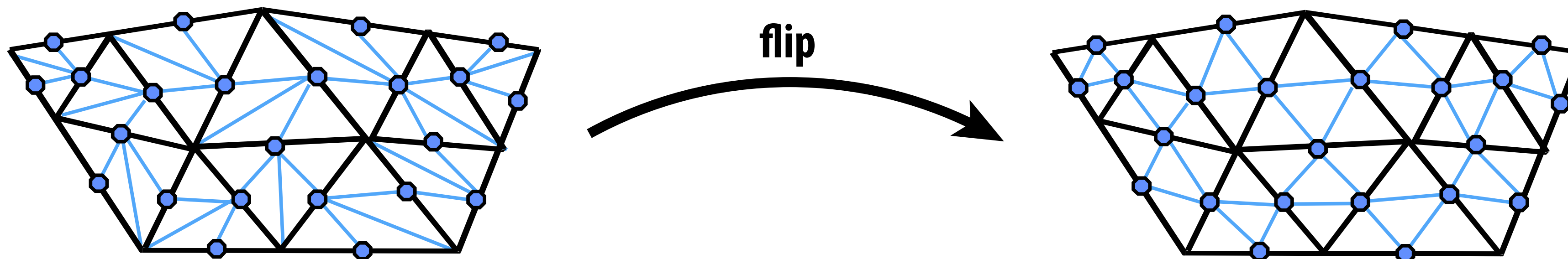
$u$ :  $3/16$  if  $n=3$ ,  $3/(8n)$  otherwise

# Loop Subdivision via Edge Operations

- First, split edges of original mesh in any order:



- Next, flip new edges that touch a new & old vertex:



**(Don't forget to update vertex positions!)**

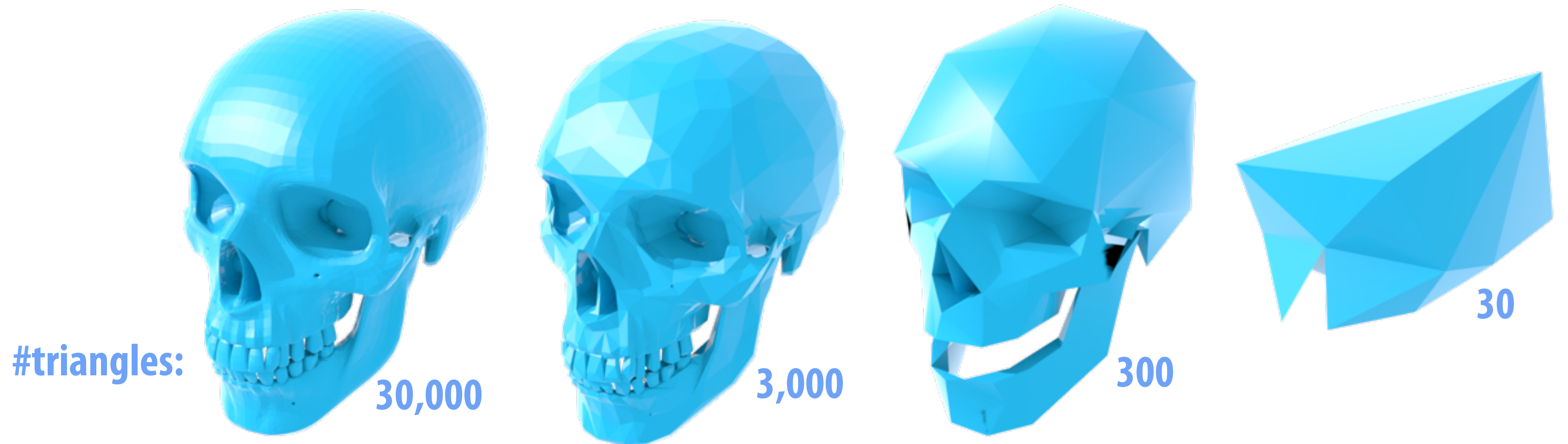


# **Downsampling**

**(i.e., what if we want fewer triangles?)**

# Simplification via Edge Collapse

- One popular scheme: iteratively collapse edges
- Greedy algorithm:
  - assign each edge a cost
  - collapse edge with least cost
  - repeat until target number of elements is reached
- Particularly effective cost function: quadric error metric\*

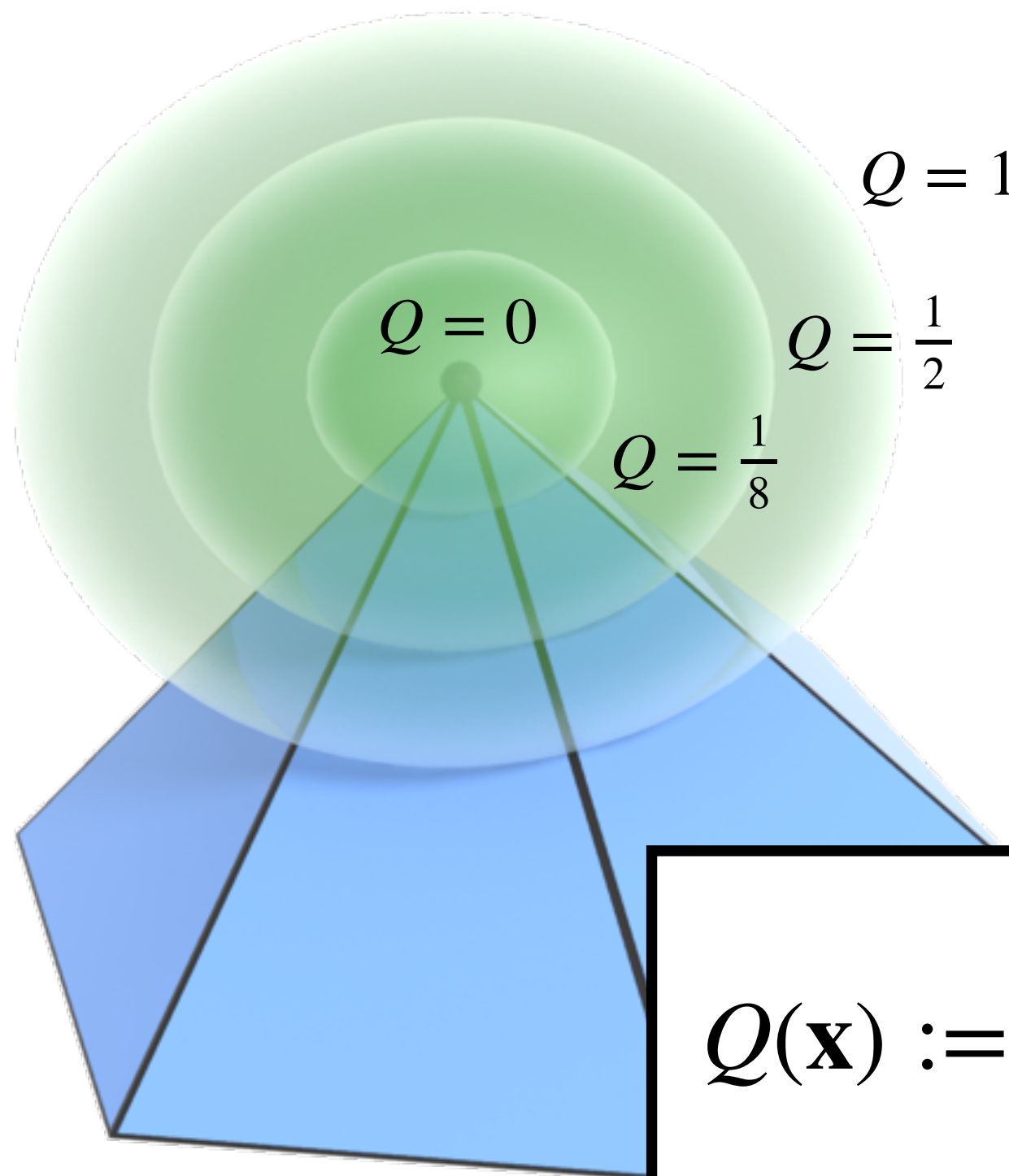
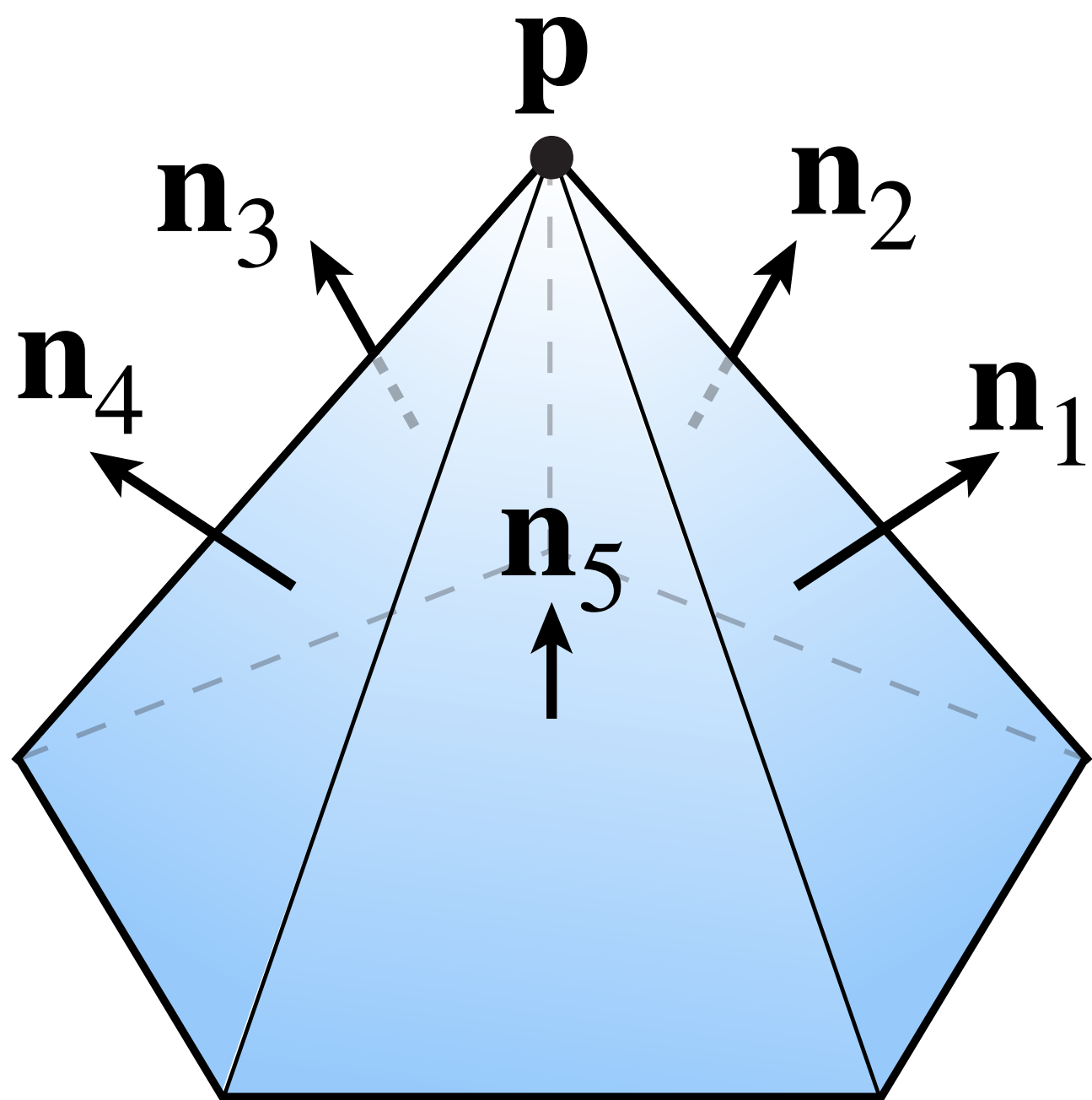
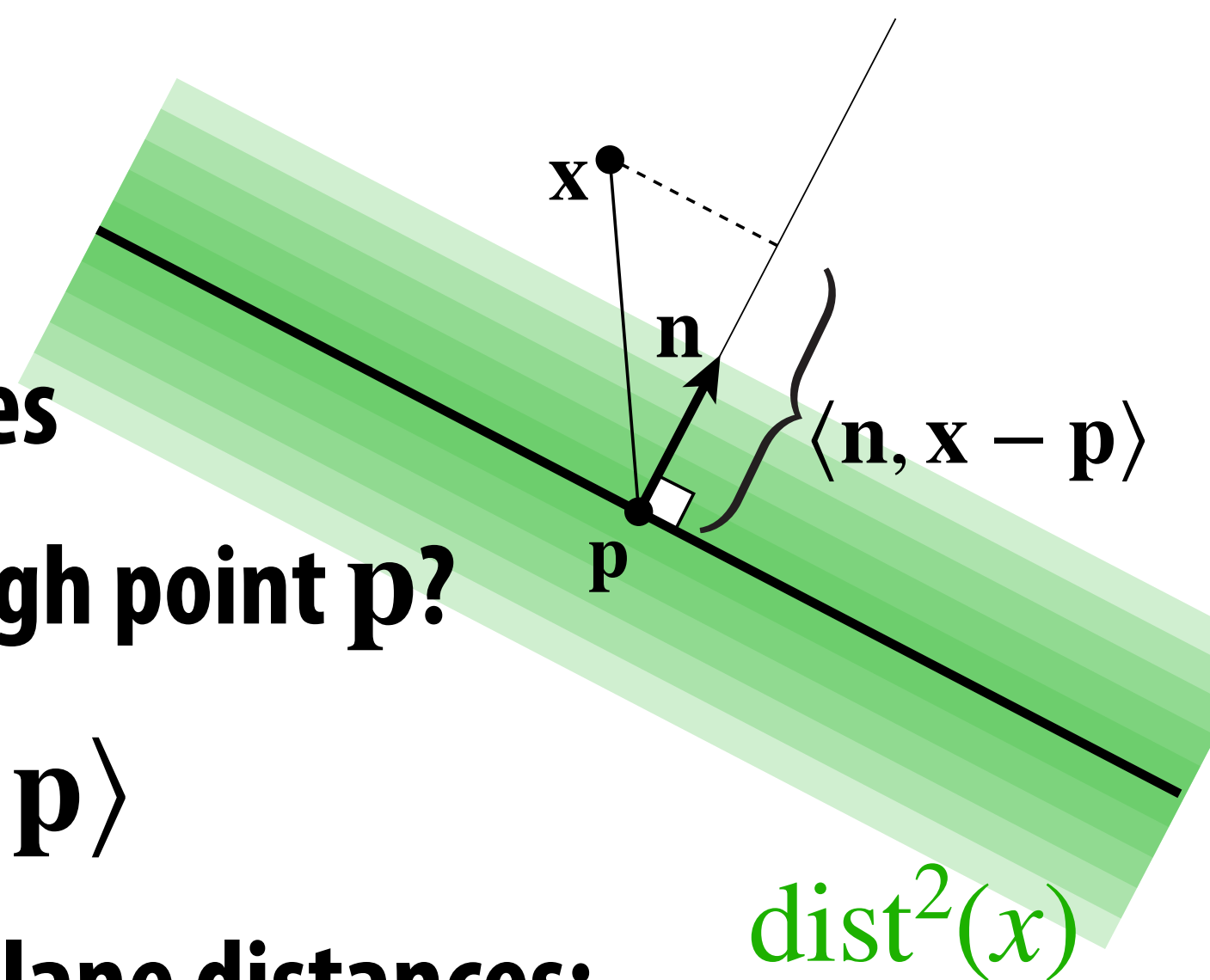


\*invented at CMU (Garland & Heckbert 1997)



# Quadric Error Metric

- Approximate distance to a collection of triangles
- $Q$ : Distance to plane w/ normal  $\mathbf{n}$  passing through point  $\mathbf{p}$ ?
- A:  $\text{dist}(\mathbf{x}) = \langle \mathbf{n}, \mathbf{x} \rangle - \langle \mathbf{n}, \mathbf{p} \rangle = \langle \mathbf{n}, \mathbf{x} - \mathbf{p} \rangle$
- Quadric error is then sum of squared point-to-plane distances:

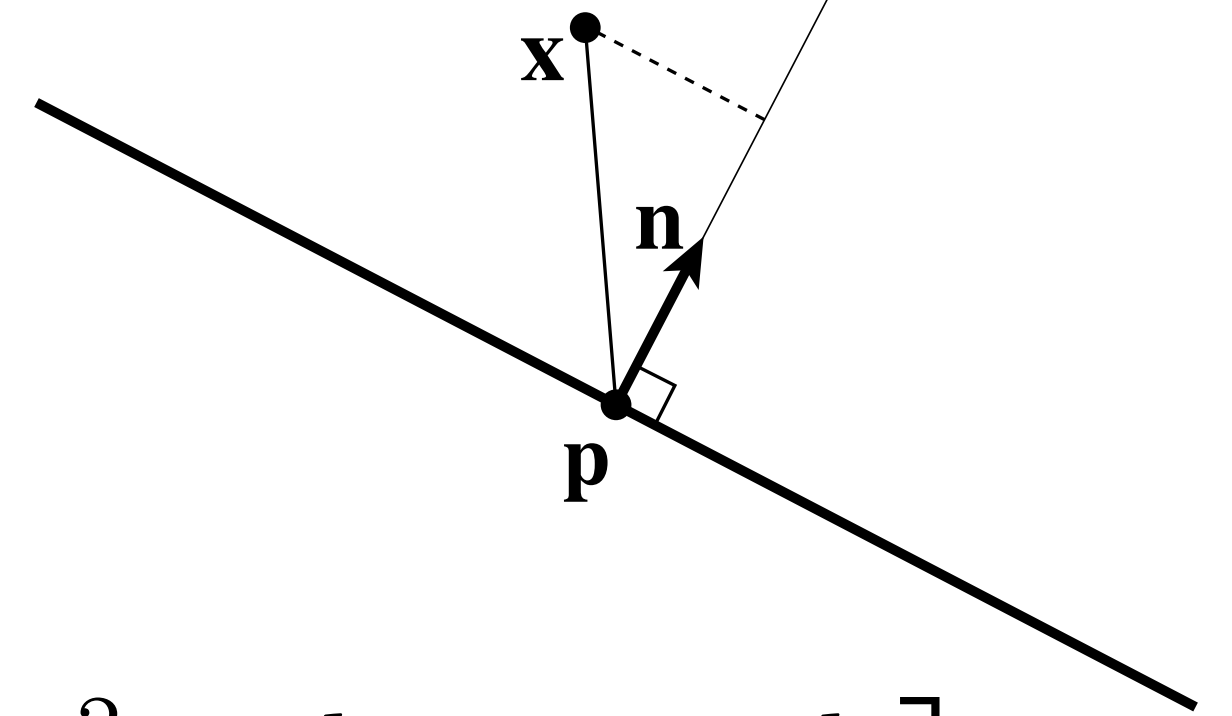


$$Q(\mathbf{x}) := \sum_{i=1}^k \langle \mathbf{n}_i, \mathbf{x} - \mathbf{p} \rangle^2$$

# Quadric Error - Homogeneous Coordinates

- Suppose in coordinates we have

- a query point  $\mathbf{x} = (x, y, z)$
- a normal  $\mathbf{n} = (a, b, c)$
- an offset  $d := -\langle \mathbf{n}, \mathbf{p} \rangle$



- In homogeneous coordinates, let

- $\mathbf{u} := (x, y, z, 1)$
- $\mathbf{v} := (a, b, c, d)$

$$K = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$

- Signed distance to plane is then just  $\langle \mathbf{u}, \mathbf{v} \rangle = ax + by + cz + d$

- Squared distance is  $\langle \mathbf{u}, \mathbf{v} \rangle^2 = \mathbf{u}^\top (\mathbf{v}\mathbf{v}^\top) \mathbf{u} =: \mathbf{u}^\top K \mathbf{u}$

- Matrix  $K = \mathbf{v}\mathbf{v}^\top$  encodes squared distance to plane

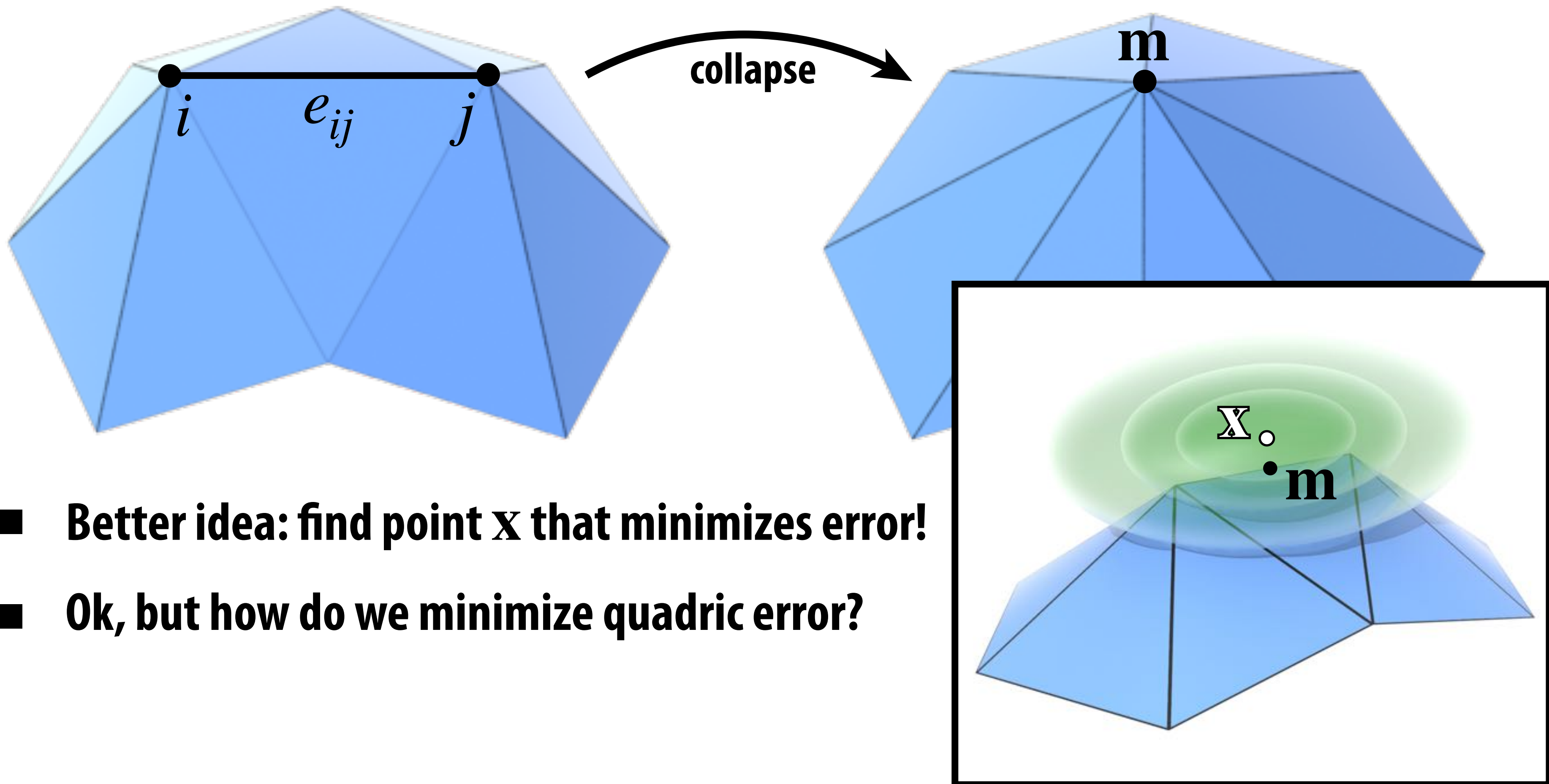
**Key idea: sum of matrices  $K \iff$  distance to union of planes**

$$\mathbf{u}^\top K_1 \mathbf{u} + \mathbf{u}^\top K_2 \mathbf{u} = \mathbf{u}^\top (K_1 + K_2) \mathbf{u}$$



# Quadric Error of Edge Collapse

- How much does it cost to collapse an edge  $e_{ij}$ ?
- Idea: compute midpoint  $\mathbf{m}$ , measure error  $Q(\mathbf{m}) = \mathbf{m}^T (K_i + K_j) \mathbf{m}$
- Error becomes "score" for  $e_{ij}$ , determining priority



- Better idea: find point  $\mathbf{x}$  that minimizes error!
- Ok, but how do we minimize quadric error?

# Review: Minimizing a Quadratic Function

■ Suppose you have a function  $f(x) = ax^2 + bx + c$

■ **Q: What does the graph of this function look like?**

■ Could also look like this!

■ **Q: How do we find the minimum?**

■ **A: Find where the function looks “flat” if we zoom in really close**

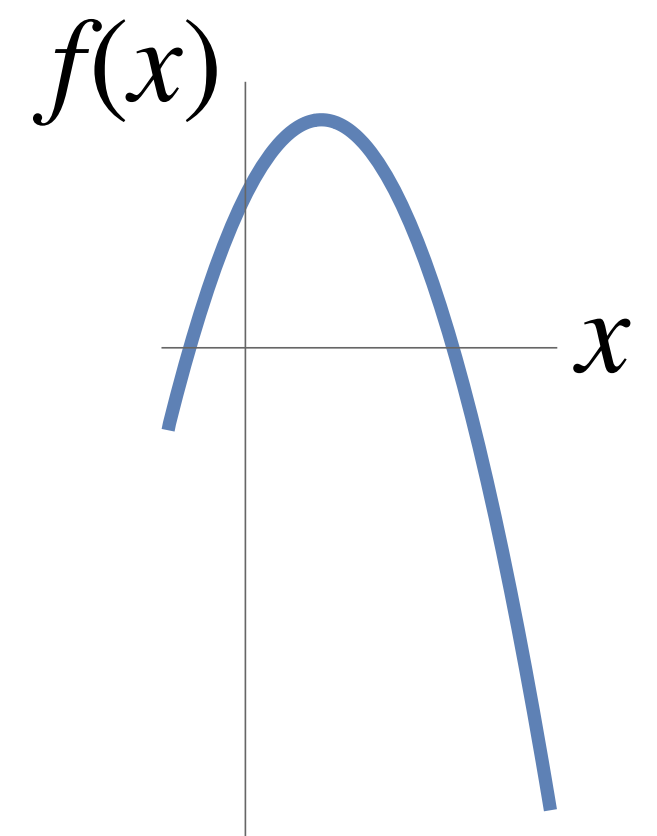
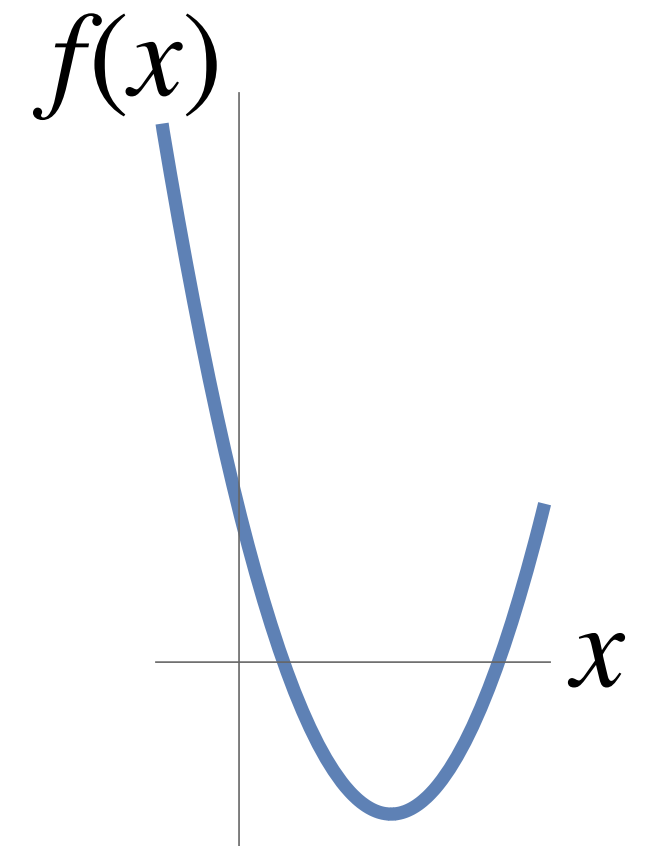
■ **I.e., find point  $x$  where 1st derivative vanishes:**

$$f'(x) = 0$$

$$2ax + b = 0$$

$$x = -b/2a$$

**(What does  $x$  describe for the second function?)**



# Minimizing Quadratic Polynomial

- Not much harder to minimize a quadratic polynomial in  $n$  variables
- Can always write in terms of a symmetric matrix  $A$
- E.g., in 2D:  $f(x, y) = ax^2 + bxy + cy^2 + dx + ey + g$

$$\mathbf{x} = \begin{bmatrix} x \\ y \end{bmatrix} \quad A = \begin{bmatrix} a & b/2 \\ b/2 & c \end{bmatrix} \quad \mathbf{u} = \begin{bmatrix} d \\ e \end{bmatrix}$$

$$f(x, y) = \mathbf{x}^T A \mathbf{x} + \mathbf{u}^T \mathbf{x} + g$$

(will have this same form for any  $n$ )

- **Q: How do we find a critical point (min/max/saddle)?**

- **A: Set derivative to zero!**

$$2A\mathbf{x} + \mathbf{u} = 0$$

$$\mathbf{x} = -\frac{1}{2}A^{-1}\mathbf{u}$$

(compare with  
our 1D solution)

$$x = -b/2a$$

(Can you show this is true, at least in 2D?)

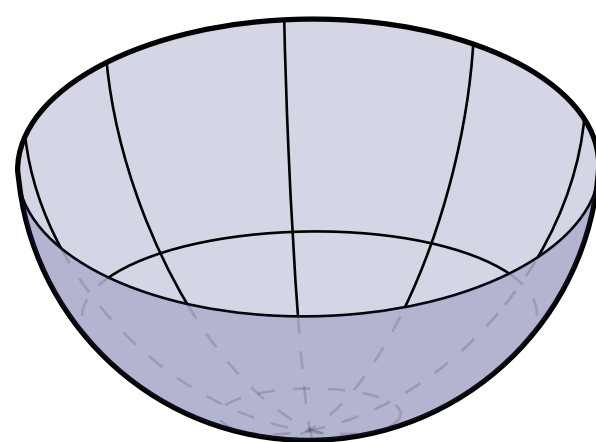


# Positive Definite Quadratic Form

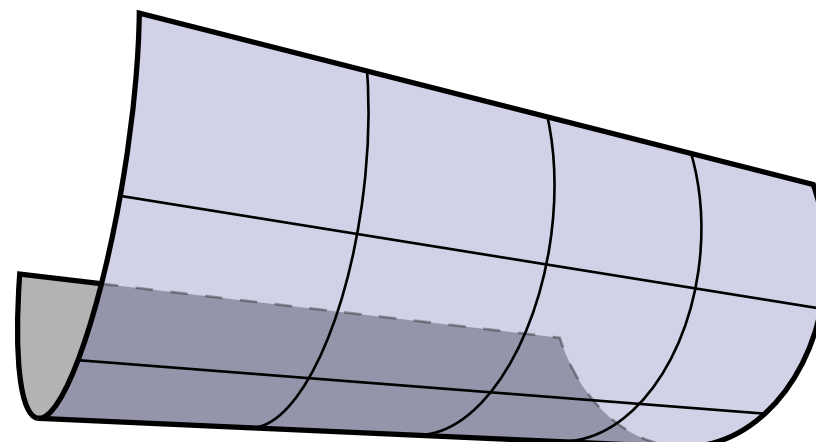
- Just like our 1D parabola, critical point is not always a min!
- **Q: In 2D, 3D, nD, when do we get a minimum?**
- **A: When matrix  $A$  is positive-definite:**

$$\mathbf{x}^T A \mathbf{x} > 0 \quad \forall \mathbf{x}$$

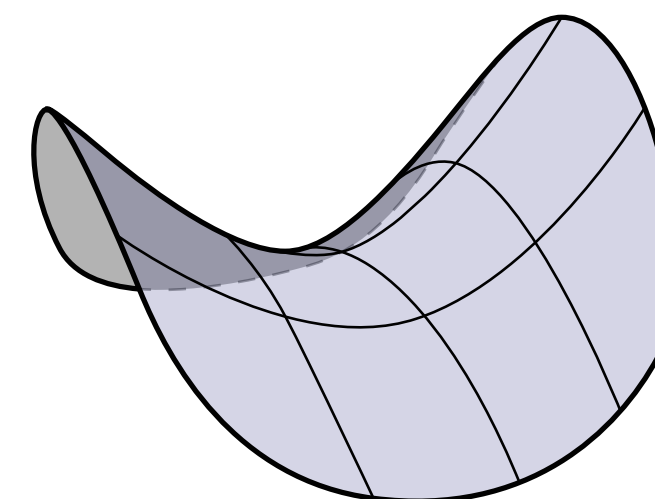
- **1D: Must have  $xax = ax^2 > 0$ . In other words:  $a$  is positive!**
- **2D: Graph of function looks like a “bowl”:**



positive definite



positive semidefinite



indefinite

Positive-definiteness **extremely important** in computer graphics: means we can find minimizers by solving linear equations. Starting point for many algorithms (geometry processing, simulation, ...)

# Minimizing Quadratic Error

- Find “best” point for edge collapse by minimizing quadratic form

$$\min_{\mathbf{u} \in \mathbb{R}^4} \mathbf{u}^T K \mathbf{u}$$

- Already know fourth (homogeneous) coordinate for a point is 1
- So, break up our quadratic function into two pieces:

$$\begin{aligned} & \begin{bmatrix} \mathbf{x}^T & 1 \end{bmatrix} \begin{bmatrix} B & \mathbf{w} \\ \mathbf{w}^T & d^2 \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ 1 \end{bmatrix} \\ & = \mathbf{x}^T B \mathbf{x} + 2\mathbf{w}^T \mathbf{x} + d^2 \end{aligned}$$

- Now we have a quadratic polynomial in the unknown position  $\mathbf{x} \in \mathbb{R}^3$
- Can minimize as before:

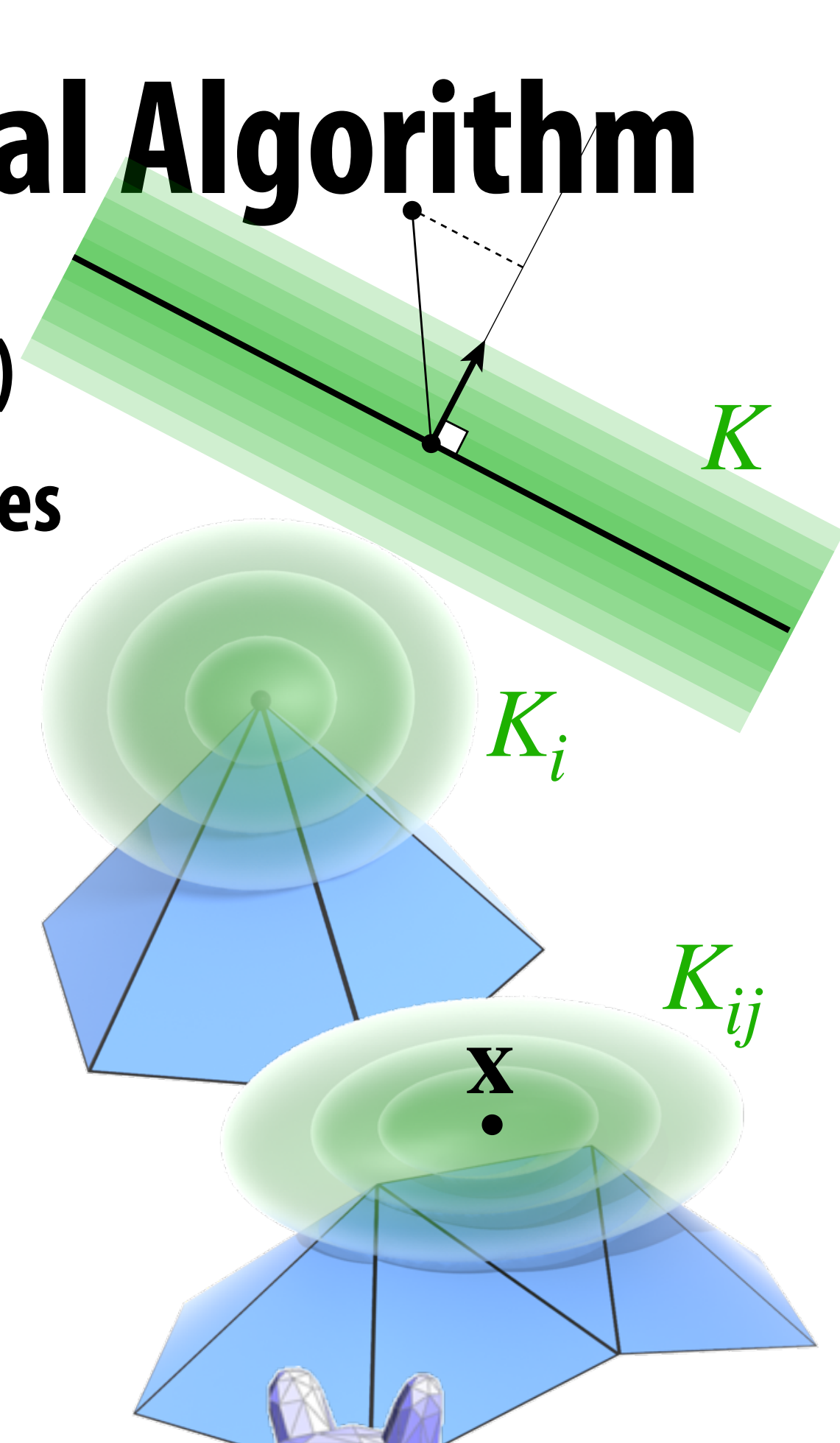
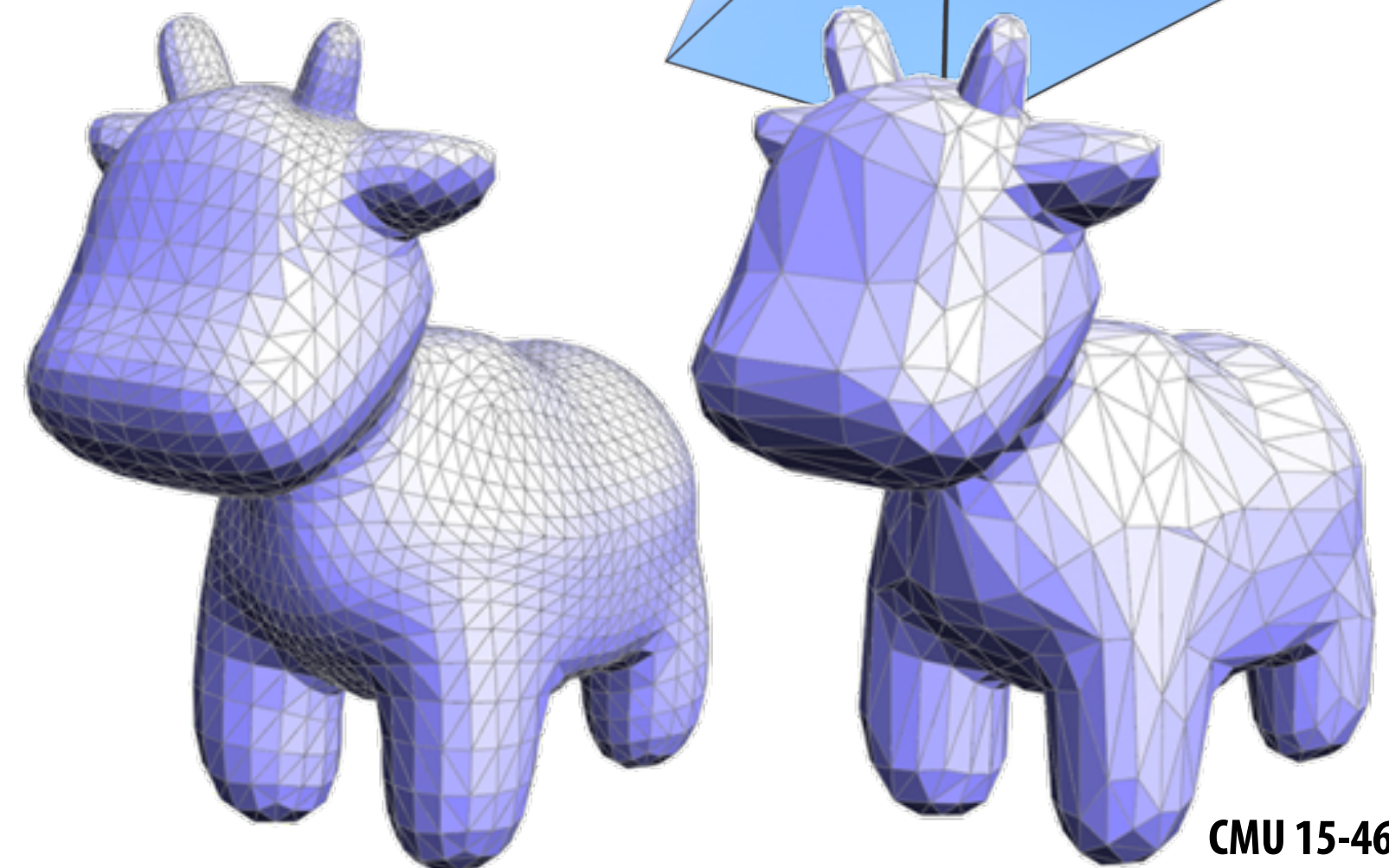
$$2B\mathbf{x} + 2\mathbf{w} = 0 \quad \iff \quad \mathbf{x} = -B^{-1}\mathbf{w}$$

**Q: Why should  $B$  be positive-definite?**



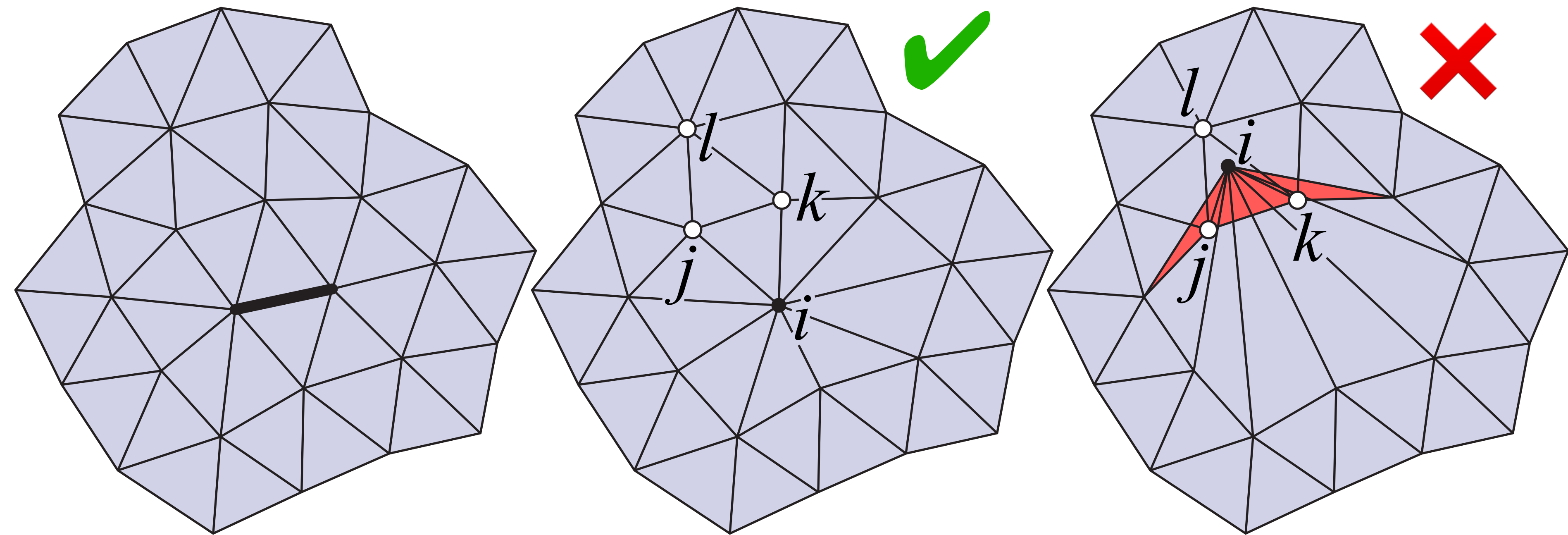
# Quadric Error Simplification: Final Algorithm

- Compute  $K$  for each triangle (squared distance to plane)
- Set  $K_i$  at each vertex to sum of  $K$ s from incident triangles
- For each edge  $e_{ij}$ :
  - set  $K_{ij} = K_i + K_j$
  - find point  $\mathbf{x}$  minimizing error, set cost to  $K_{ij}(\mathbf{x})$
- Until we reach target number of triangles:
  - collapse edge  $e_{ij}$  with smallest cost to optimal point  $\mathbf{x}$
  - set quadric at new vertex to  $K_{ij}$
  - update cost of edges touching new vertex
- **More details in assignment writeup!**



# Quadric Simplification—Flipped Triangles

- Depending on where we put the new vertex, one of the new triangles might be “flipped” (normal points in instead of out):



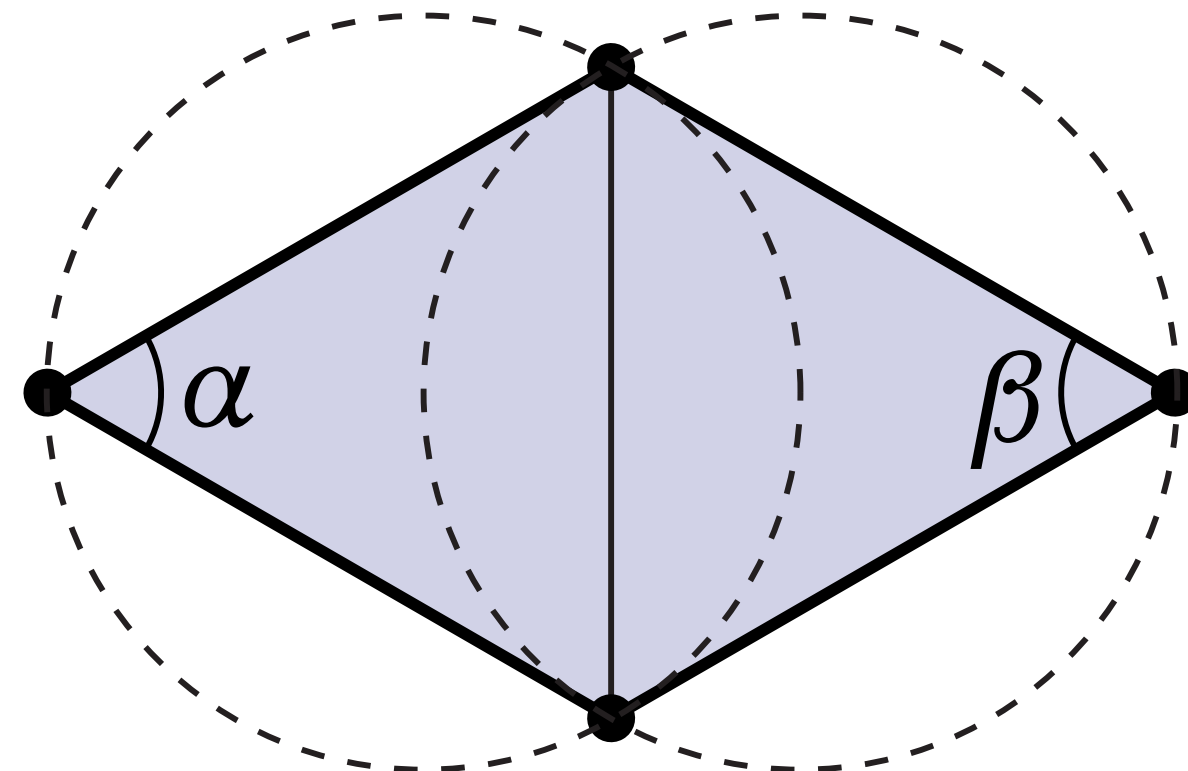
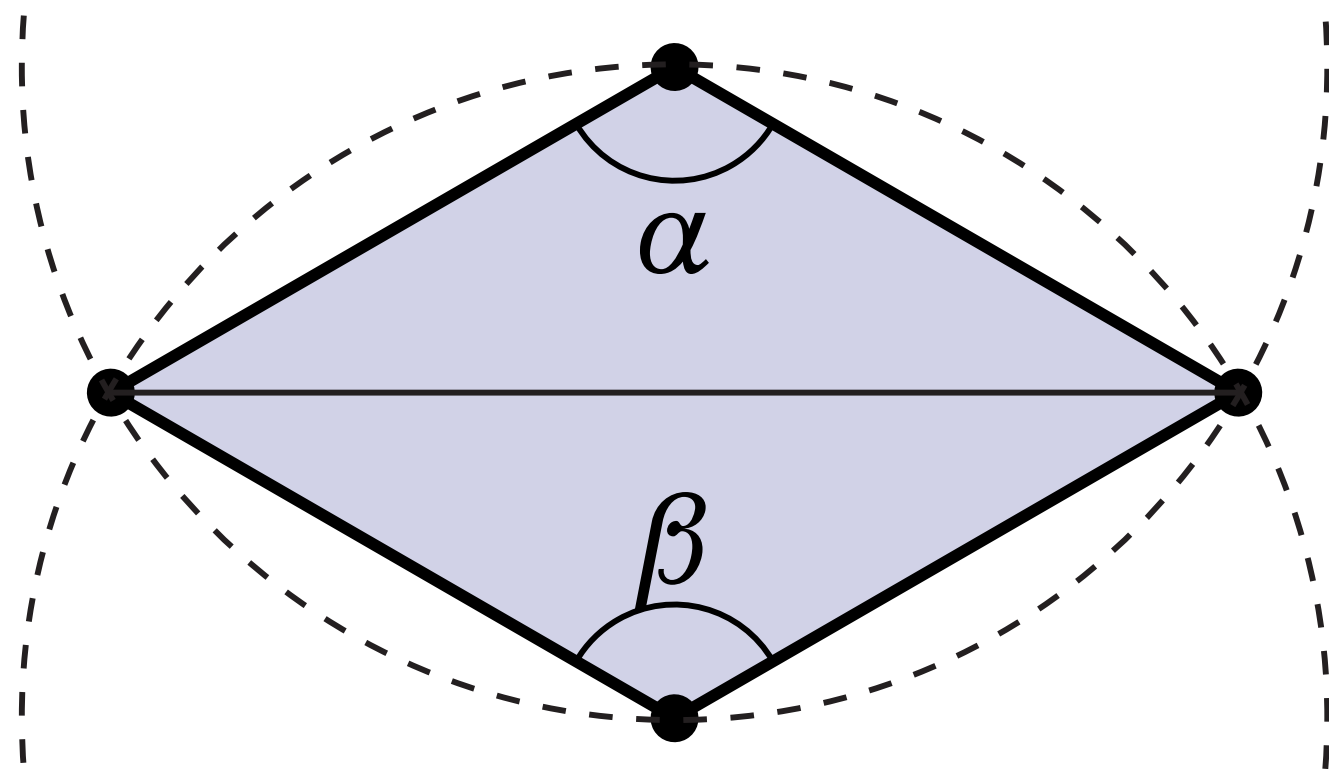
- **Easy solution:** for each triangle  $ijk$  touching collapsed vertex  $i$ , consider normals  $N_{ijk}$  and  $N_{kjl}$  (where  $kjl$  is other triangle containing edge  $jk$ )
- If  $\langle N_{ijk}, N_{kjl} \rangle$  is negative, don't collapse this edge!

**What if we're happy with the number of triangles, but want to improve quality?**



# How do we make a mesh “more Delaunay”?

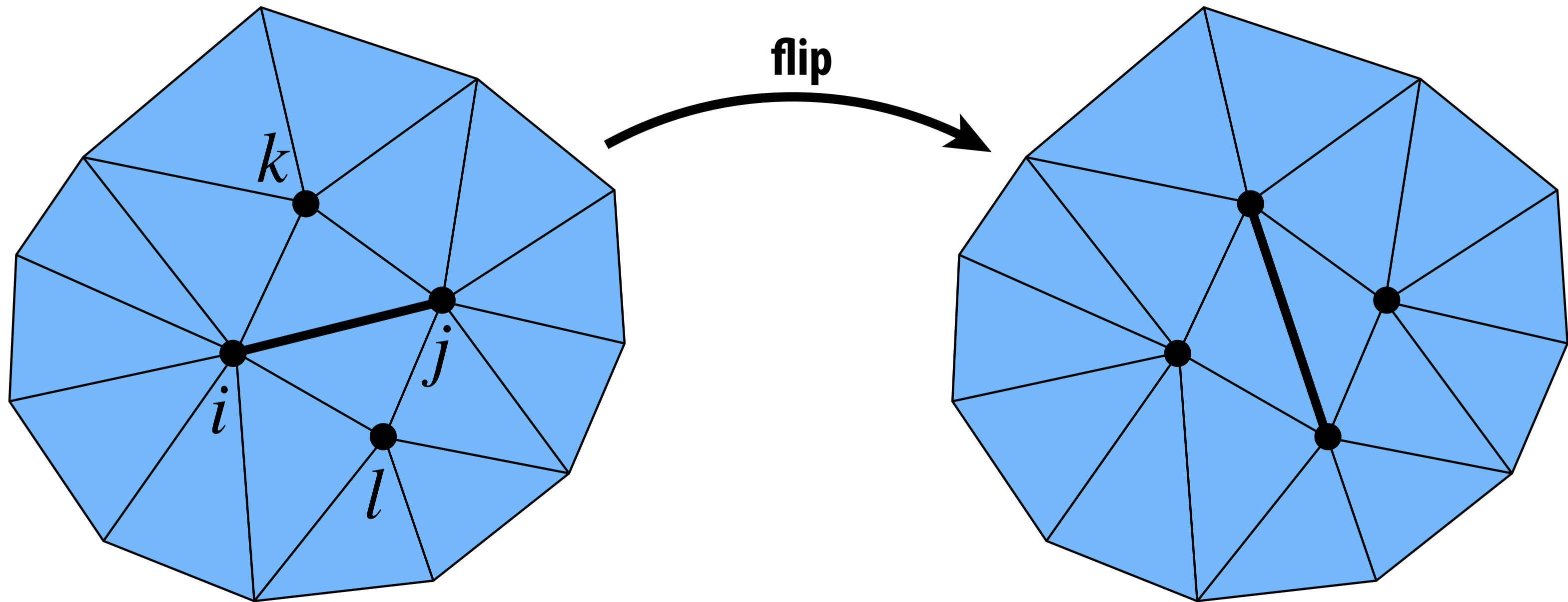
- Already have a good tool: edge flips!
- If  $\alpha + \beta > \pi$ , flip it!



- **FACT:** in 2D, flipping edges eventually yields Delaunay mesh
- **Theory:** worst case  $O(n^2)$ ; doesn't always work for surfaces in 3D
- **Practice:** simple, effective way to improve mesh quality

# Alternatively: how do we improve degree?

- Same tool: edge flips!
- If total deviation from degree-6 gets smaller, flip it!

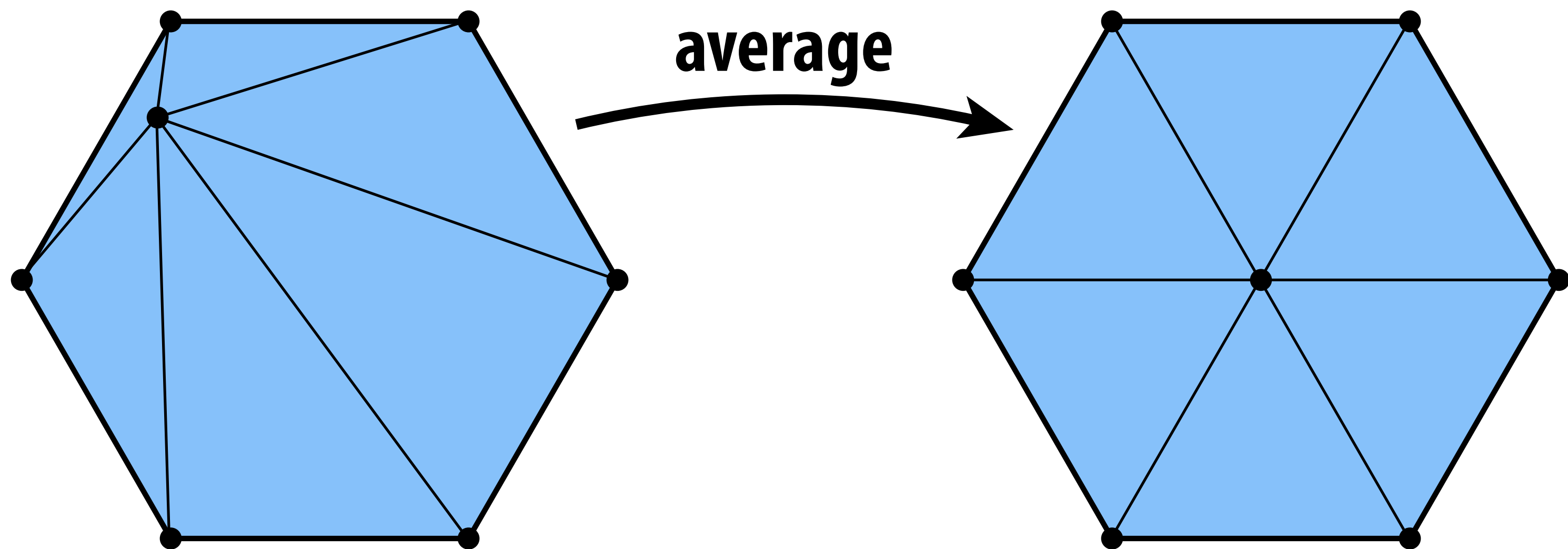


total deviation:  $|d_i - 6| + |d_j - 6| + |d_k - 6| + |d_l - 6|$

- **FACT: average degree approaches 6 as number of elements increases**
- **Iterative edge flipping acts like “discrete diffusion” of degree**
- **No (known) guarantees; works well in practice**

# How do we make a triangles “more round”?

- Delaunay doesn't guarantee triangles are “round” (angles near  $60^\circ$ )
- Can often improve shape by centering vertices:

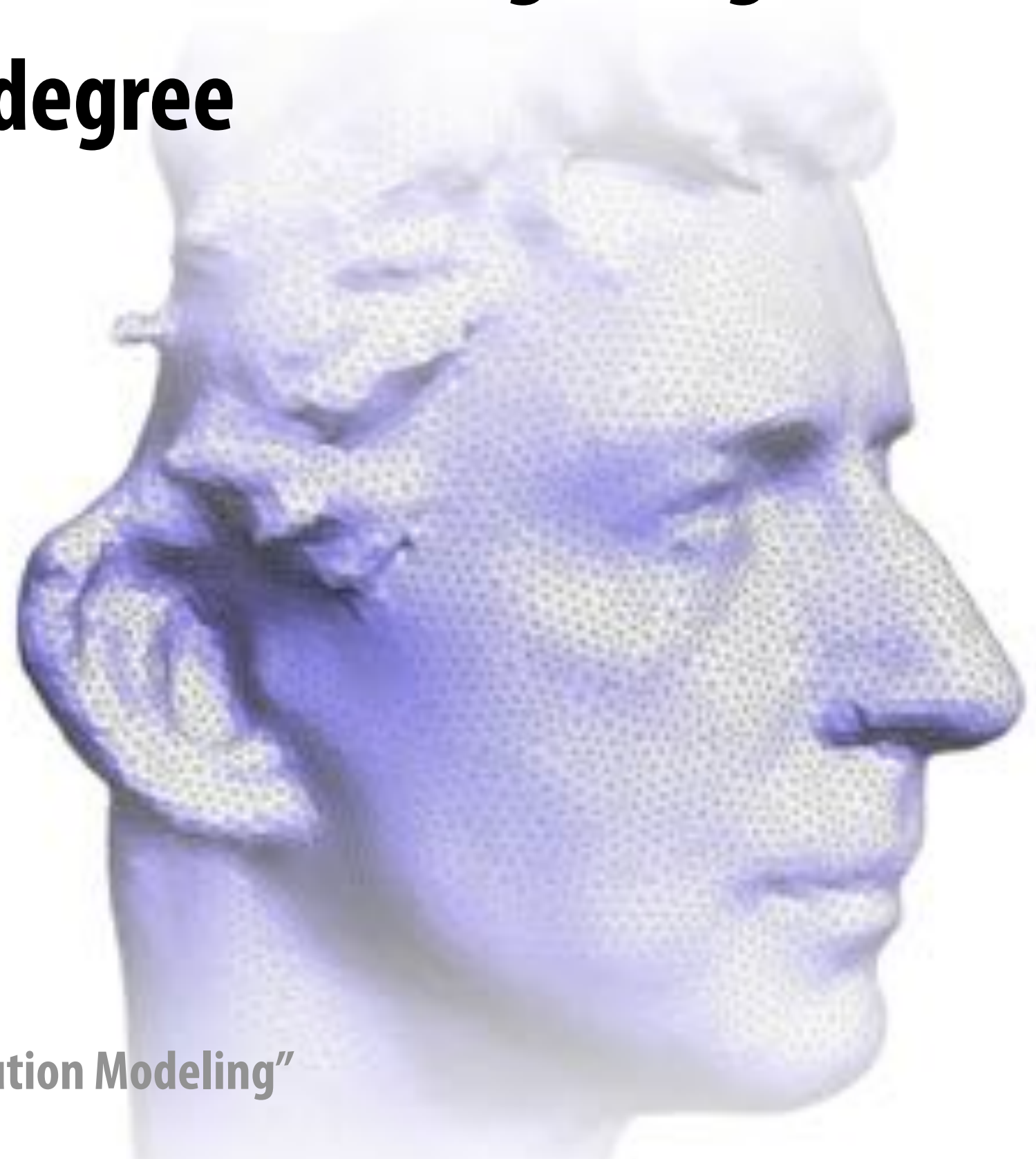
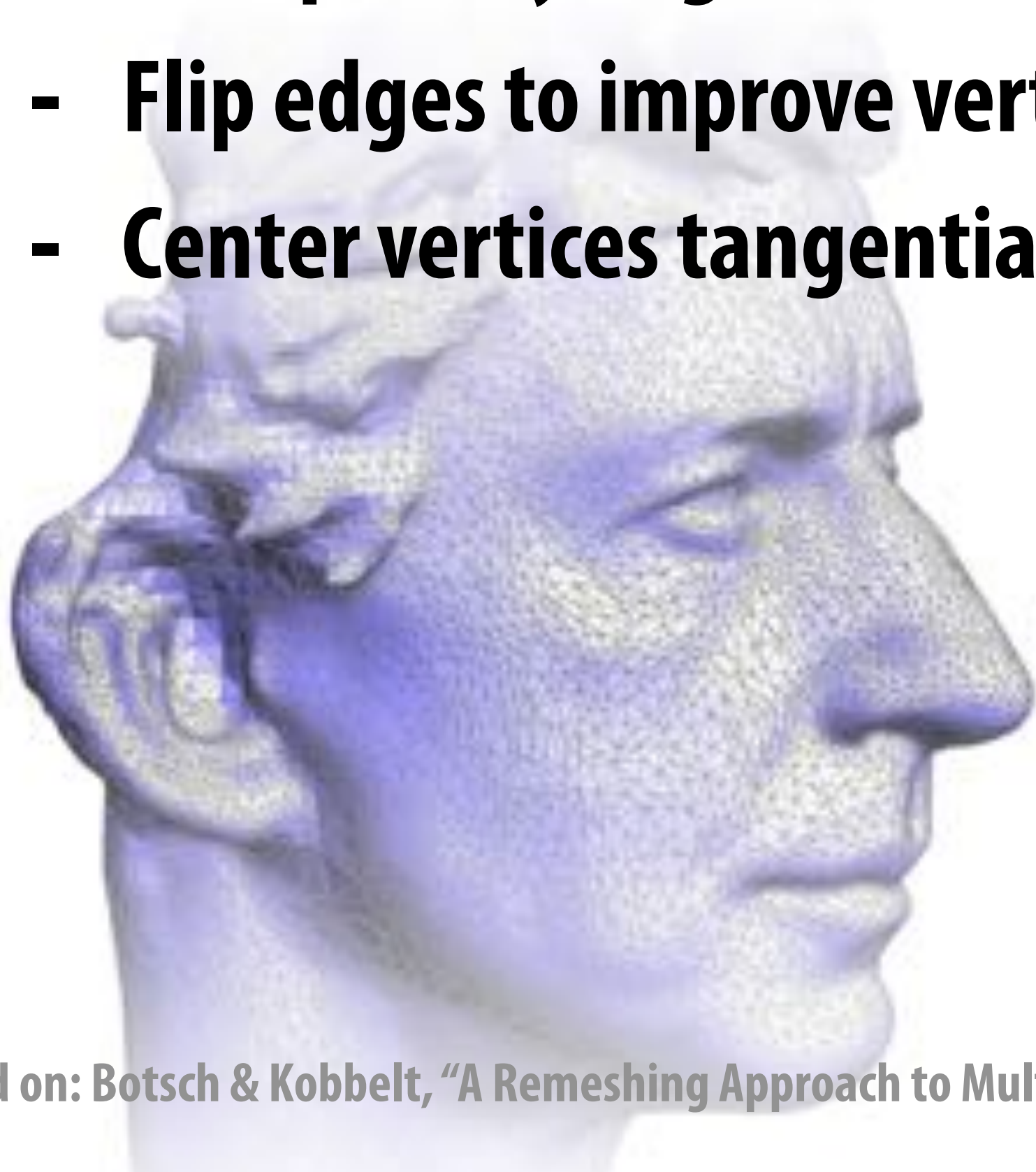


- Simple version of technique called “Laplacian smoothing”
- On surface: move only in tangent direction
- How? Remove normal component from update vector



# Isotropic Remeshing Algorithm

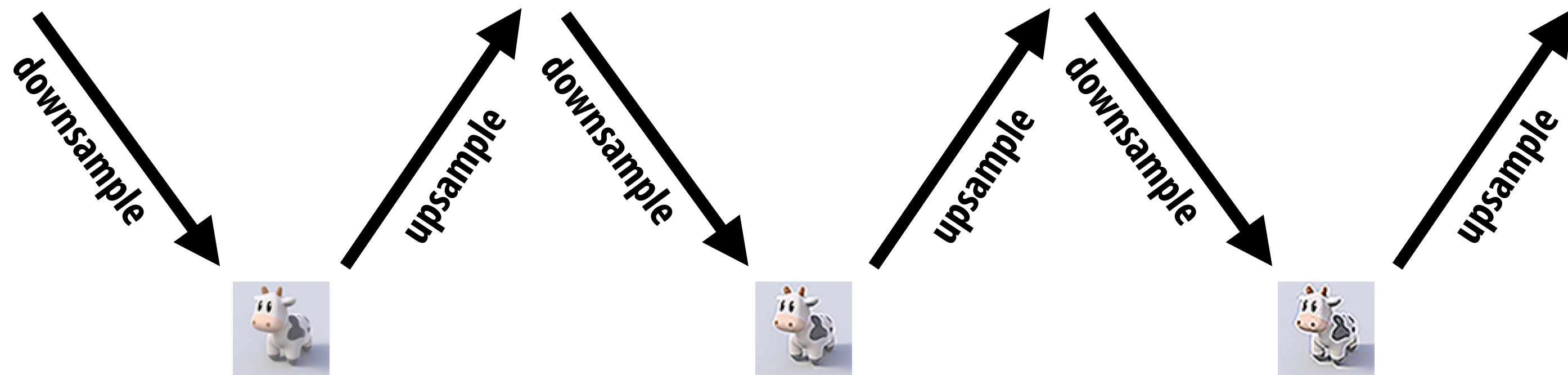
- Try to make triangles uniform shape & size
- Repeat four steps:
  - Split any edge over  $\frac{4}{3}$  mean edge length
  - Collapse any edge less than  $\frac{4}{5}$  mean edge length
  - Flip edges to improve vertex degree
  - Center vertices tangentially



**What can go wrong when  
you resample a signal?**

# Danger of Resampling

**Q: What happens if we repeatedly resample an image?**

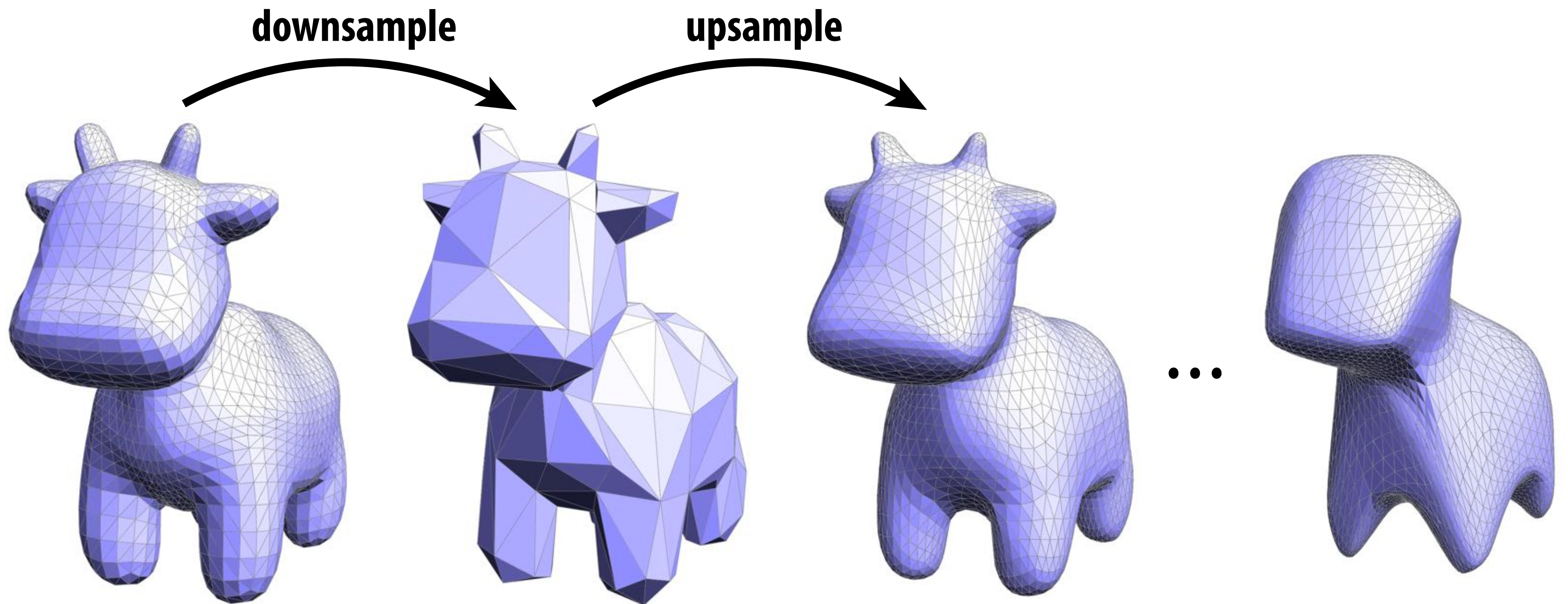


**A: Signal quality degrades!**



# Danger of Resampling

**Q: What happens if we repeatedly resample a mesh?**



**A: Signal also degrades!**

**But wait: we have the original signal (mesh).**  
**Why not just project each new sample point**  
**onto the closest point of the original mesh?**



# Next Time: Geometric Queries

- **Q: Given a point, in space, how do we find the closest point on a surface? Are we inside or outside the surface? How do we find intersection of two triangles? Etc.**
- **Do implicit/explicit representations make such tasks easier?**
- **What's the cost of the naïve algorithm, and how do we accelerate such queries for large meshes?**
- **So many questions!**

