# Meshes and Manifolds 

## Computer Graphics <br> CMU 15-462/15-662

## Last time: overview of geometry

- Many types of geometry in nature
- Demand sophisticated representations
- Two major categories:
- IMPLICIT - "tests" if a point is in shape
- EXPLICIT - directly"lists" points
- Lots of representations for both
- Today:
- what is a surface, anyway?
- nuts \& bolts of polygon meshes
- geometry processing/resampling

Geometry


## Manifold Assumption

- Today we're going to introduce the idea of manifold geometry
- Can be hard to understand motivation at first!
- So first, let's revisit a more familiar example...



## Bitmap Images, Revisited

To encode images, we used a regular grid of pixels:


# But images are not fundamentally made of little squares: 



Goyō Hashiguchi, Kamisuki (ca 1920)

## So why did we choose a square grid?


...rather than dozens of alternatives?

## Regular grids make life easy

- One reason: SIMPLICITY / EFFICIENCY
- E.g., always have four neighbors
- Easy to index, easy to filter...
- Storage is just a list of numbers
- Another reason: GENERALITY
- Can encode basically any image

|  | $(i, j-1)$ |  |
| :--- | :--- | :--- |
| $(i-1, j)$ | $(i, j)$ | $(i+1, j)$ |
|  | $(i, j+1)$ |  |

- Are regular grids always the best choice for bitmap images?
- No! E.g., suffer from anisotropy, don't capture edges, ...
- But more often than not are a pretty good choice
- Will see a similar story with geometry...


## So, how should we encode surfaces?

## Smooth Surfaces

■ Intuitively, a surface is the boundary or "shell" of an object

- (Think about the candy shell, not the chocolate.)
- Surfaces are manifold:
- If you zoom in far enough (at any point) looks like a plane*
- E.g., the Earth from space vs. from the ground



## Isn't every shape manifold?

- No, for instance:


Center point never looks like the plane, no matter how close we get.

## More Examples of Smooth Surfaces

- Which of these shapes are manifold?



## A manifold polygon mesh has fans, not fins

- For polygonal surfaces just two easy conditions to check:

1. Every edge is contained in only two polygons (no "fins")
2. The polygons containing each vertex make a single "fan"


## What about boundary?

- The boundary is where the surface "ends."
- E.g., waist \& ankles on a pair of pants.
- Locally, looks like a half disk
- Globally, each boundary forms a loop

- Polygon mesh:

- one polygon per boundary edge
- boundary vertex looks like "pacman"


## Ok, but why is the manifold assumption useful?

## Keep it Simple!

- Same motivation as for images:
- make some assumptions about our geometry to keep data structures/algorithms simple and efficient
- in many common cases, doesn't fundamentally limit what we can do with geometry

|  | $(i, j-1)$ |  |
| :--- | :--- | :--- |
| $(i-1, j)$ | $(i, j)$ | $(i+1, j)$ |
|  | $(i, j+1)$ |  |



## How do we actually encode all this data?

## Warm up: storing numbers

- Q: What data structures can we use to store a list of numbers?

■ One idea: use an array (constant time lookup, coherent access)

| 1.7 | 2.9 | 0.3 | 7.5 | 9.2 | 4.8 | 6.0 | 0.1 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

■ Alternative: use a linked list (linear lookup, incoherent access)


- Q: Why bother with the linked list?
- A: For one, we can easily insert numbers wherever we like...


## Polygon Soup (Array-like)

- Store triples of coordinates ( $x, y, z$ ), tuples of indices
- E.g., tetrahedron:

|  | VERTICES |  |  |
| ---: | ---: | ---: | ---: |
|  | $\mathbf{x}$ | $\mathbf{y}$ | $\mathbf{z}$ |
| $\mathbf{0}:$ | -1 | -1 | -1 |
| $1:$ | 1 | -1 | 1 |
| $\mathbf{2}:$ | 1 | 1 | -1 |
| $3:$ | -1 | 1 | 1 |

- Q: How do we find all the polygons touching vertex 2?
- Ok, now consider a more complicated mesh:


POLYGONS

| $\mathbf{i}$ | $\mathbf{j}$ | $\mathbf{k}$ |
| :--- | :--- | :--- |
| 0 | 2 | 1 |
| 0 | 3 | 2 |
| 3 | 0 | 1 |
| 3 | 1 | 2 |



Very expensive to find the neighboring triangles! (What's the cost?)

## Incidence Matrices

- If we want to answer neighborhood queries, why not simply store a list of neighbors?
- Can encode all neighbor information via incidence matrices
- E.g., tetrahedron: VERTEX $\leftrightarrow$ EDGE

| EDGE $\leftrightarrow$ FACE |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | e1 | e2 | e3 | e4 | e5 |
| f0 | 1 | 0 | 0 | 1 | 0 | 1 |
| f1 | 0 | 1 | 0 | 0 | 1 | 1 |
| f2 | 1 | 1 | 1 | 0 | 0 | 0 |
| f3 | 0 | 0 | 1 | 1 | 1 | 0 |

- 1 means "touches"; 0 means "does not touch"
- Instead of storing lots of 0's, use sparse matrices
- Still large storage cost, but finding neighbors is now 0 (1)
- Hard to change connectivity, since we used fixed indices

■ Bonus feature: mesh does not have to be manifold

## Halfedge Data Structure (Linked-list-like)

- Store some information about neighbors
- Don't need an exhaustive list; just a few key pointers

■ Key idea: two halfedges act as "glue" between mesh elements:


Each vertex, edge face points to just one of its halfedges.

## Halfedge makes mesh traversal easy

- Use"twin" and "next" pointers to move around mesh

■ Use "vertex", "edge", and "face" pointers to grab element
■ Example: visit all vertices of a face:


- Example: visit all neighbors of a vertex:

```
Halfedge* h = v->halfedge;
do {
        h = h->twin->next;
}
while( h != v->halfedge );
```

■ Note: only makes sense if mesh is manifold!


## Halfedge meshes are always manifold

- Consider simplified halfedge data structure
- Require only "common-sense" conditions

```
struct Halfedge {
    Halfedge *next, *twin;
};
```

```
twin->twin == this
next != this
twin != this
```

- Keep following next, and you'll get faces.
- Keep following twin and you'll get edges.
- Keep following next->twin and you'll get vertices.


Q: Why, therefore, is it impossible to encode the red figures?

## Halfedge meshes are easy to edit

- Remember key feature of linked list: insert/delete elements
- Same story with halfedge mesh ("linked list on steroids")
- E.g., for triangle meshes, several atomic operations:

- How? Allocate/delete elements; reassigning pointers.
- Must be careful to preserve manifoldness!


## Edge Flip (Triangles)

- Triangles ( $\mathbf{a}, \mathrm{b}, \mathrm{c}$ ), ( $\mathbf{b}, \mathrm{d}, \mathrm{c}$ ) become ( $\mathbf{a}, \mathrm{d}, \mathrm{c}$ ), ( $\mathbf{a}, \mathrm{b}, \mathrm{d}$ ):


■ Long list of pointer reassignments (edge->halfedge = ...)

- However, no elements created/destroyed.
- Q: What happens if we flip twice?
- Challenge: can you implement edge flip such that pointers are unchanged after two flips?


## Edge Split (Triangles)

■ Insert midpoint $m$ of edge ( $\mathbf{c}, \mathrm{b}$ ), connect to get four triangles:


- This time, have to add new elements.
- Lots of pointer reassignments.

■ Q: Can we "reverse" this operation?

## Edge Collapse (Triangles)

- Replace edge ( $b, \mathrm{c}$ ) with a single vertex m :

- Now have to delete elements.
- Still lots of pointer assignments!

■ Q: How would we implement this with a polygon soup?

- Any other good way to do it? (E.g., different data structure?)


## Comparison of Polygon Mesh Data Strucutres

| Case study: <br> triangles. | Polygon Soup | Incidence <br> Matrices | Halfedge Mesh |
| :---: | :---: | :---: | :---: |
| storage cost* | $\sim 3$ x \#vertices | $\sim 33$ x \#vertices | $\sim 36$ x \#vertices |
| constant-time <br> neighborhood access? | NO | YES | YES |
| easy to add/remove <br> mesh elements? | NO | NO | YES |
| nonmanifold <br> geometry? | YES | YES | NO |

## Conclusion: pick the right data structure for the job!

*number of integer values and/or pointers required to encode connectivity (all data structures require same amount of storage for vertex positions)

## Alternatives to Halfedge

- Many very similar data structures:
- winged edge
- corner table
- quadedge
- ....


cubeeroct

Paul Heckbert (former CMU prof.) quadedge code - http://bit.ly/1QZLHos

dodec $\leftrightarrow$ ices

- Each stores local neighborhood information
- Similar tradeoffs relative to simple polygon list:
- CONS: additional storage, incoherent memory access
- PROS: better access time for individual elements, intuitive traversal of local neighborhoods
- (Food for thought: can you design a halfedge-like data structure with reasonably coherent data storage?)


## Ok, but what can we actually do with our fancy new data structure?

## Subdivision Modeling



## Subdivision Modeling

- Common modeling paradigm in modern 3D tools:
- Coarse "control cage"
- Perform local operations to control/edit shape
- Global subdivision process determines final surface



## Subdivision Modeling—Local Operations

- For general polygon meshes, we can dream up lots of local mesh operations that might be useful for modeling:

...and many, many more!


## Global Subdivision

- Start with coarse polygon mesh ("control cage")
- Subdivide each element
- Update vertices via local averaging
- Many possible rule:
- Catmull-Clark (quads)
- Loop (triangles)
-     - • -
- Common issues:

- interpolating or approximating?
- continuity at vertices?

- Easier than splines for modeling; harder to evaluate pointwise


## Next Time: Digital Geometry Processing

- Extend traditional digital signal processing (audio, video, etc.) to deal with geometric signals:
- upsampling/downsampling / resampling / filtering ...
- aliasing (reconstructed surface gives "false impression")
- Also some new challenges (very recent field!):
- over which domain is a geometric signal expressed?
- no terrific sampling theory, no fast Fourier transform, ...
- Often need new data structures \& new algorithms


