# Meshes and Manifolds

**Computer Graphics CMU 15-462/15-662** 

# Last time: overview of geometry

- Many types of geometry in nature
- **Demand sophisticated representations**
- **Two major categories:** 
  - IMPLICIT "tests" if a point is in shape
  - EXPLICIT directly "lists" points
- Lots of representations for both

### Today:

- what is a surface, anyway?
- nuts & bolts of polygon meshes
- geometry processing / resampling

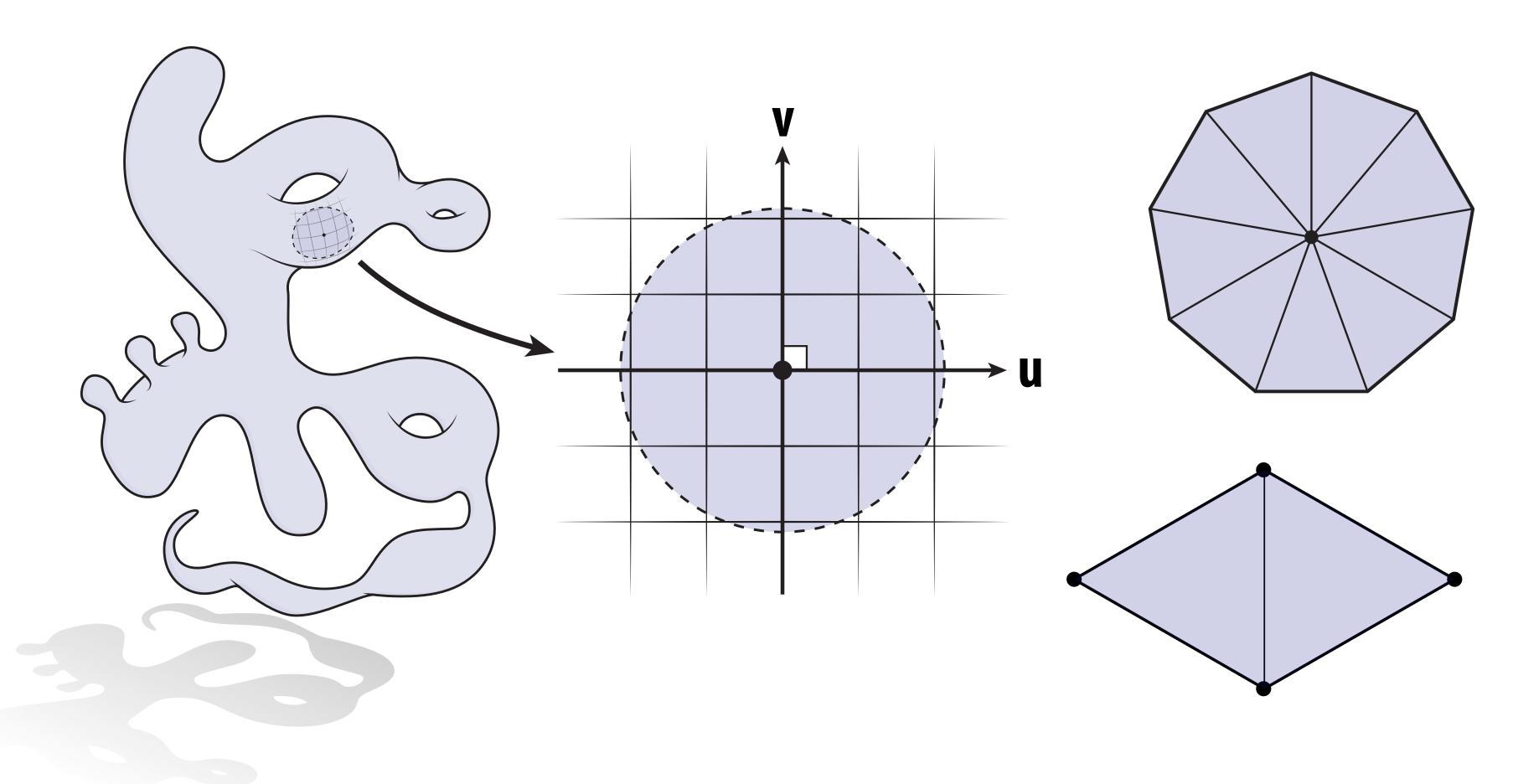
### Geometry





## Manifold Assumption

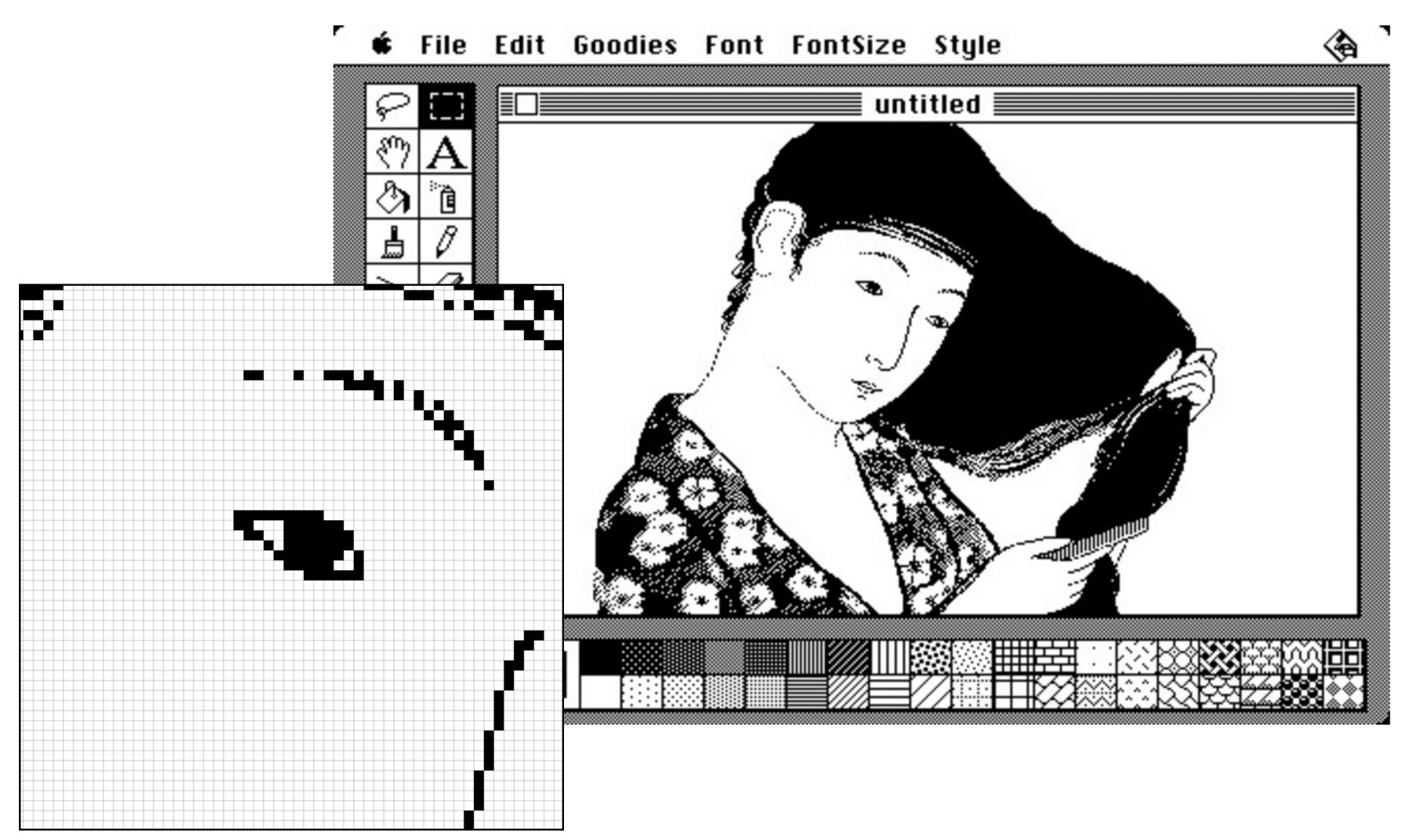
- Can be hard to understand motivation at first!
- So first, let's revisit a more familiar example...



# Today we're going to introduce the idea of *manifold* geometry



## Bitmap Images, Revisited To encode images, we used a *regular grid* of pixels:

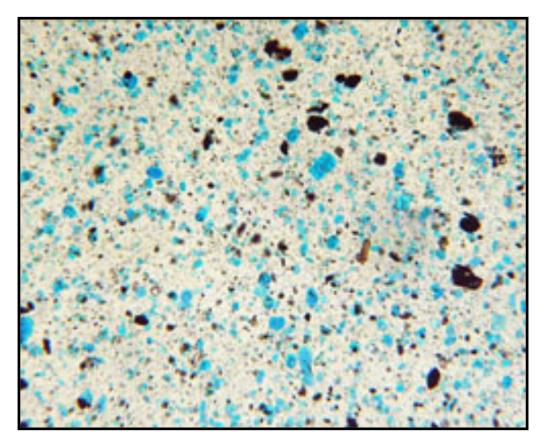




# But images are not fundamentally made of little squares:



Goyō Hashiguchi, *Kamisuki* (ca 1920)



photomicrograph of paint

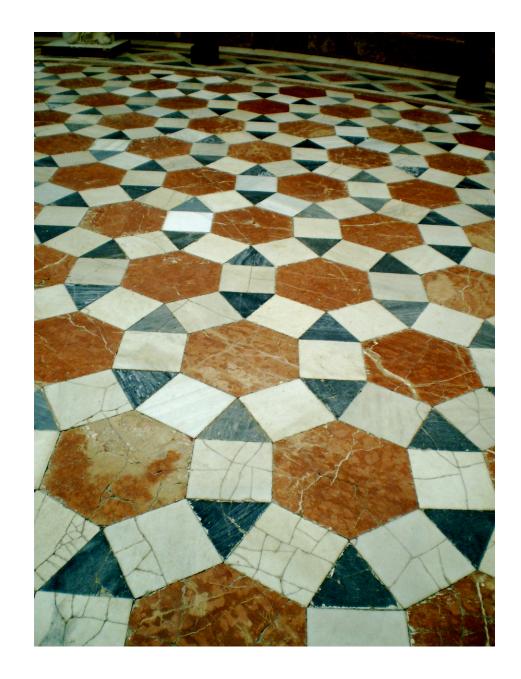


# So why did we choose a square grid?





## ... rather than dozens of possible alternatives?





# **Regular grids make life easy**

- **One reason: SIMPLICITY / EFFICIENCY** 
  - E.g., always have four neighbors
  - Easy to index, easy to filter...
  - Storage is just a list of numbers
- **Another reason: GENERALITY** 
  - Can encode basically any image
- Are regular grids always the best choice for bitmap images?
  - No! E.g., suffer from anisotropy, don't capture edges, ...
  - But more often than not are a pretty good choice
- Will see a similar story with geometry...



(i,j-1)

(i,j)

(i,j+1)

(i+1,j)

(i-1,j)

## So, how should we encode surfaces?



## **Smooth Surfaces**

- (Think about the candy shell, not the chocolate.)
- Surfaces are *manifold*:

  - E.g., the Earth from space vs. from the ground

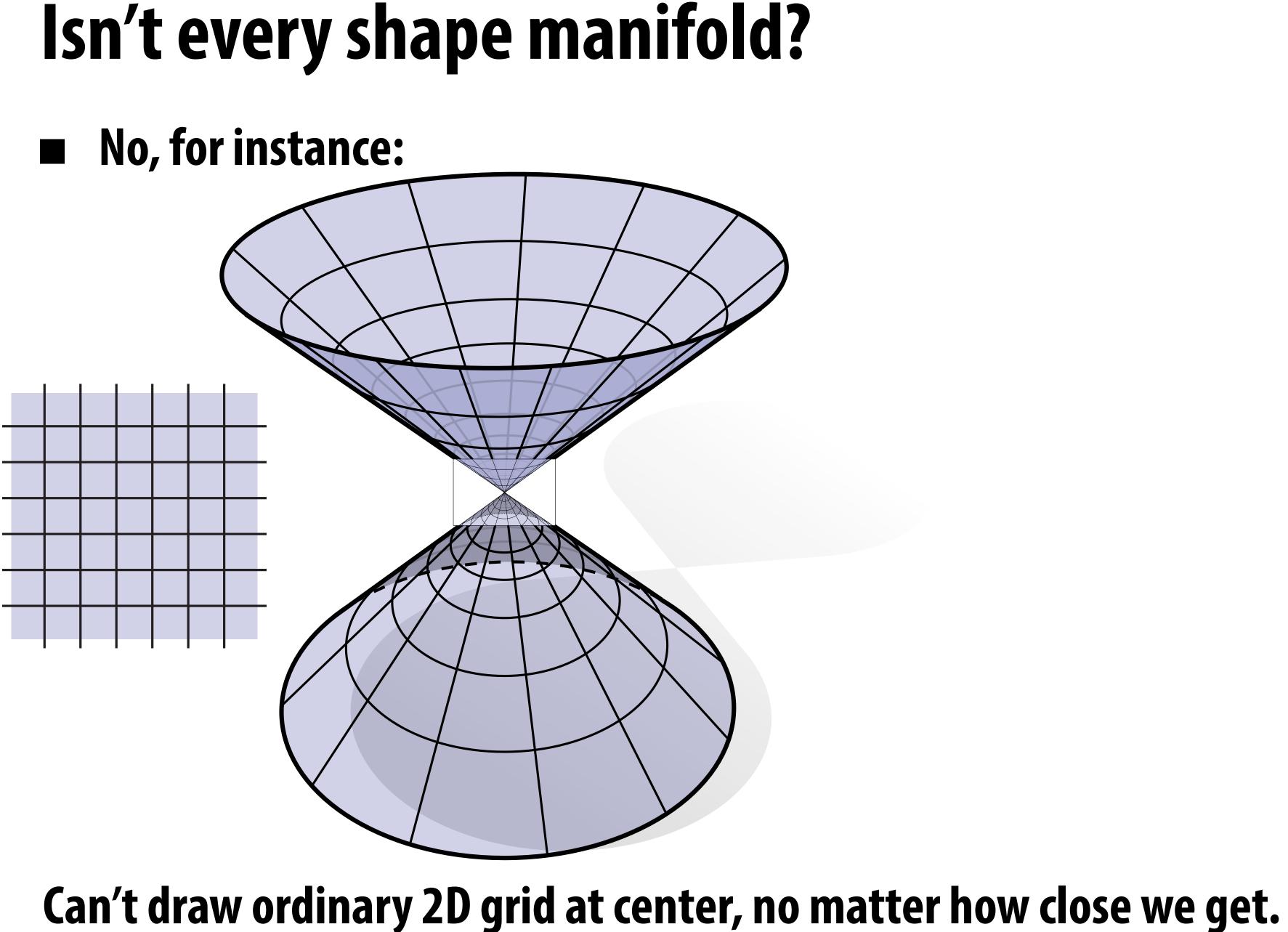


# Intuitively, a *surface* is the boundary or "shell" of an object

# - If you zoom in far enough, can draw a regular coordinate grid

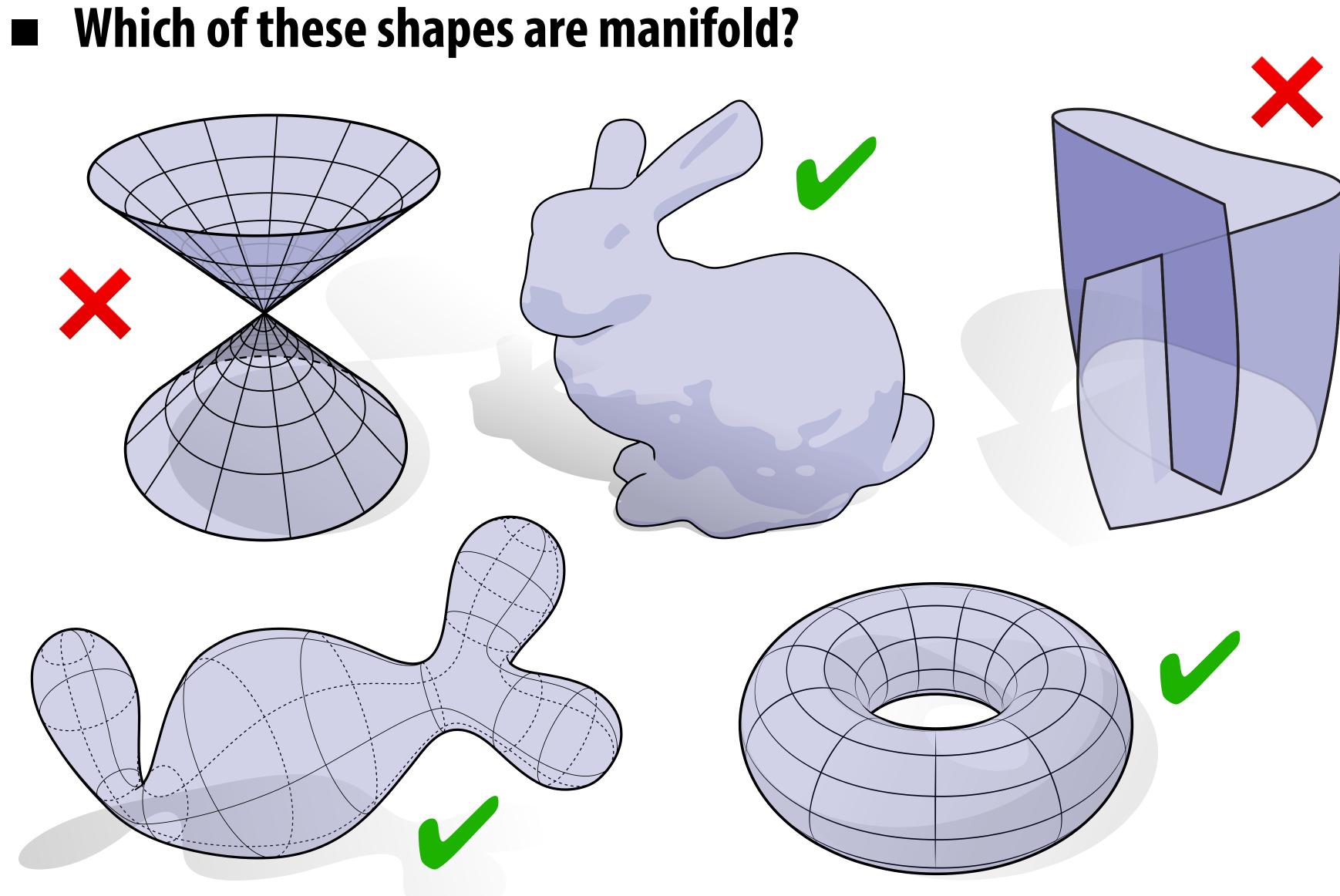








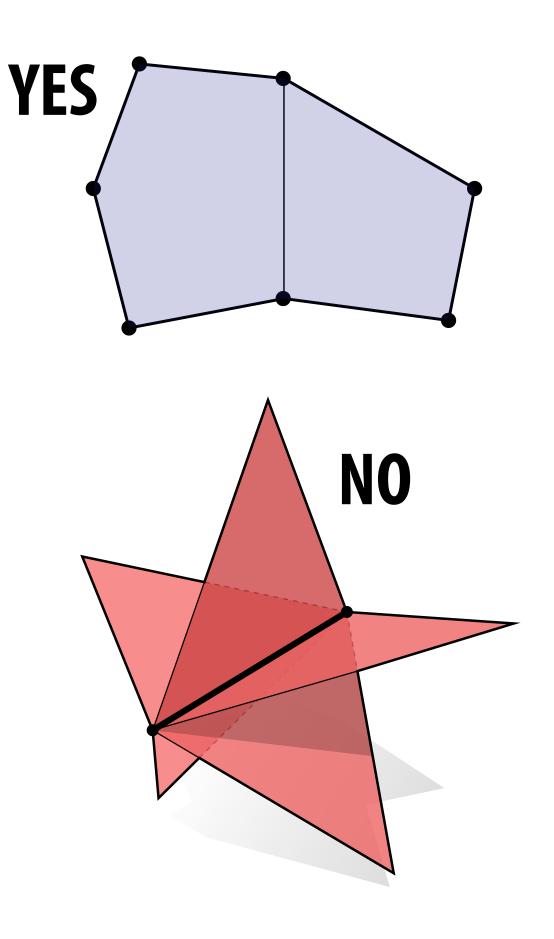
# Examples—Manifold vs. Nonmanifold

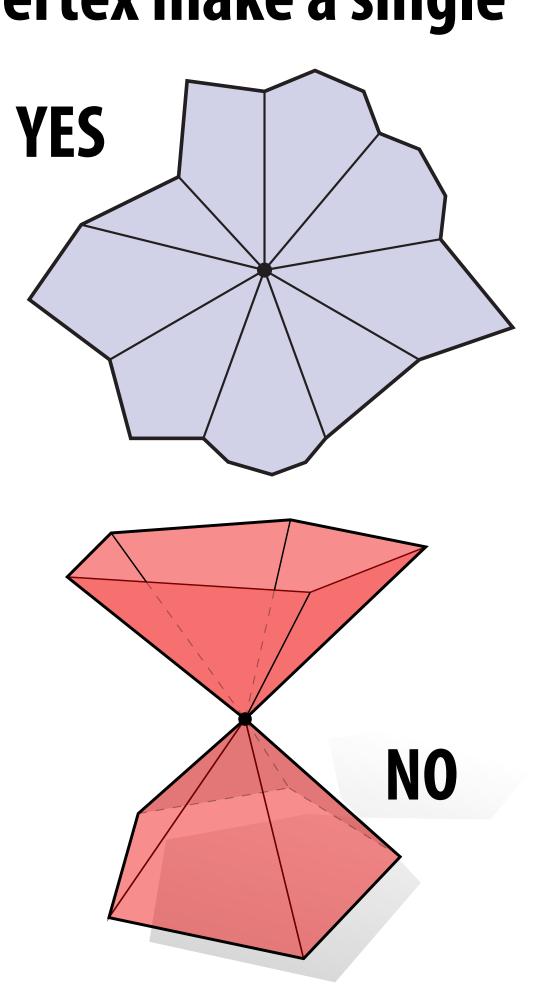




### A manifold polygon mesh has fans, not fins For polygonal surfaces just two easy conditions to check: 1. Every edge is contained in only two polygons (no "fins")

2. The polygons containing each vertex make a single "fan"

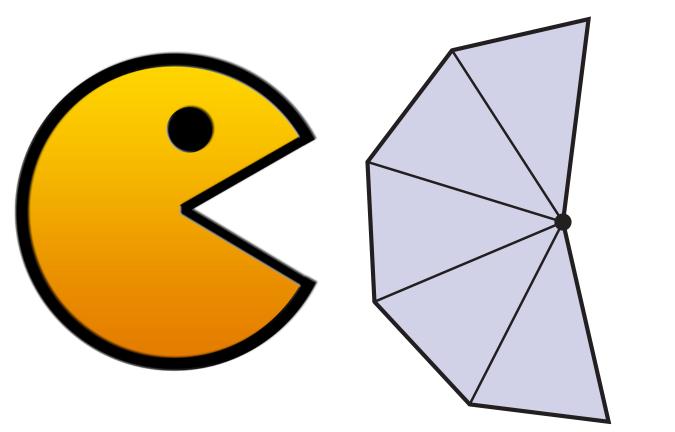






## What about boundary?

- The boundary is where the surface "ends."
- E.g., waist & ankles on a pair of pants.
- Locally, looks like a *half* disk
- Globally, each boundary forms a loop



### Polygon mesh:

- one polygon per boundary edge
- boundary vertex looks like "pacman"

YES





# Ok, but why is the manifold assumption useful?



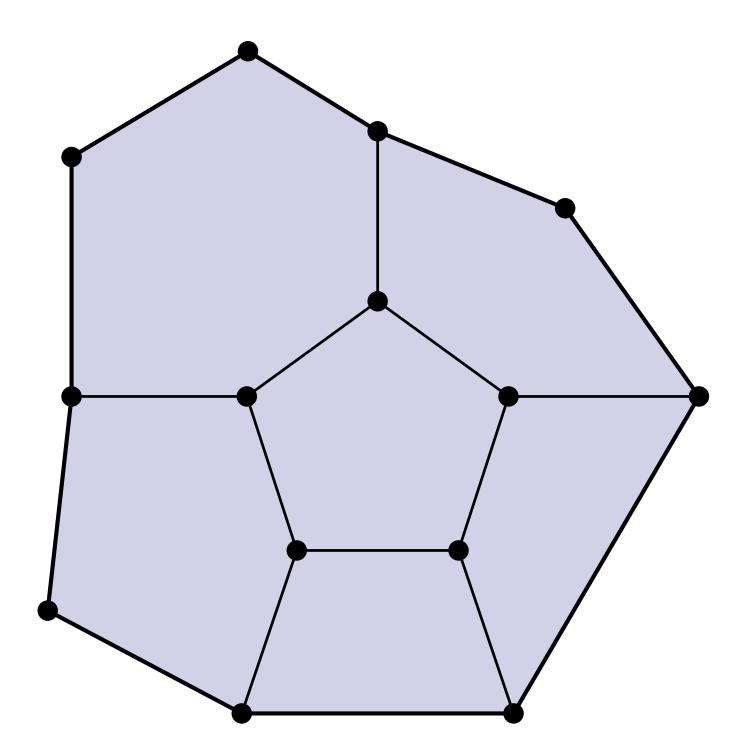
# **Keep it Simple!**

### Same motivation as for images:

	(i,j-1)	
(i-1,j)	(i,j)	(i+1,j)
	(i,j+1)	

# make some assumptions about our geometry to keep data structures/algorithms simple and efficient

# in *many common cases,* doesn't fundamentally limit what we can do with geometry



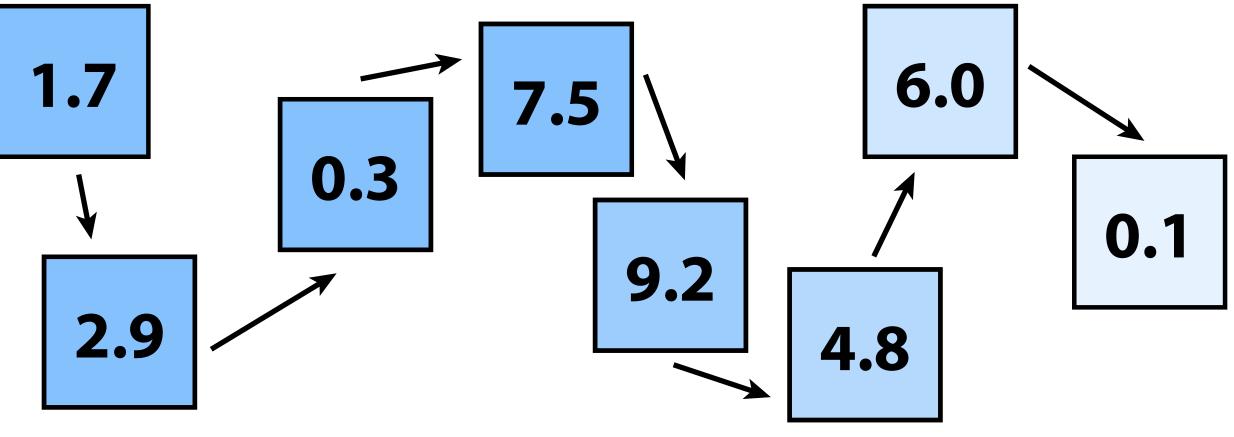


## How do we actually encode all this data?



## Warm up: storing numbers

### 



Q: Why bother with the linked list?

Q: What data structures can we use to store a list of numbers? One idea: use an array (constant time lookup, coherent access)

Alternative: use a linked list (linear lookup, incoherent access)

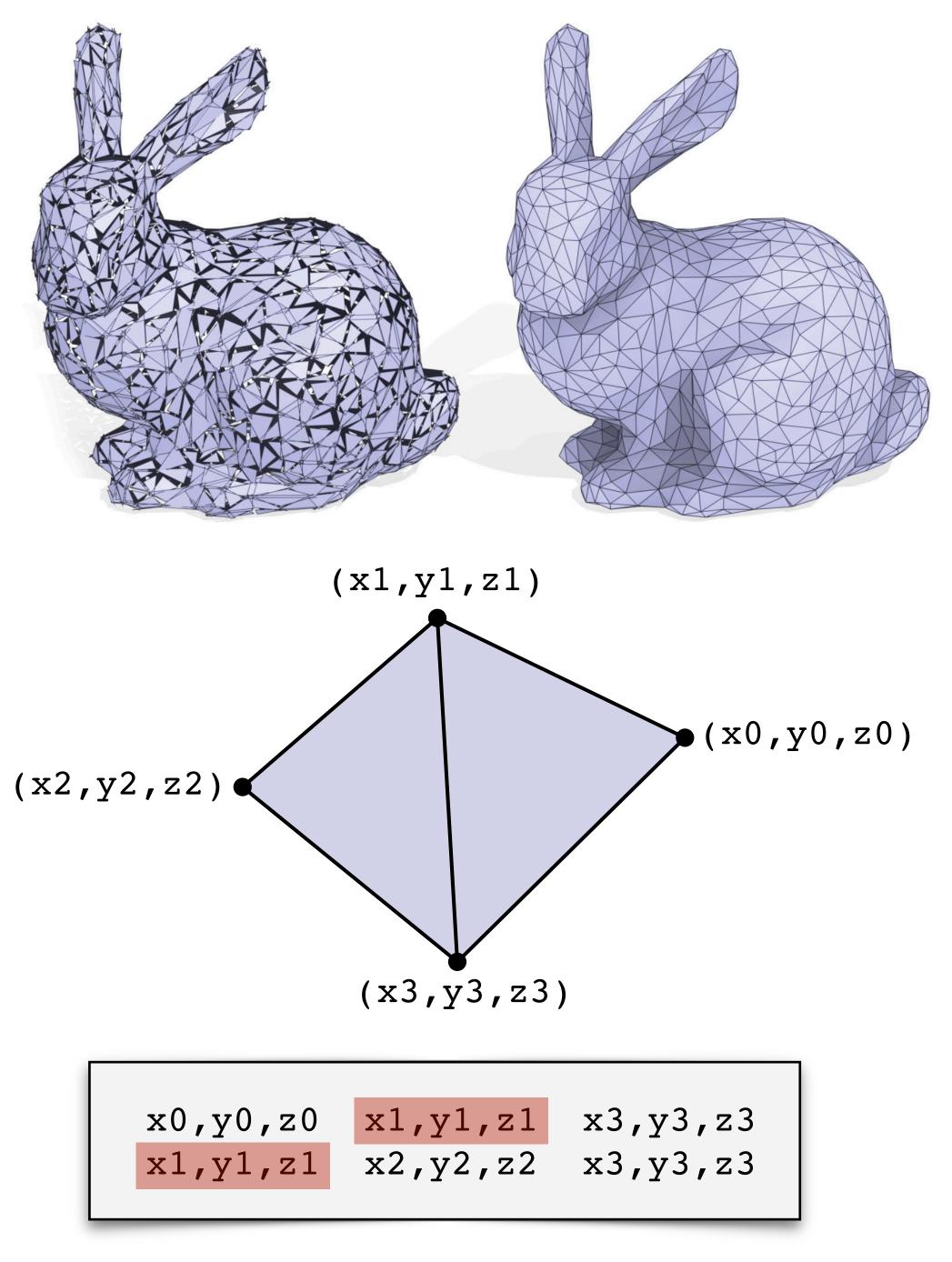
- A: For one, we can easily insert numbers wherever we like...



# Polygon Soup

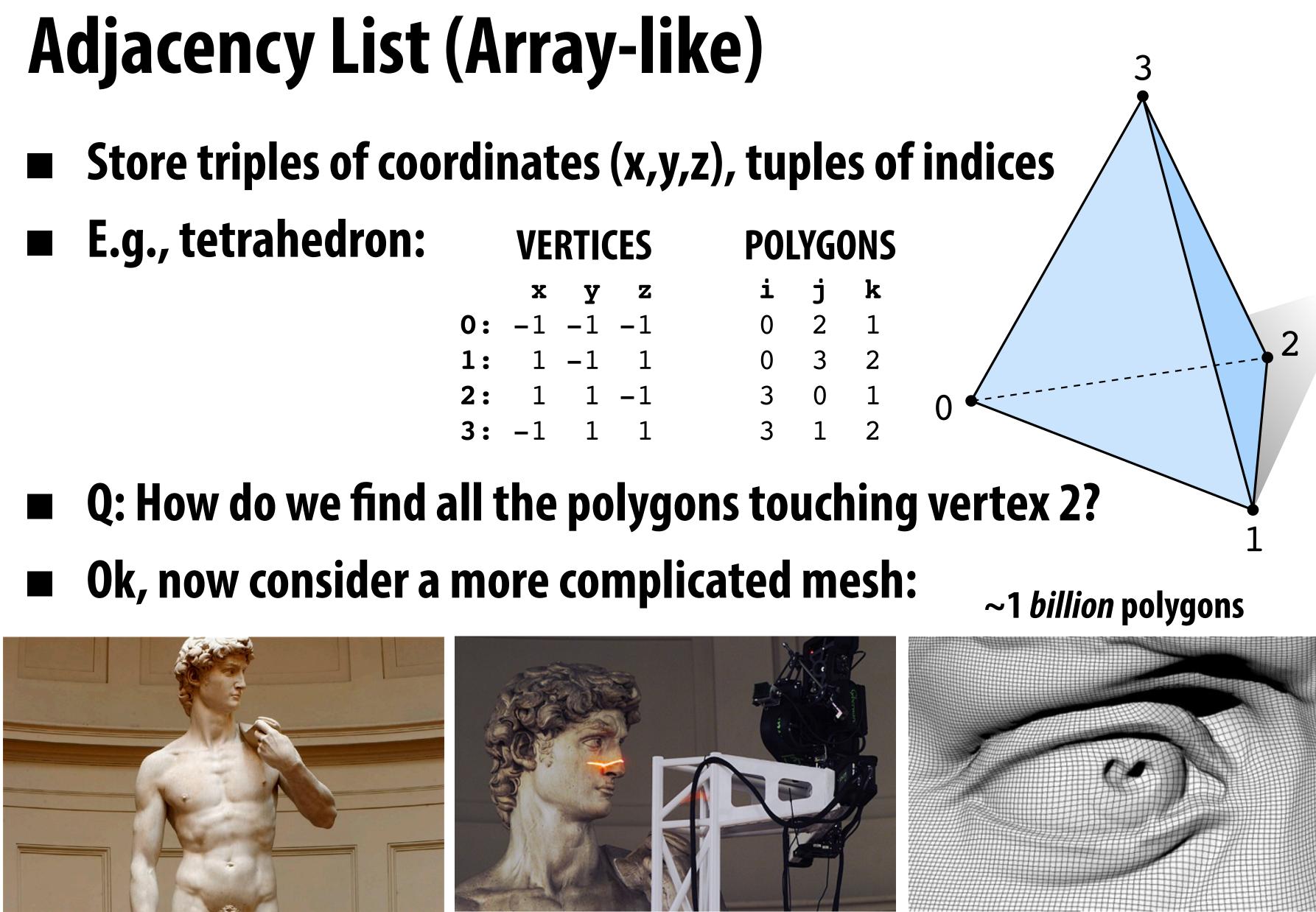
### Most basic idea:

- For each triangle, just store three coordinates
- No other information about connectivity
- Not much different from point cloud! ("Triangle cloud?")
- **Pros:** 
  - **Really stupidly simple**
- Cons:
  - **Redundant storage**
  - Hard to do much beyond simply drawing the mesh on screen
  - **Need spatial data structures** (later) to find neighbors





**3:** –1



### Very expensive to find the neighboring polygons! (What's the cost?)



## **Incidence** Matrices

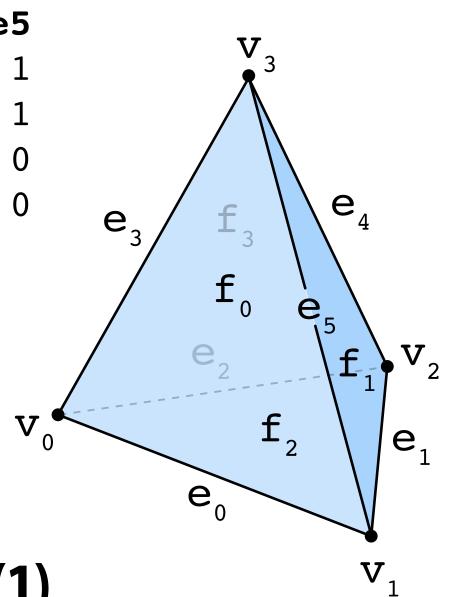
- **neighbors**?
- Can encode all neighbor information via *incidence matrices*

E.g., tetrahedron:	<u>VERTEX⇔EDC</u>				
5.	7	70	<b>v1</b>	<b>v2</b>	v
	e0	1	1	0	
	e1	0	1	1	
	e2	1	0	1	
	<b>e</b> 3	1	0	0	
	e4	0	0	1	
	e5	0	1	0	

- 1 means "touches"; 0 means "does not touch"
- Instead of storing lots of 0's, use sparse matrices
- Still large storage cost, but finding neighbors is now O(1)
- Hard to change connectivity, since we used fixed indices
- Bonus feature: mesh does not have to be manifold

### If we want to know who our neighbors are, why not just store a list of

**EDGE** ↔ **FACE** <u>GE</u> e1 e2 e3 73 e4 e5 **f3** 0

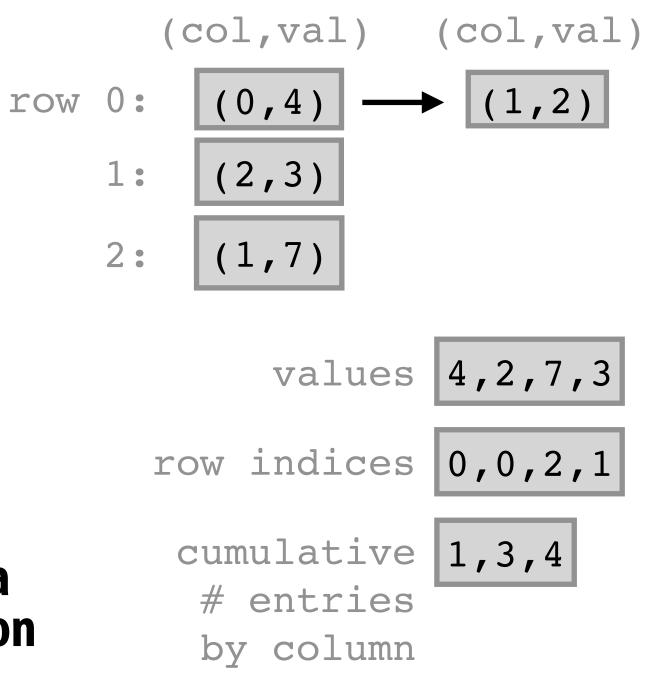




## **Aside: Sparse Matrix Data Structures**

- Ok, but how do we actually store a "sparse matrix"?
- Lots of possible data structures:
  - <u>Associative array</u> from (row, column) to value
    - easy to lookup/set entries, fast (e.g., hash table)
    - harder to do matrix operations (e.g., multiplication)
  - <u>Array of linked lists</u> (one per row)
    - conceptually simple
    - slow access time, incoherent memory access
  - <u>Compressed column format</u>—pack entries in list
    - hard to add/modify entries
    - fast for actual matrix operations
- In practice: often build up entries using an "easier" data structure, convert to compressed format for computation

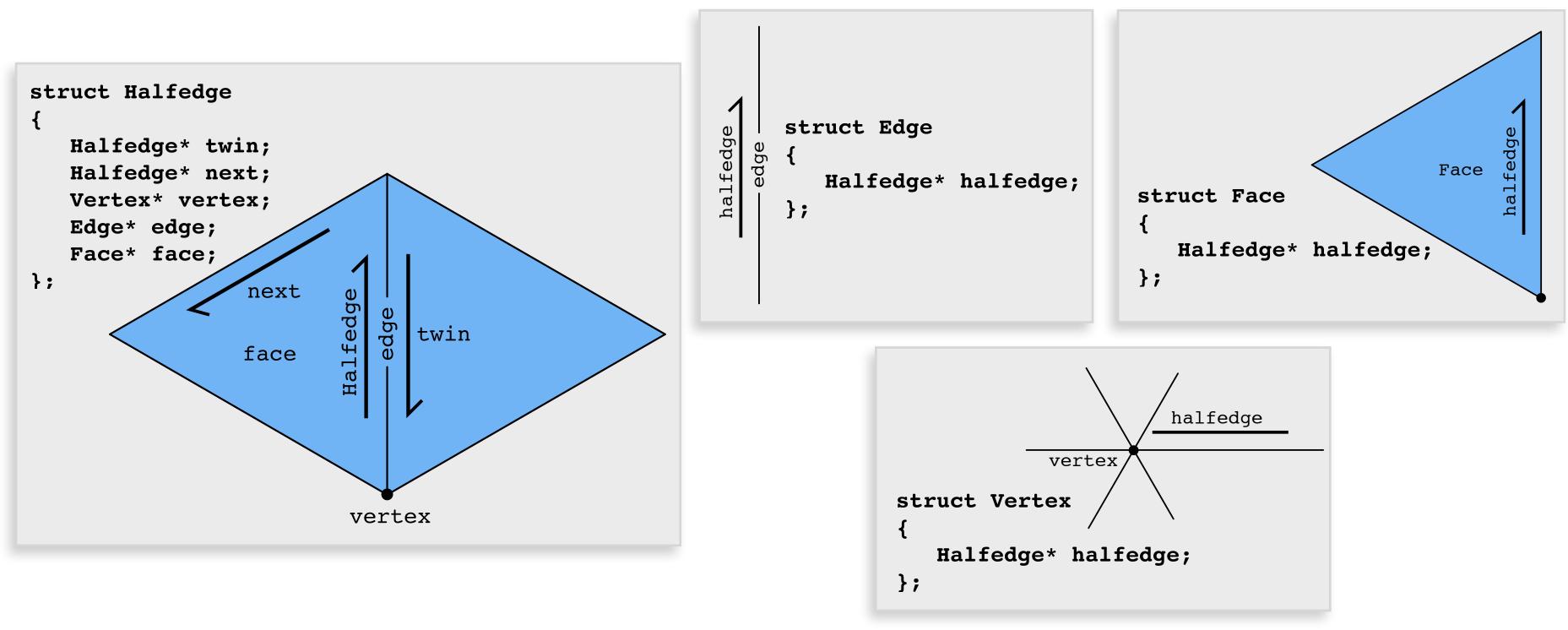
(row,col) val (0,0) -> 4(0,1) -> 2(1,2) -> 3(2,1) -> 7





# Halfedge Data Structure (Linked-list-like)

### Store some information about neighbors Don't need an exhaustive list; just a few key pointers Key idea: two *halfedges* act as "glue" between mesh elements:



### Each vertex, edge face points to just one of its halfedges.



## Halfedge makes mesh traversal easy

## Use "twin" and "next" pointers to move around mesh

### **Example: visit all vertices of a face:**

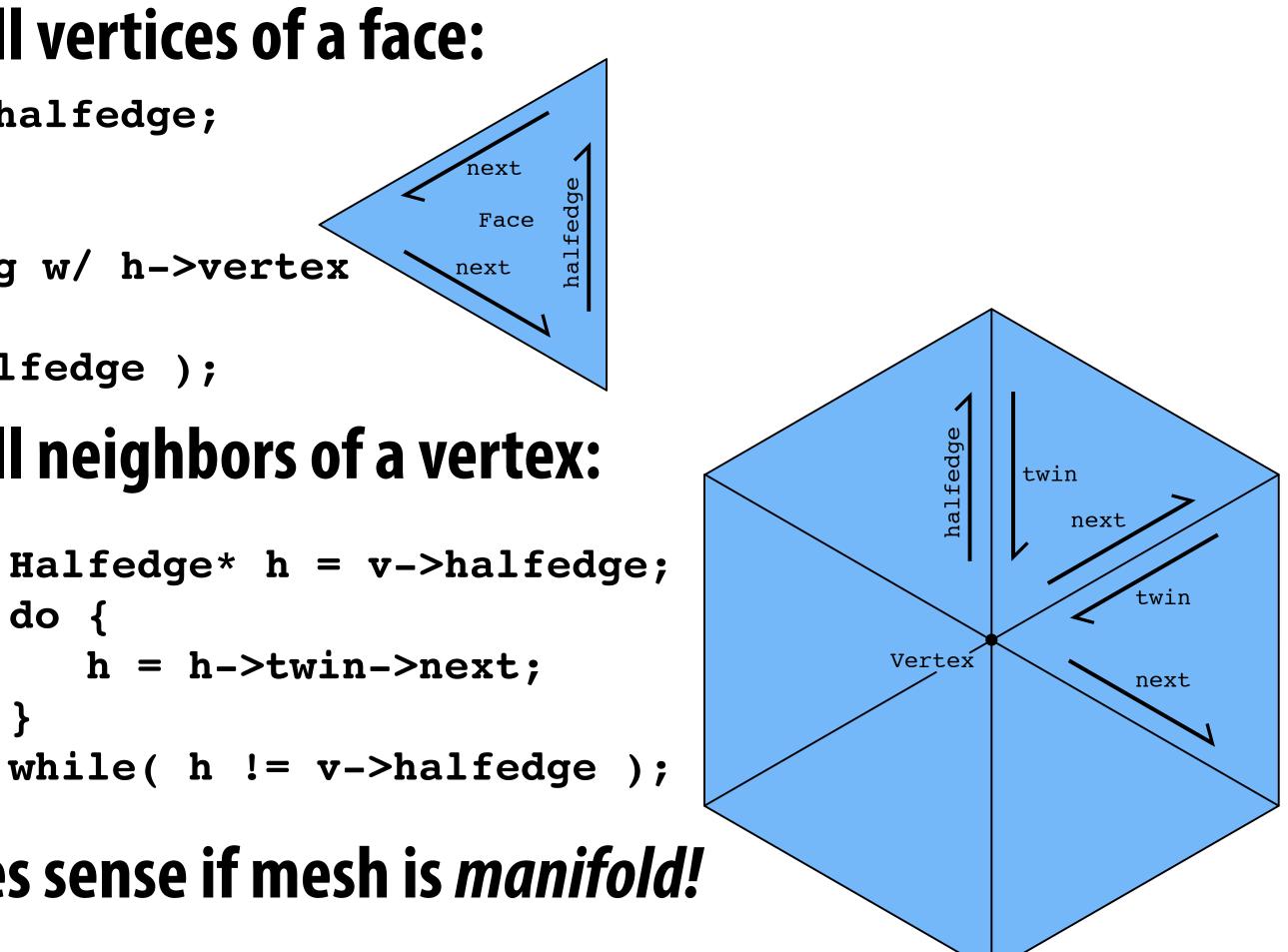
Halfedge\* h = f->halfedge; do h = h - next;// do something w/ h->vertex while( h != f->halfedge );

### Example: visit all neighbors of a vertex:

do

### Note: only makes sense if mesh is manifold!

### Use "vertex", "edge", and "face" pointers to grab element

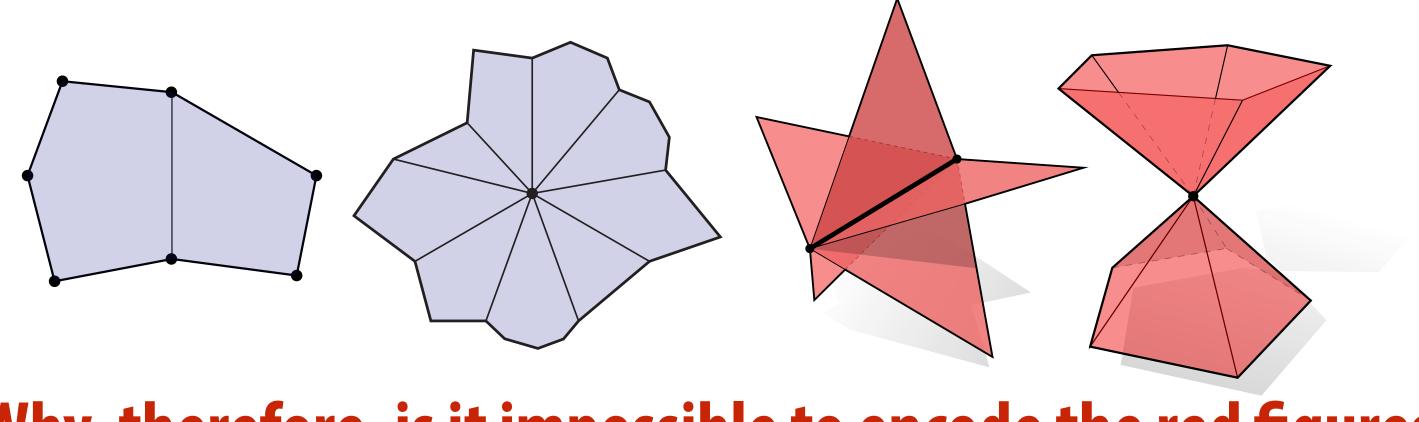




## Halfedge connectivity is always manifold Consider simplified halfedge data structure **Require only "common-sense" conditions**

struct Halfedge { Halfedge \*next, \*twin; };

### Keep following next, and you'll get faces. Keep following twin and you'll get edges. Keep following next->twin and you'll get vertices.



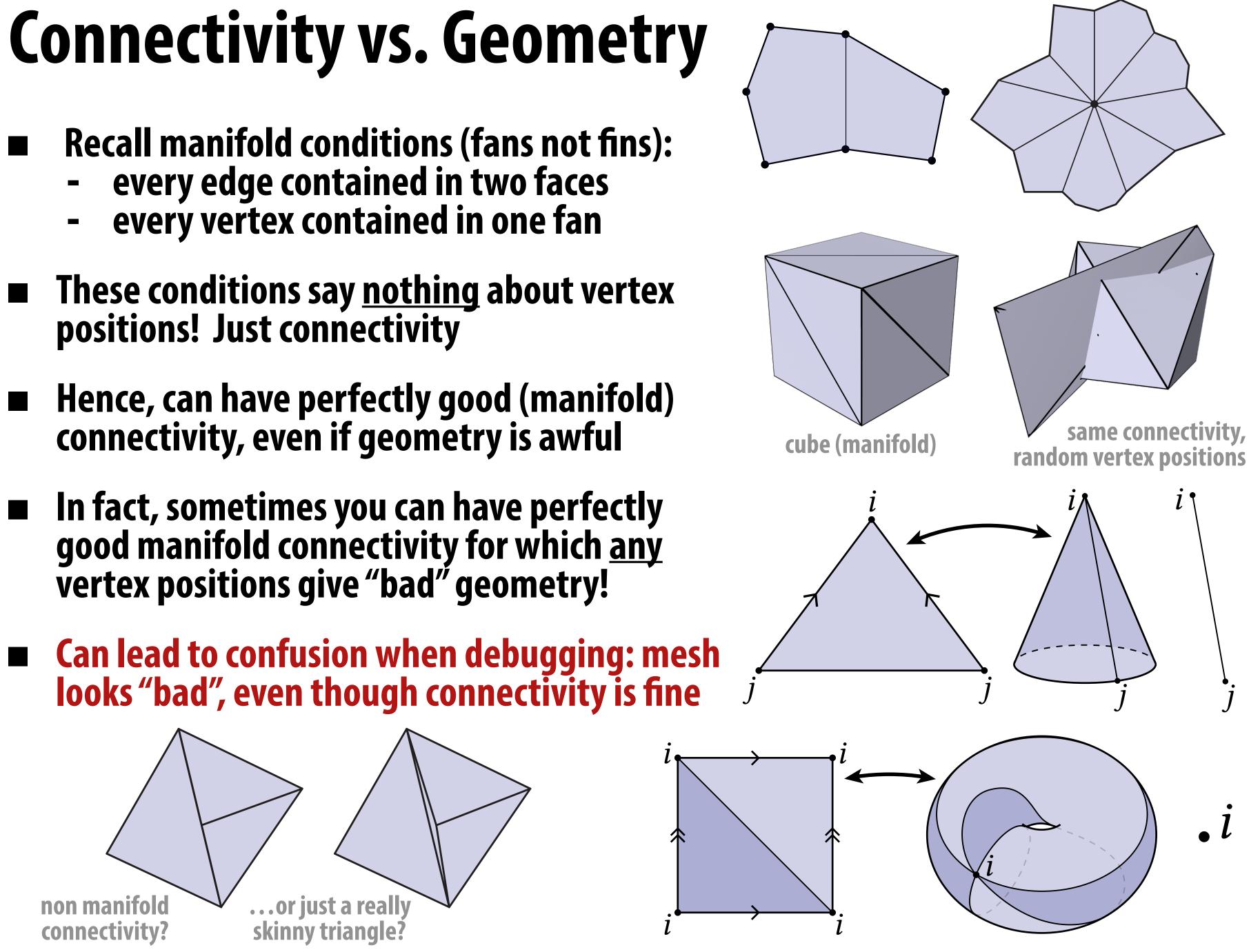
Q: Why, therefore, is it impossible to encode the red figures?

(pointer to yourself!)

twin->twin == this twin != this every he is someone's "next"

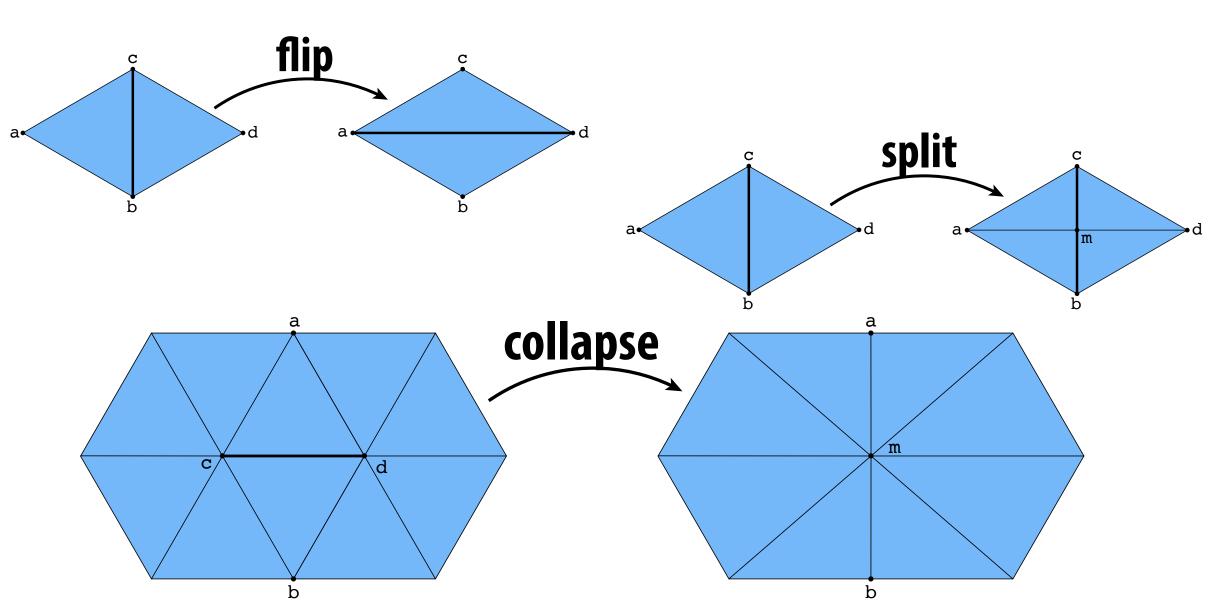


- positions! Just connectivity



## Halfedge meshes are easy to edit

- **Remember key feature of linked list: insert/delete elements** Same story with halfedge mesh ("linked list on steroids") E.g., for triangle meshes, several atomic operations:

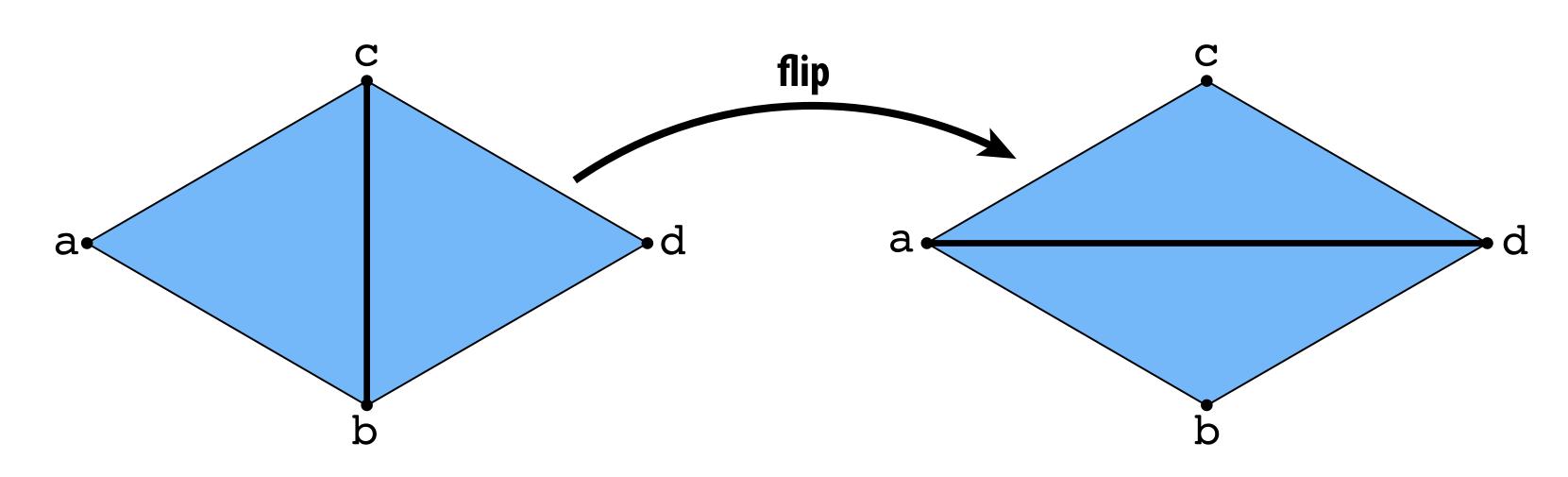


How? Allocate/delete elements; reassigning pointers. Must be careful to preserve manifoldness!



# **Edge Flip**

**Triangles (a,b,c), (b,d,c) become (a,d,c), (a,b,d):** 

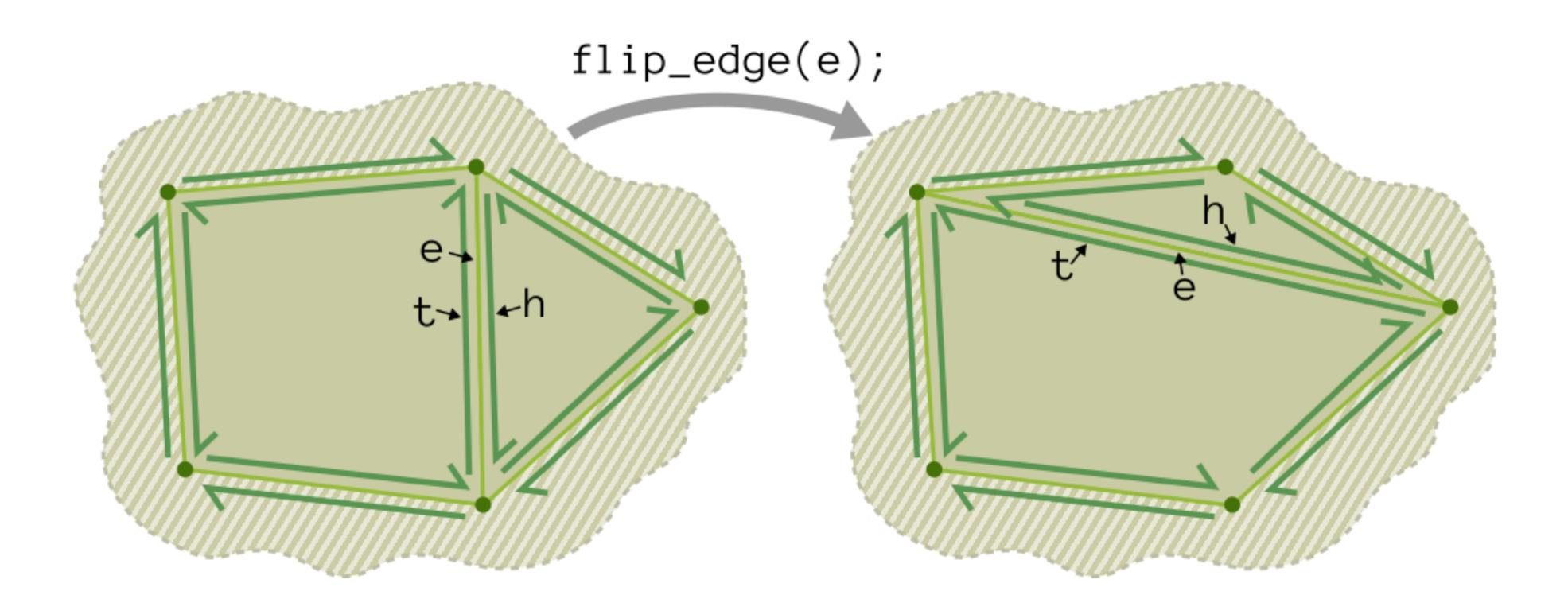


- Long list of pointer reassignments (edge->halfedge = ...)
- However, no elements created/destroyed.
- Q: What happens if we flip twice?

Challenge: can you implement edge flip such that pointers are *unchanged* after two flips? 



# Edge Flip Can generalize:

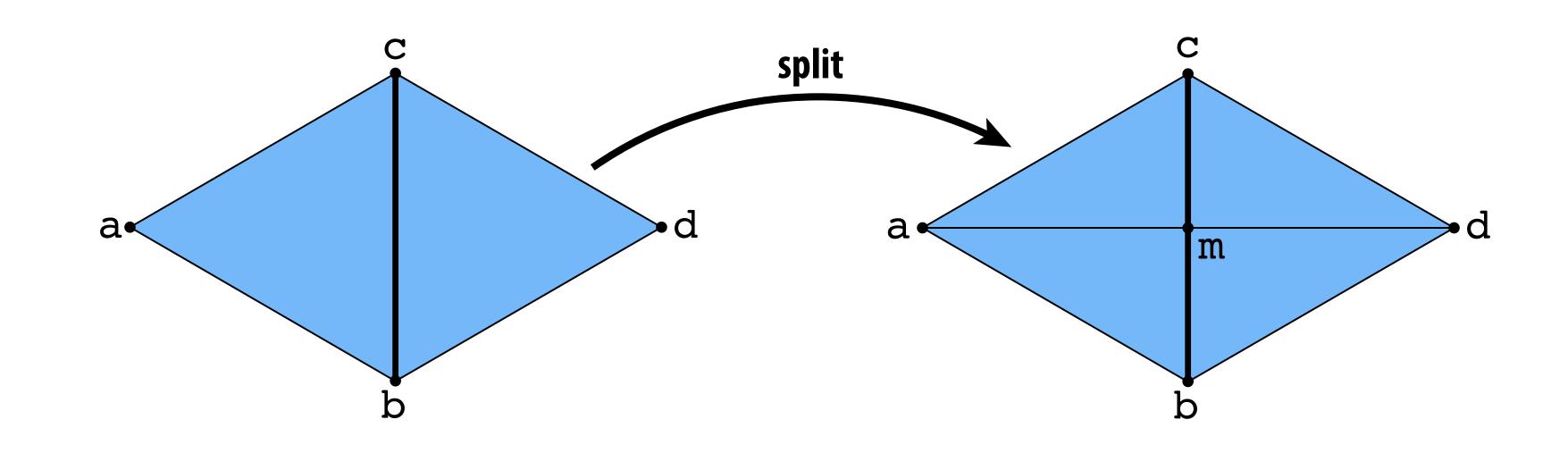


### How many flips until back at the start?



## **Edge Split**

### Insert midpoint *m* of edge (*c*, *d*), connect to get four faces:

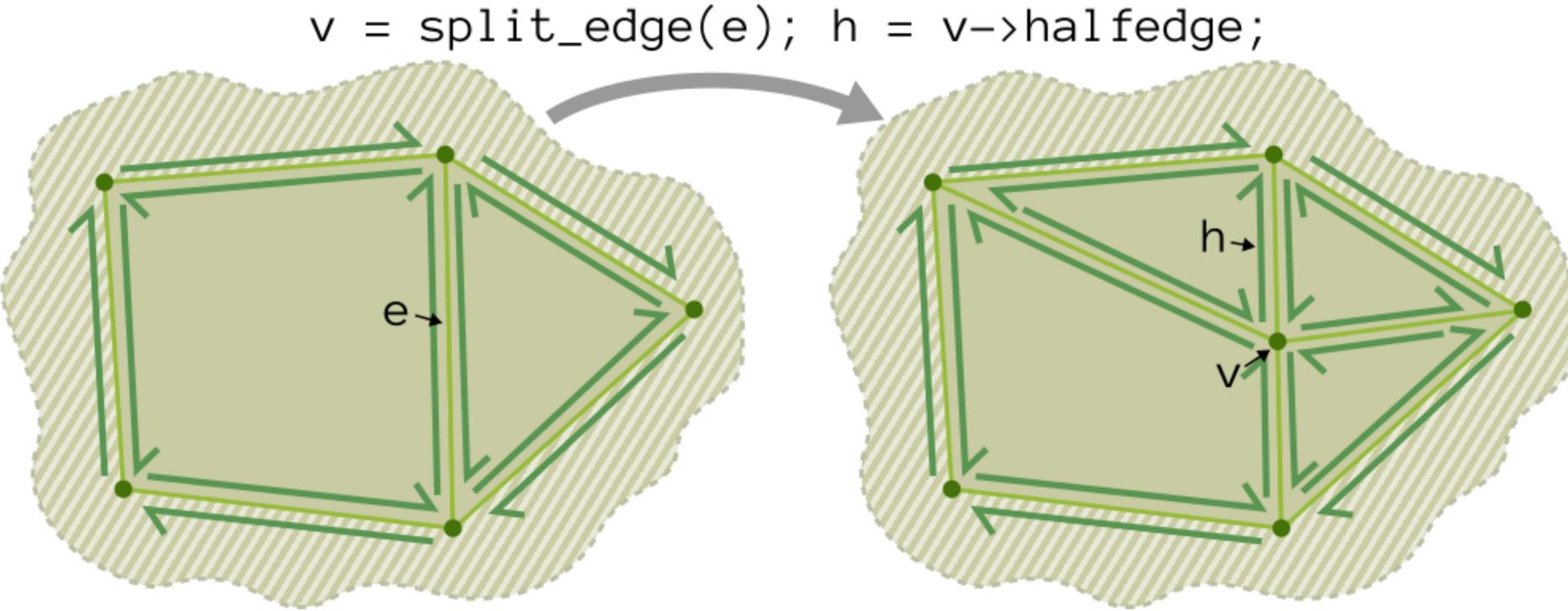


### This time, have to *add* new elements. Lots of pointer reassignments. Q: Can we "reverse" this operation?



## Edge Split

### Works on faces of any size:

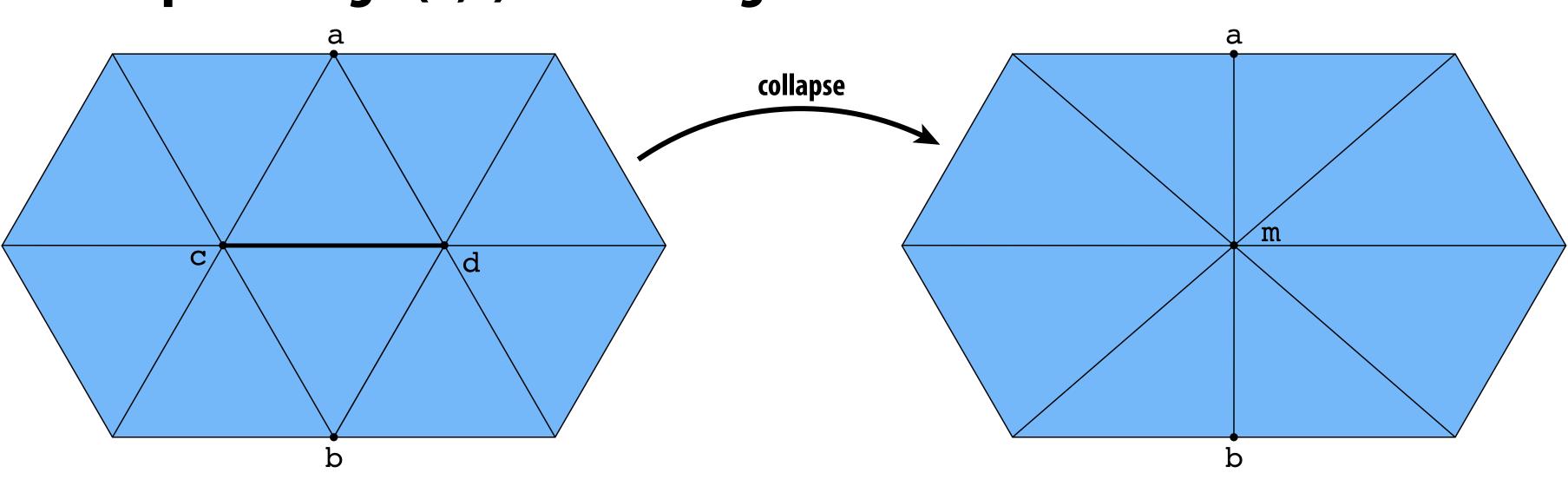


Where to split?



# **Edge Collapse**

### **Replace edge (b,c) with a single vertex m:**



- Now have to *delete* elements.
- **Still lots of pointer assignments!**

### Q: How would we implement this with an adjacency list?

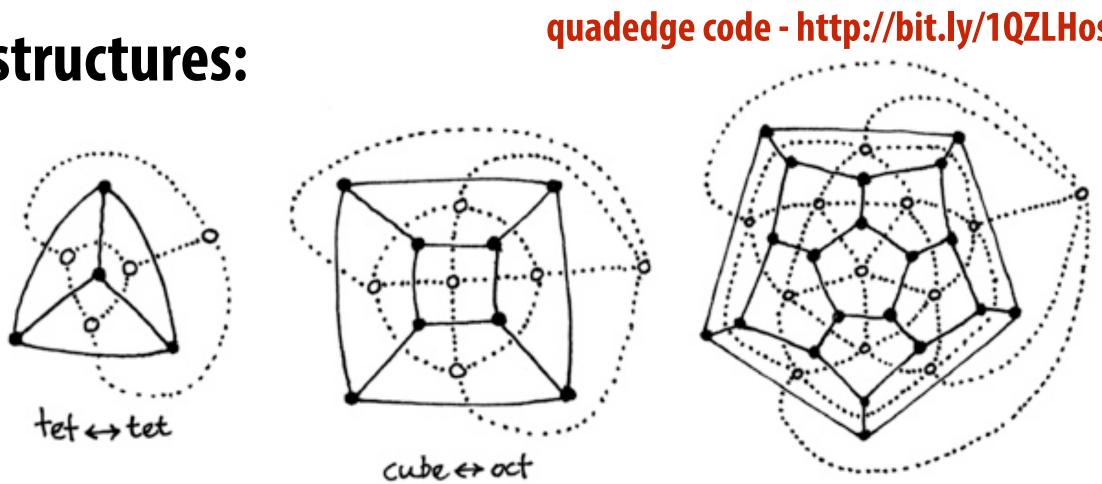
### Any other good way to do it? (E.g., different data structure?)



# **Alternatives to Halfedge**

### Many very similar data structures:

- winged edge
- corner table
- quadedge



- **Each stores local neighborhood information**
- Similar tradeoffs relative to simple polygon list:
  - **CONS**: additional storage, incoherent memory access
  - traversal of local neighborhoods

\*see for instance http://geometry-central.net/

### Paul Heckbert (former CMU prof.) quadedge code - http://bit.ly/1QZLHos

dodec (+ icos

# **PROS**: better access time for individual elements, intuitive

# With some thought\*, <u>can</u> design halfedge-type data structures with coherent data storage, support for non manifold connectivity, etc.



## **Comparison of Polygon Mesh Data Strucutres**

	Adjacency List	Incidence Matrices	Halfedge Mesh
constant-time neighborhood access?	ΝΟ	YES	YES
easy to add/remove mesh elements?	NO	NO	YES
nonmanifold geometry?	YES	YES	NO

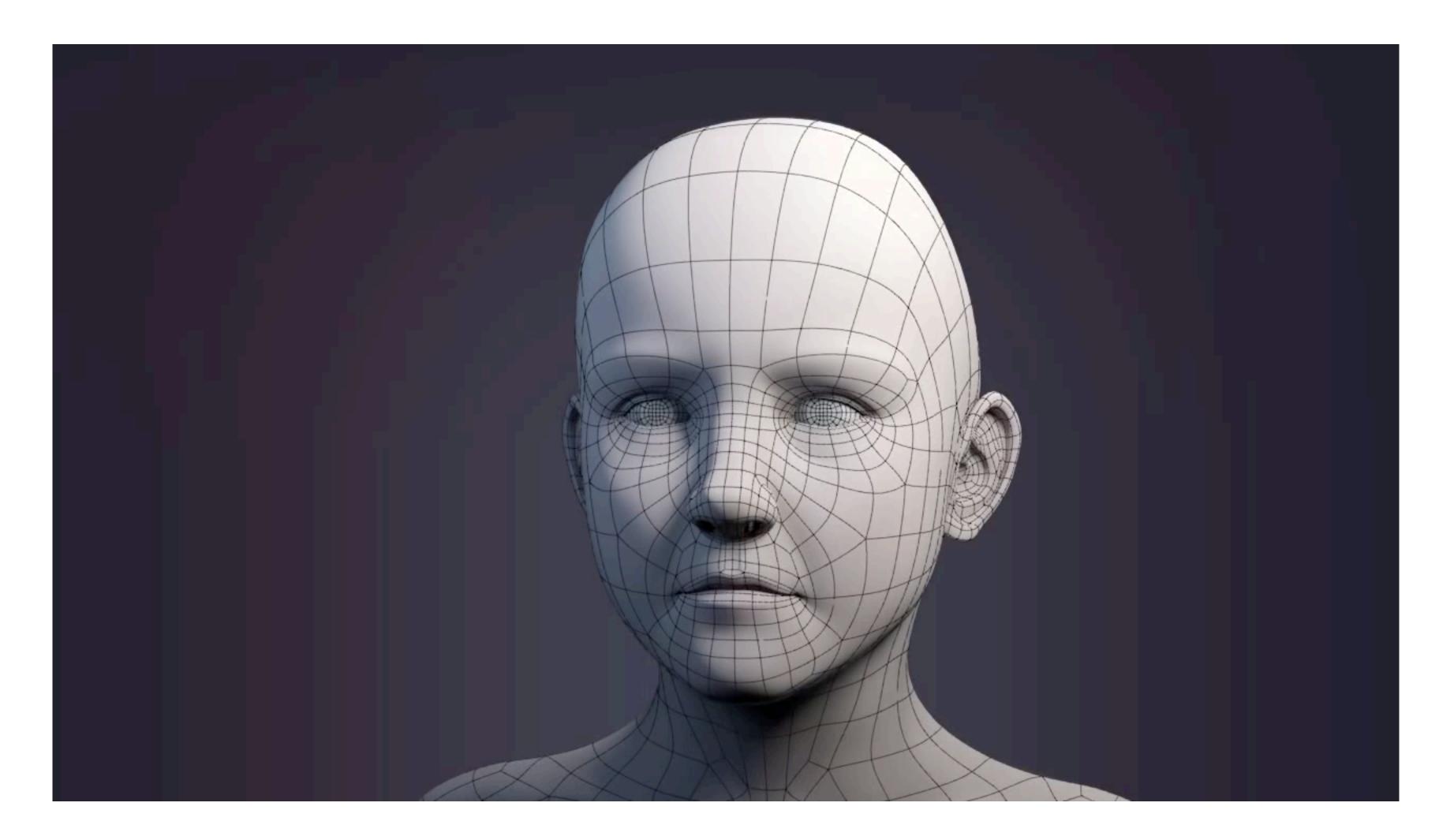
### **Conclusion: pick the right data structure for the job!**



## Ok, but what can we actually do with our fancy new data structures?



## Subdivision Modeling

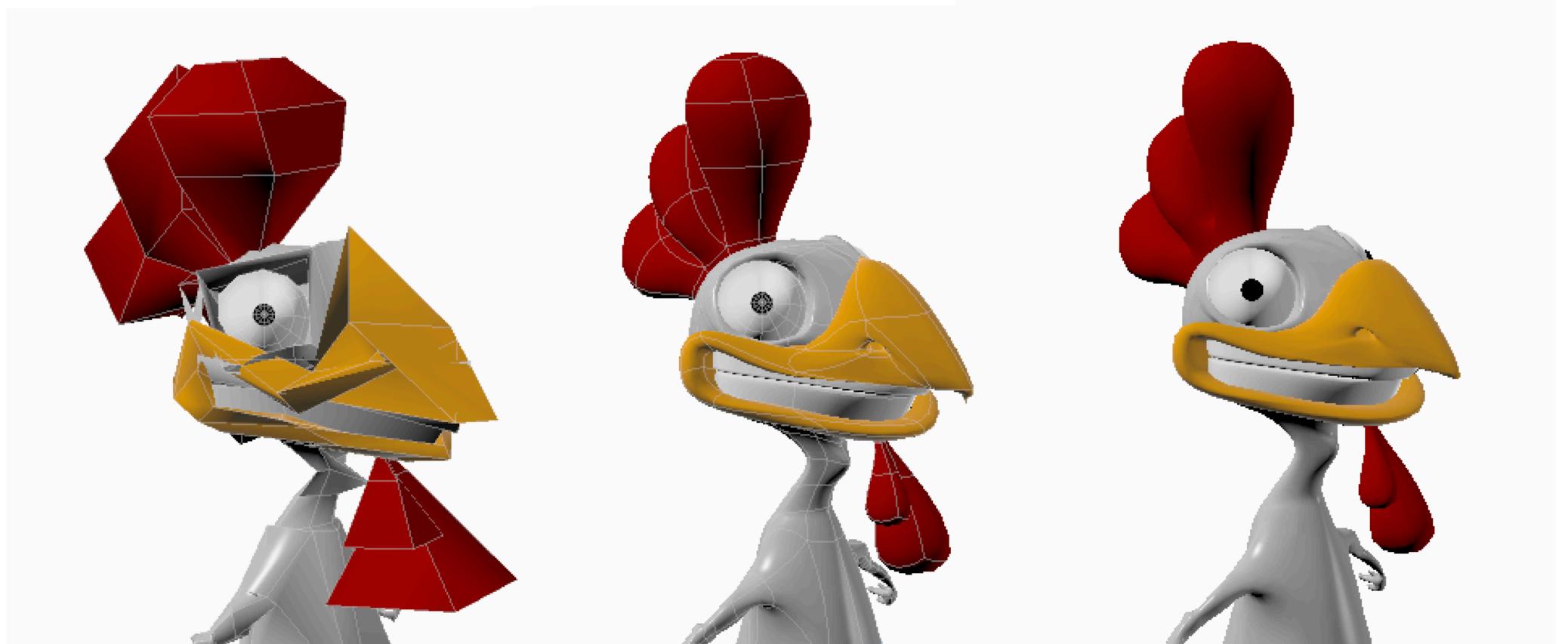






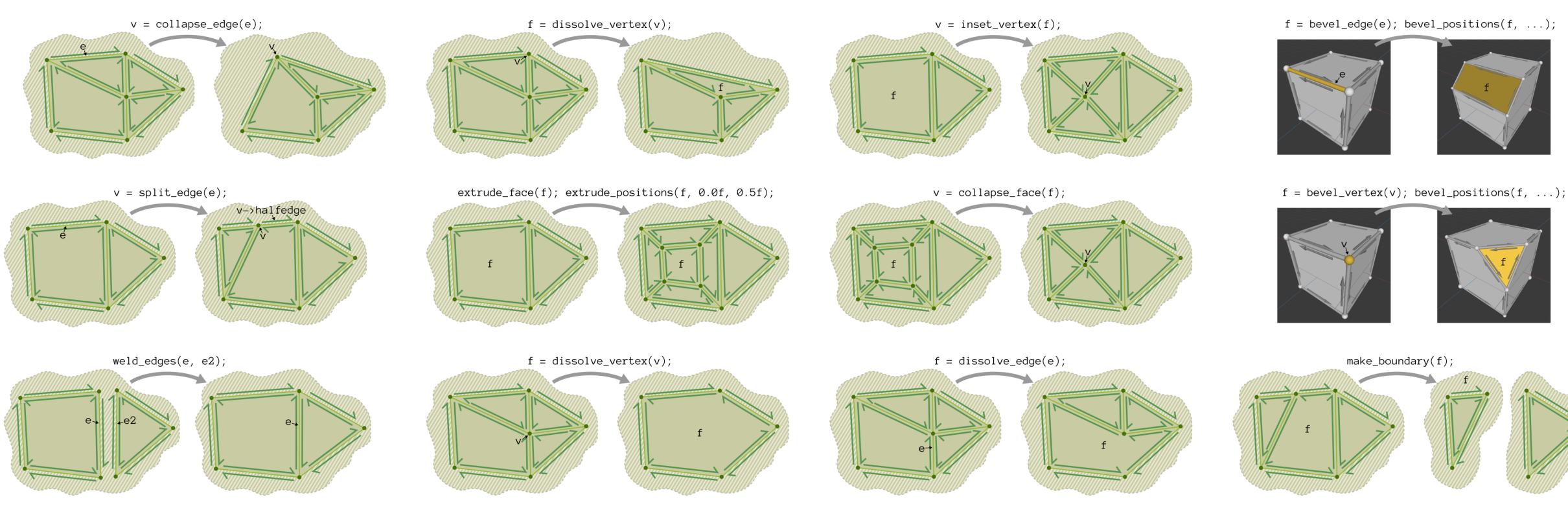
## Subdivision Modeling Common modeling paradigm in modern 3D tools:

- Coarse "control cage"
- Perform local operations to control/edit shape
- Global subdivision process determines final surface





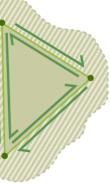
# Subdivision Modeling—Local Operations For general polygon meshes, we can dream up lots of local mesh operations that might be useful for modeling:



### ...and many, many more!









# **Next Time: Digital Geometry Processing**

- Extend traditional digital signal processing (audio, video, etc.) to deal with *geometric* signals:
  - upsampling / downsampling / resampling / filtering ...
  - aliasing (reconstructed surface gives "false impression")

