

# **Meshes and Manifolds**

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**Computer Graphics**  
**CMU 15-462/15-662**

# Last time: overview of geometry

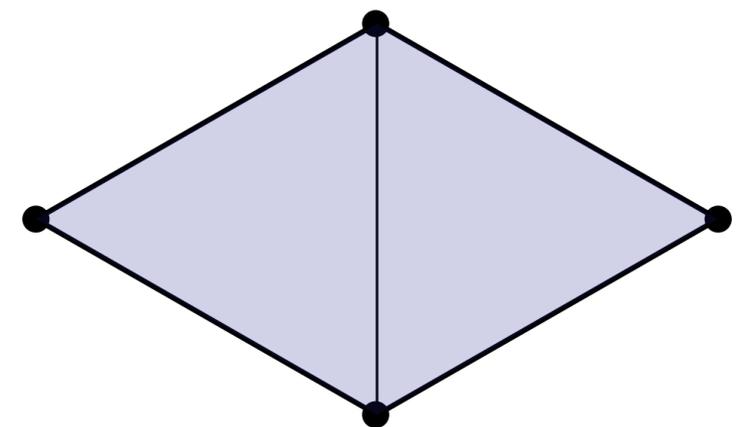
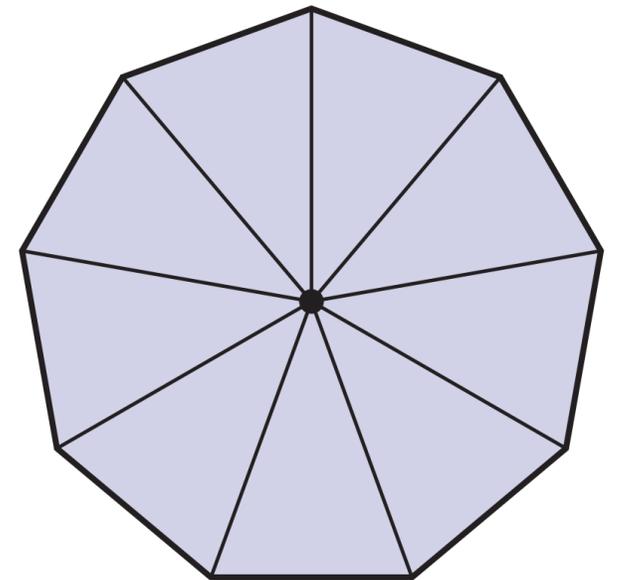
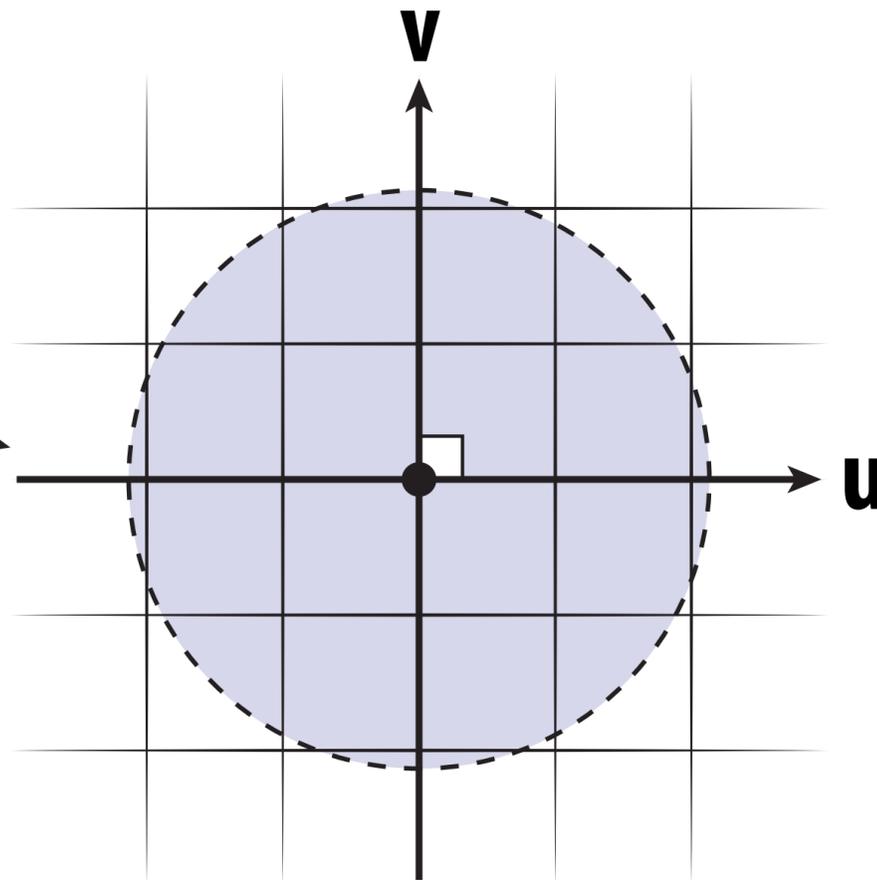
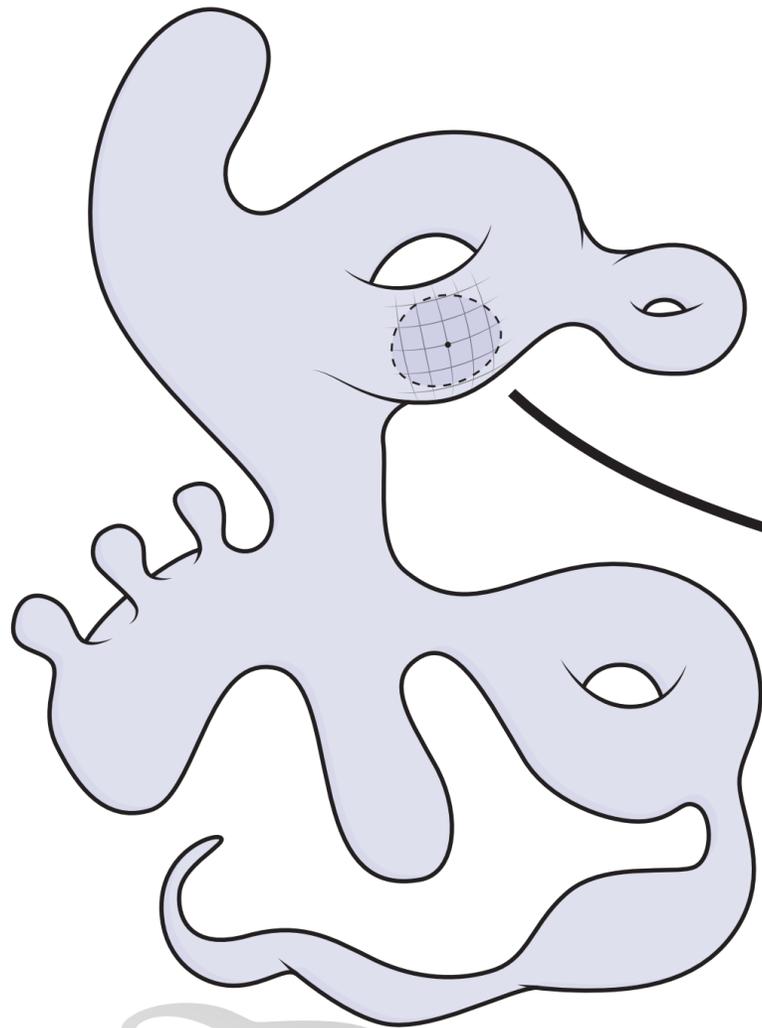
- Many types of geometry in nature
- Demand sophisticated representations
- Two major categories:
  - **IMPLICIT** - “tests” if a point is in shape
  - **EXPLICIT** - directly “lists” points
- Lots of representations for both
- Today:
  - what is a surface, anyway?
  - nuts & bolts of polygon meshes
  - geometry processing / resampling

## Geometry



# Manifold Assumption

- Today we're going to introduce the idea of *manifold* geometry
- Can be hard to understand motivation at first!
- So first, let's revisit a more familiar example...



# Bitmap Images, Revisited

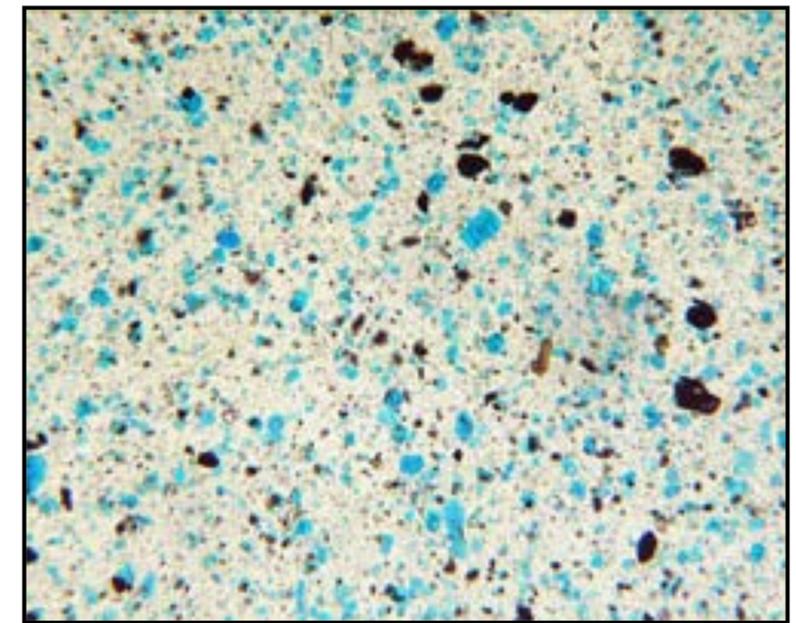
To encode images, we used a *regular grid of pixels*:



# But images are not fundamentally made of little squares:

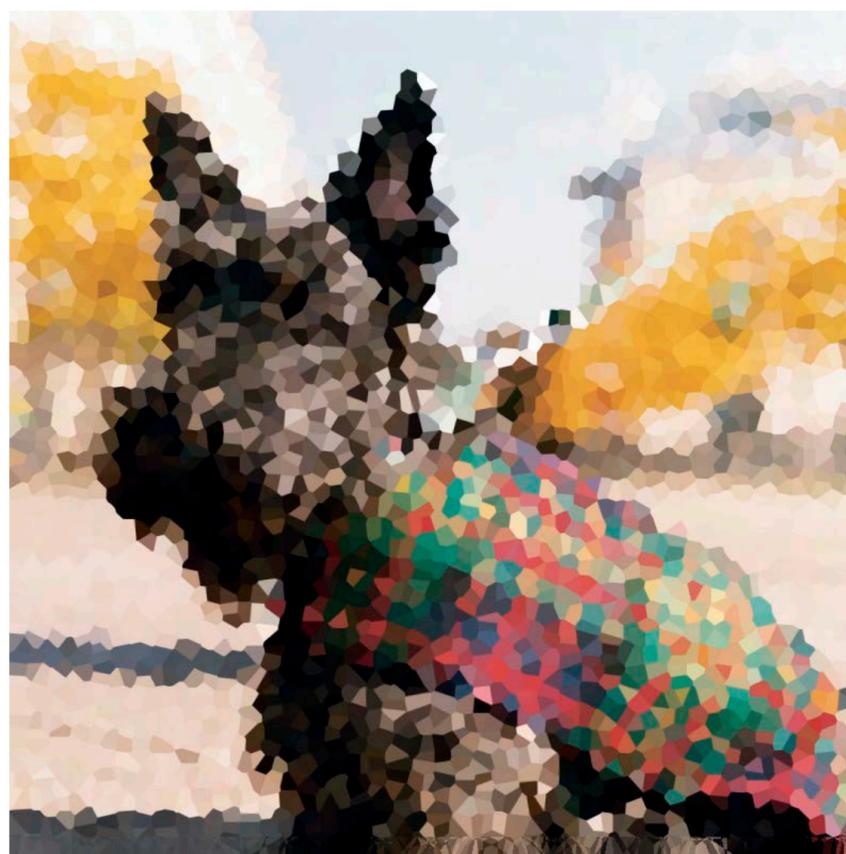
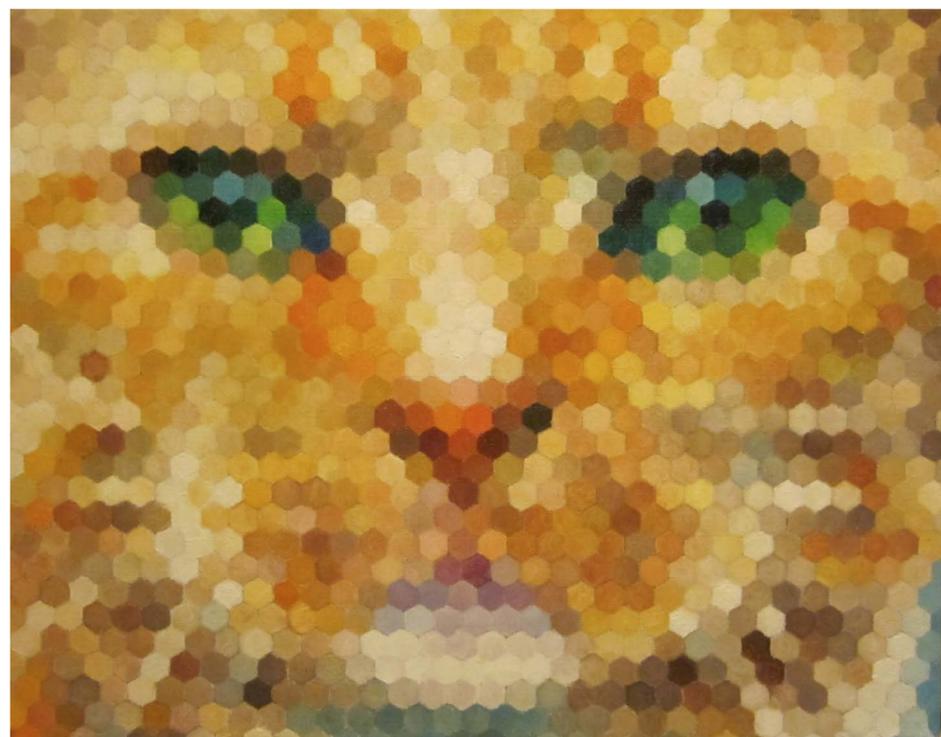


Goyō Hashiguchi, *Kamisuki* (ca 1920)



photomicrograph of paint

# So why did we choose a square grid?



...rather than dozens of possible alternatives?

# Regular grids make life easy

- **One reason: SIMPLICITY / EFFICIENCY**

- E.g., always have four neighbors
- Easy to index, easy to filter...
- Storage is just a list of numbers

- **Another reason: GENERALITY**

- Can encode basically any image

- **Are regular grids *always* the best choice for bitmap images?**

- No! E.g., suffer from anisotropy, don't capture edges, ...
- But *more often than not* are a pretty good choice

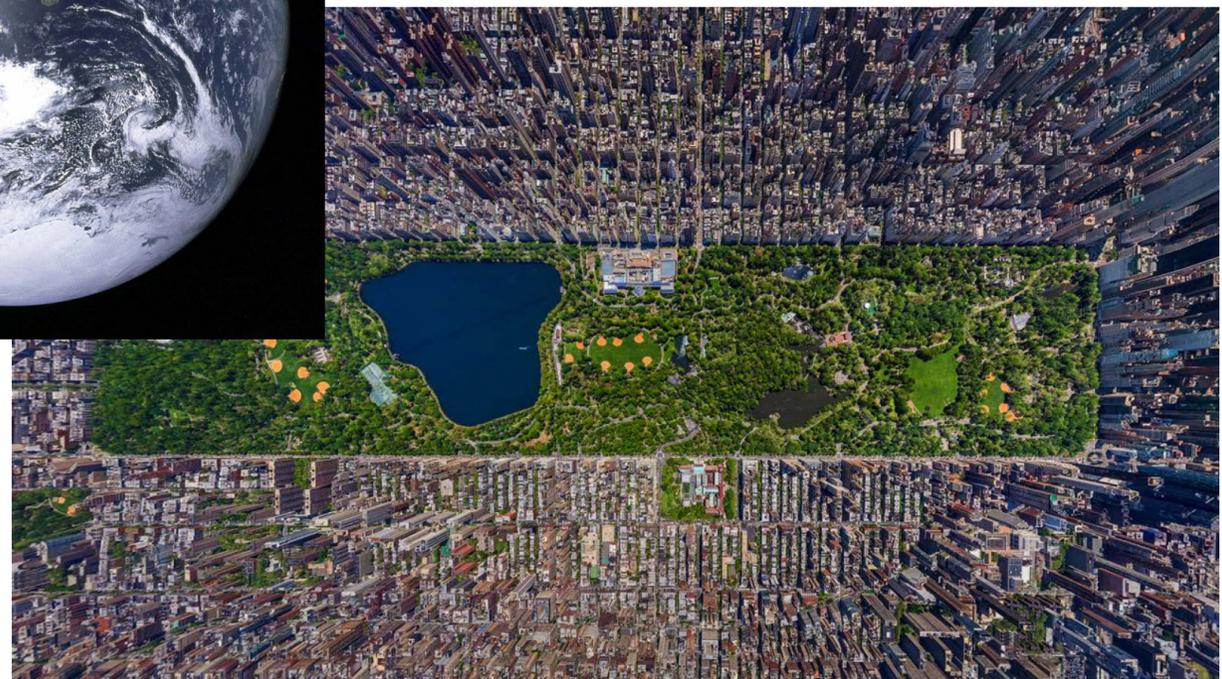
- **Will see a similar story with geometry...**

	$(i, j-1)$	
$(i-1, j)$	$(i, j)$	$(i+1, j)$
	$(i, j+1)$	

**So, how should we encode surfaces?**

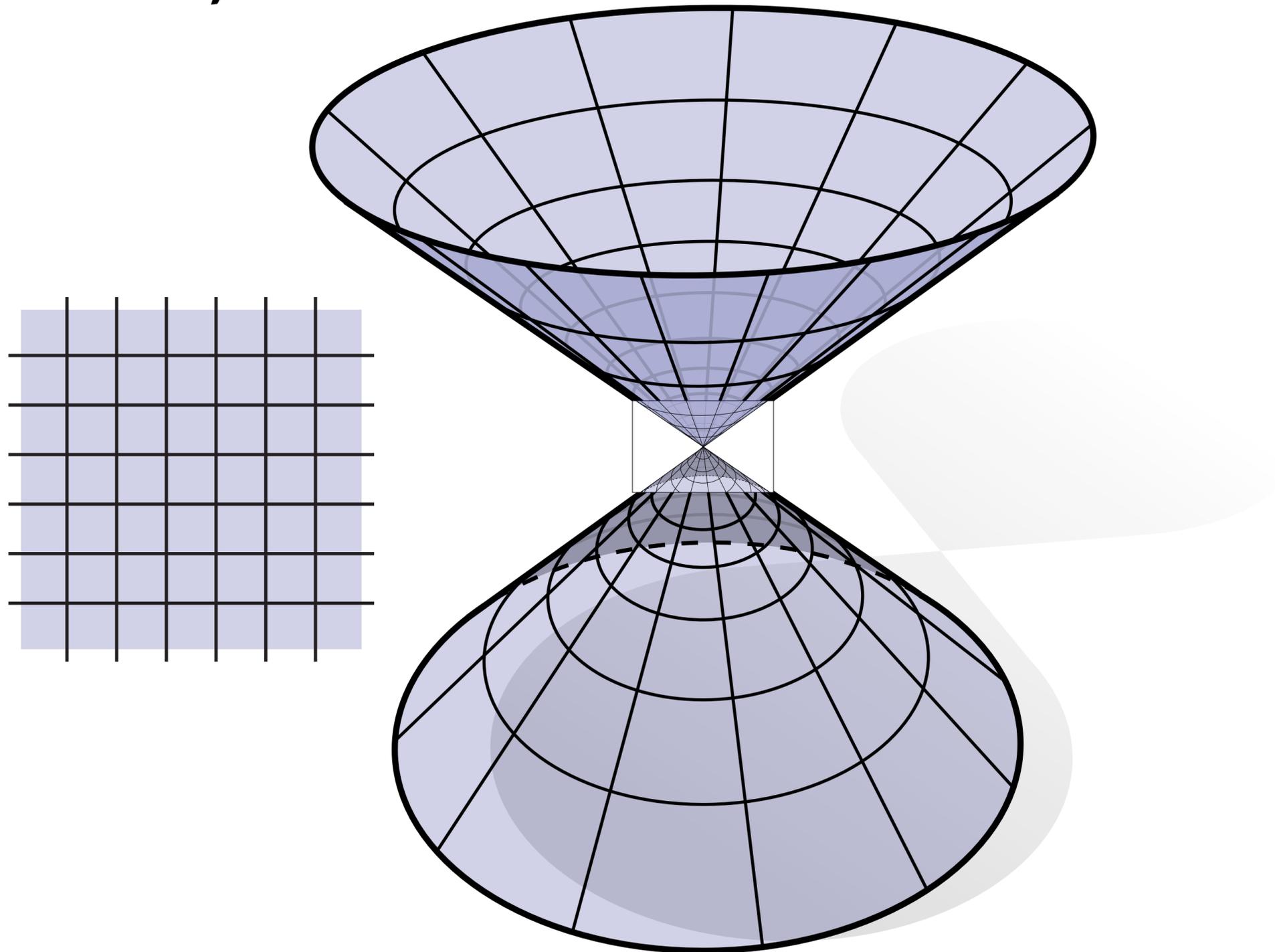
# Smooth Surfaces

- Intuitively, a surface is the boundary or “shell” of an object
- (Think about the candy shell, not the chocolate.)
- Surfaces are *manifold*:
  - If you zoom in far enough, can draw a regular coordinate grid
  - E.g., the Earth from space vs. from the ground



# Isn't every shape manifold?

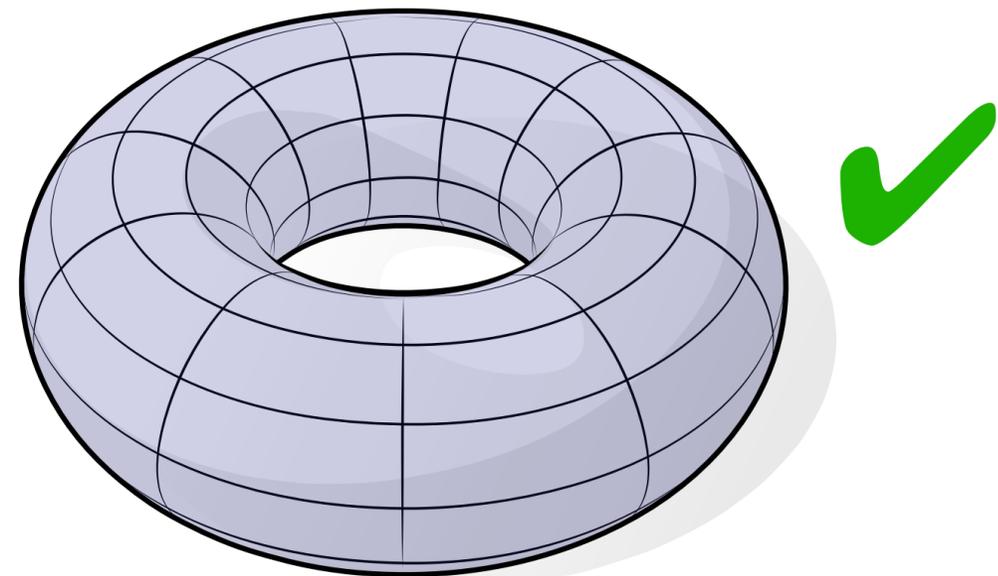
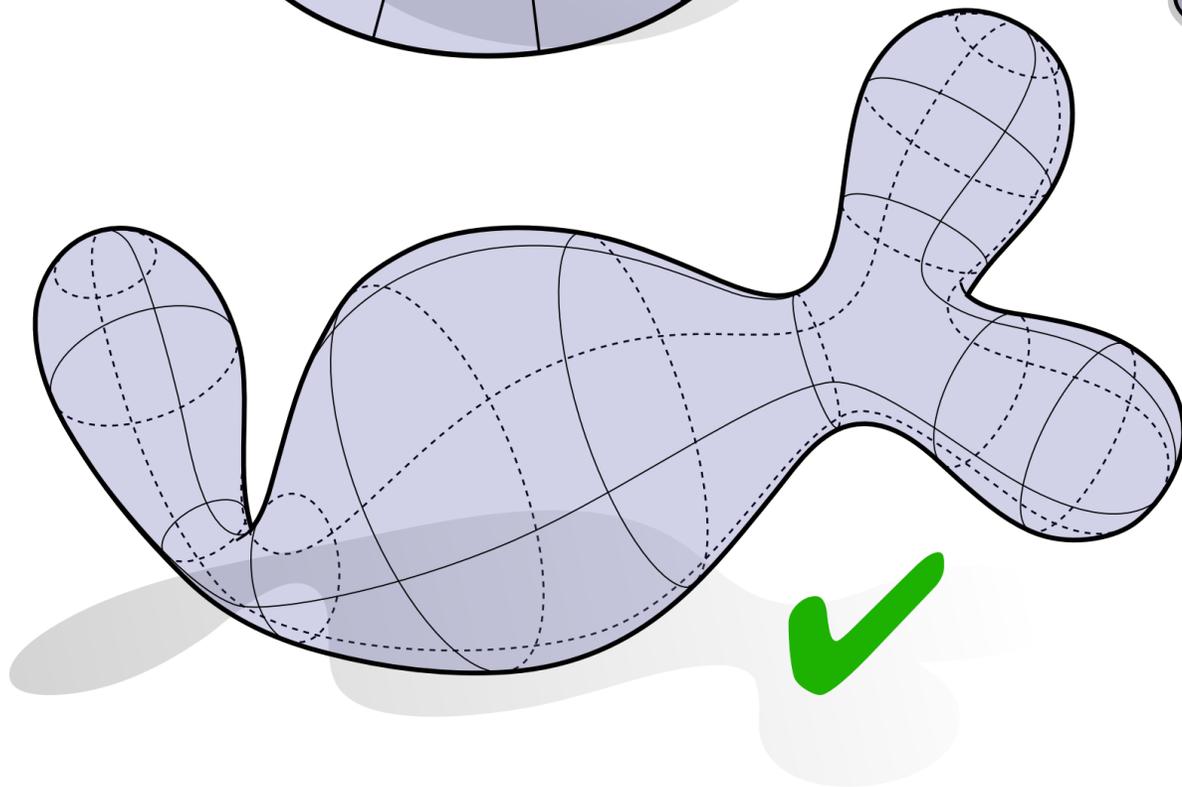
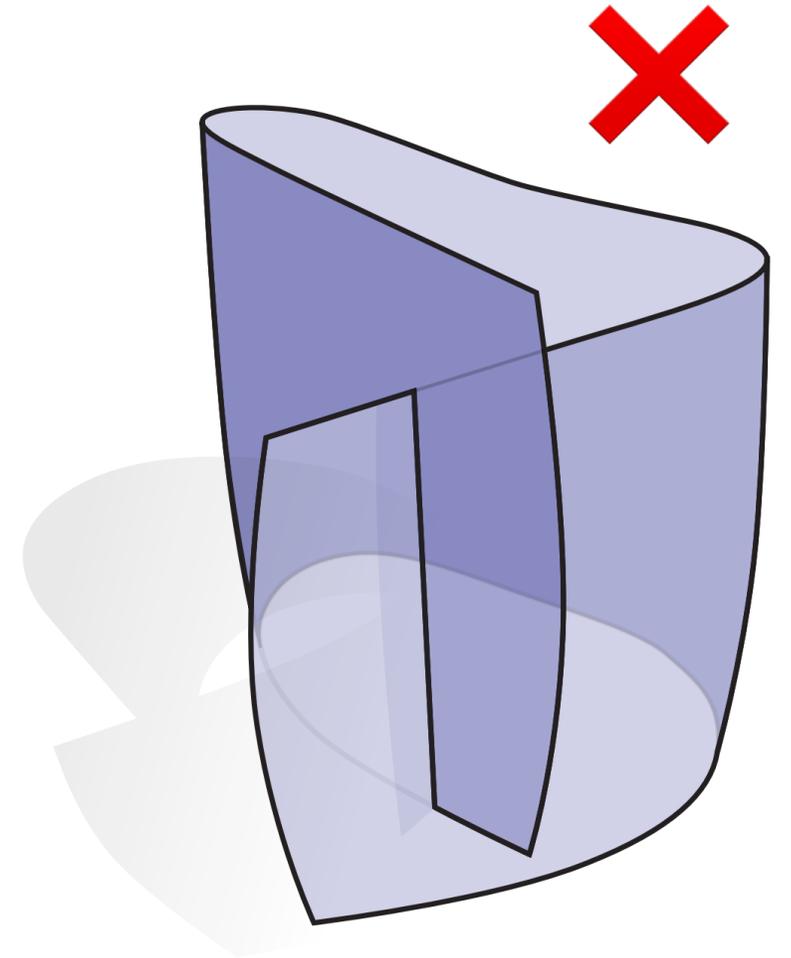
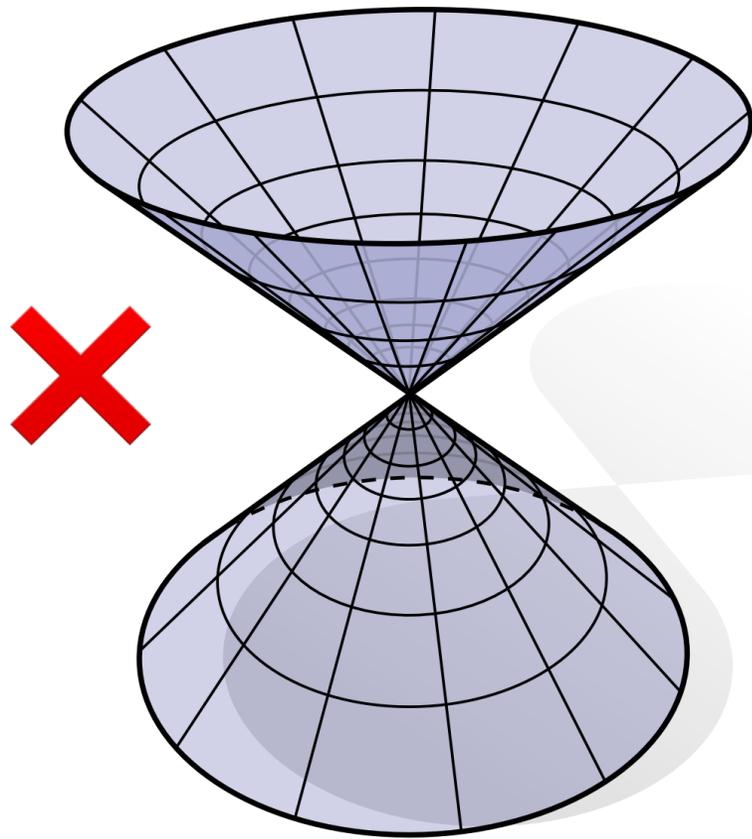
- No, for instance:



**Can't draw ordinary 2D grid at center, no matter how close we get.**

# Examples—Manifold vs. Nonmanifold

- Which of these shapes are manifold?



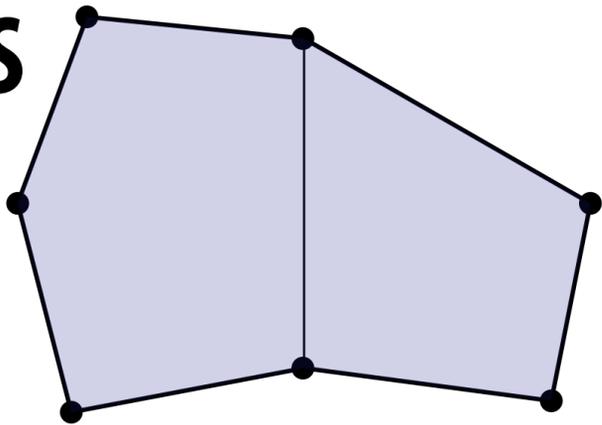
# A manifold polygon mesh has fans, not fins

- For polygonal surfaces just two easy conditions to check:

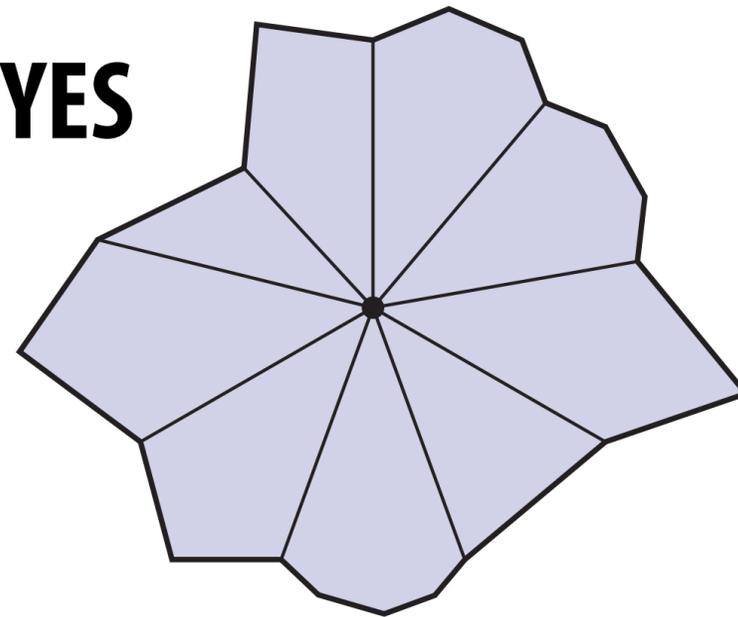
1. Every edge is contained in only two polygons (no “fins”)

2. The polygons containing each vertex make a single “fan”

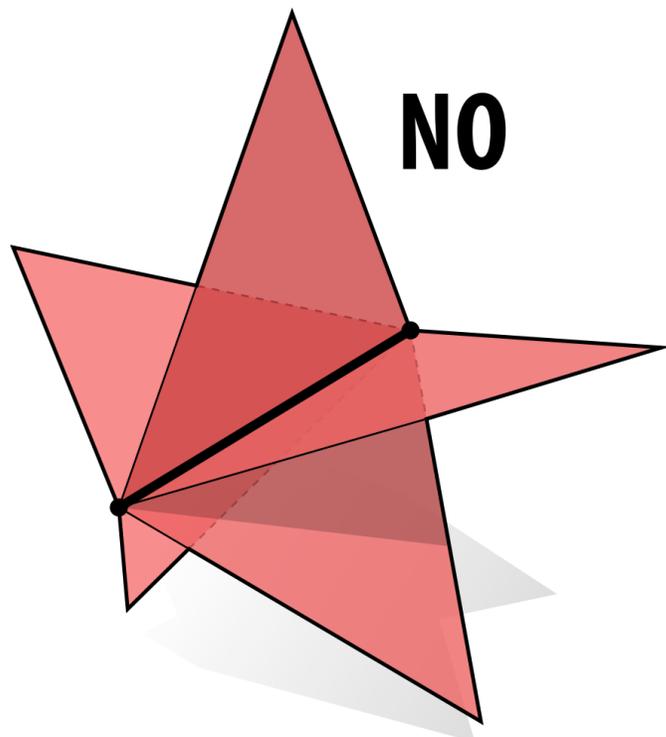
YES



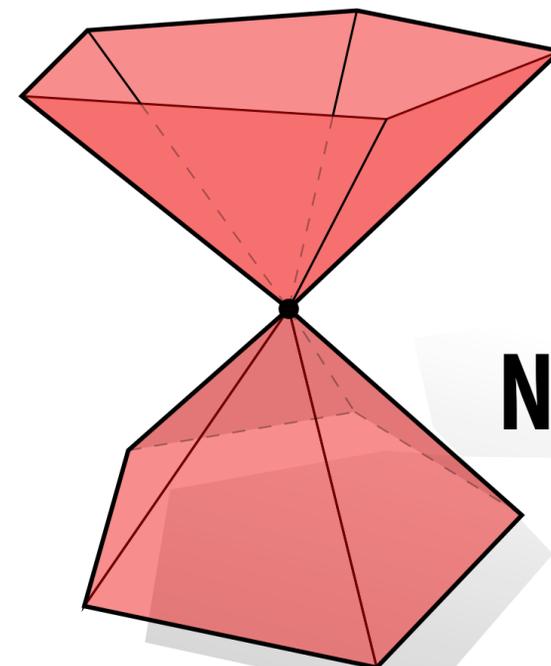
YES



NO

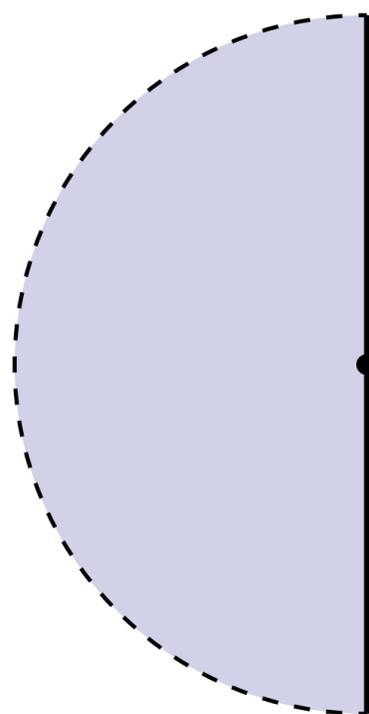
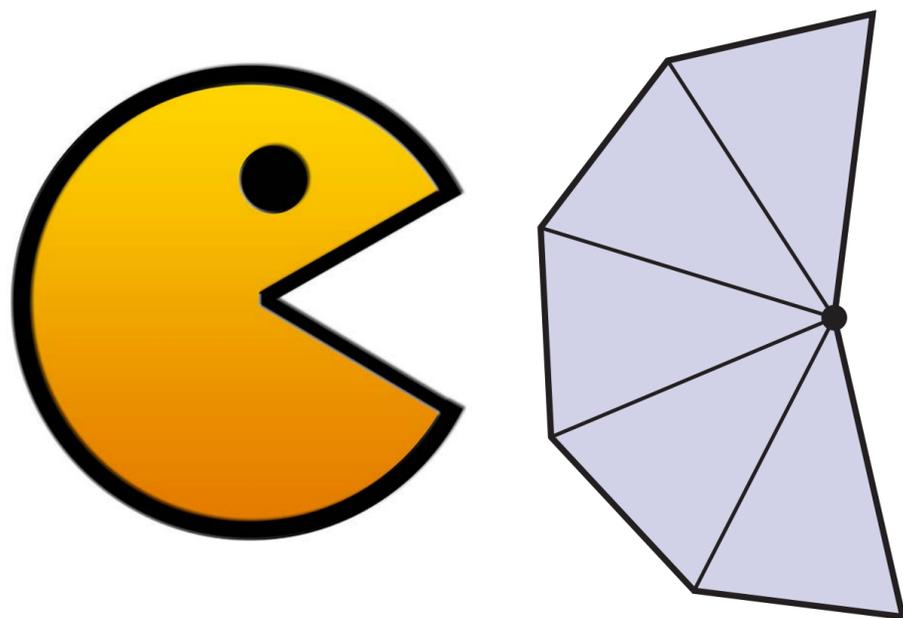


NO



# What about boundary?

- The boundary is where the surface “ends.”
- E.g., waist & ankles on a pair of pants.
- Locally, looks like a *half* disk
- Globally, each boundary forms a loop



YES



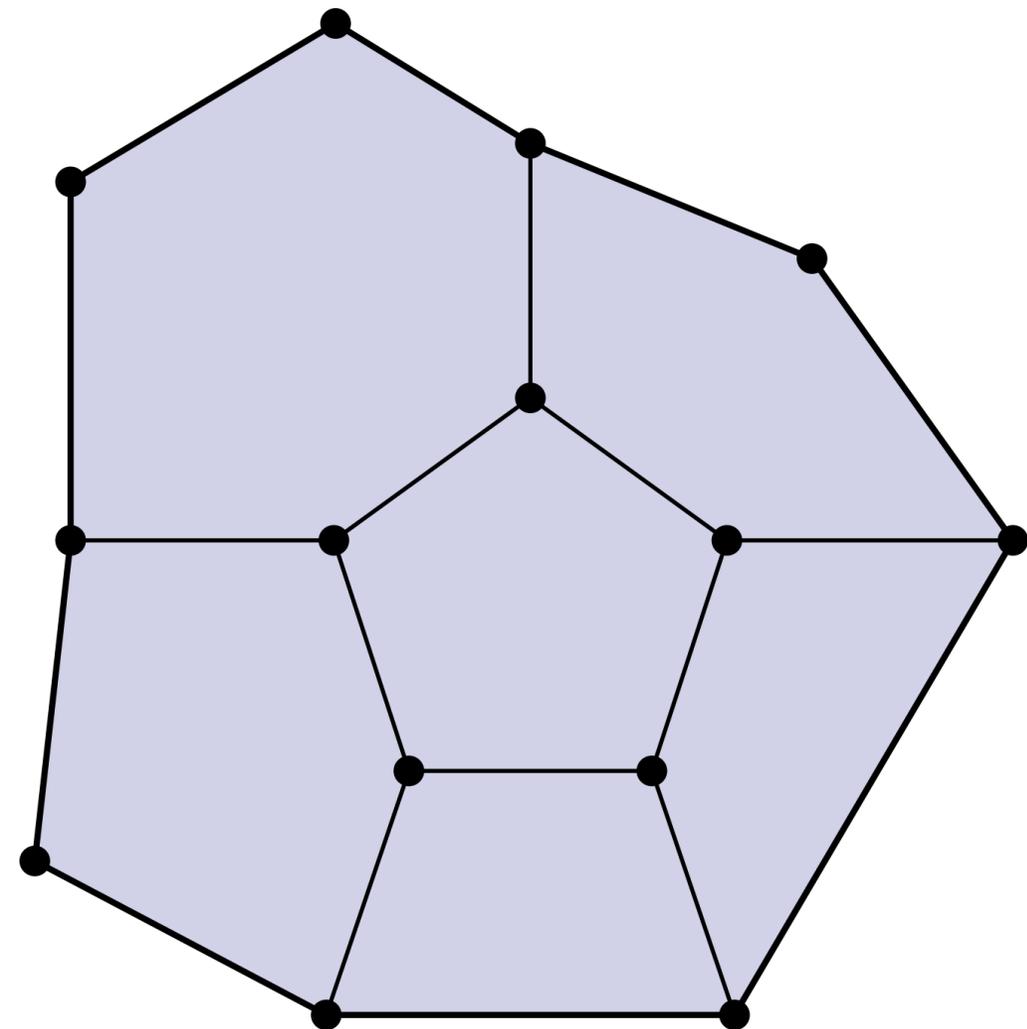
- Polygon mesh:
  - one polygon per boundary edge
  - boundary vertex looks like “pacman”

**Ok, but why is the manifold  
assumption *useful*?**

# Keep it Simple!

- **Same motivation as for images:**
  - **make some assumptions about our geometry to keep data structures/algorithms simple and efficient**
  - ***in many common cases*, doesn't fundamentally limit what we can do with geometry**

	$(i, j-1)$	
$(i-1, j)$	$(i, j)$	$(i+1, j)$
	$(i, j+1)$	



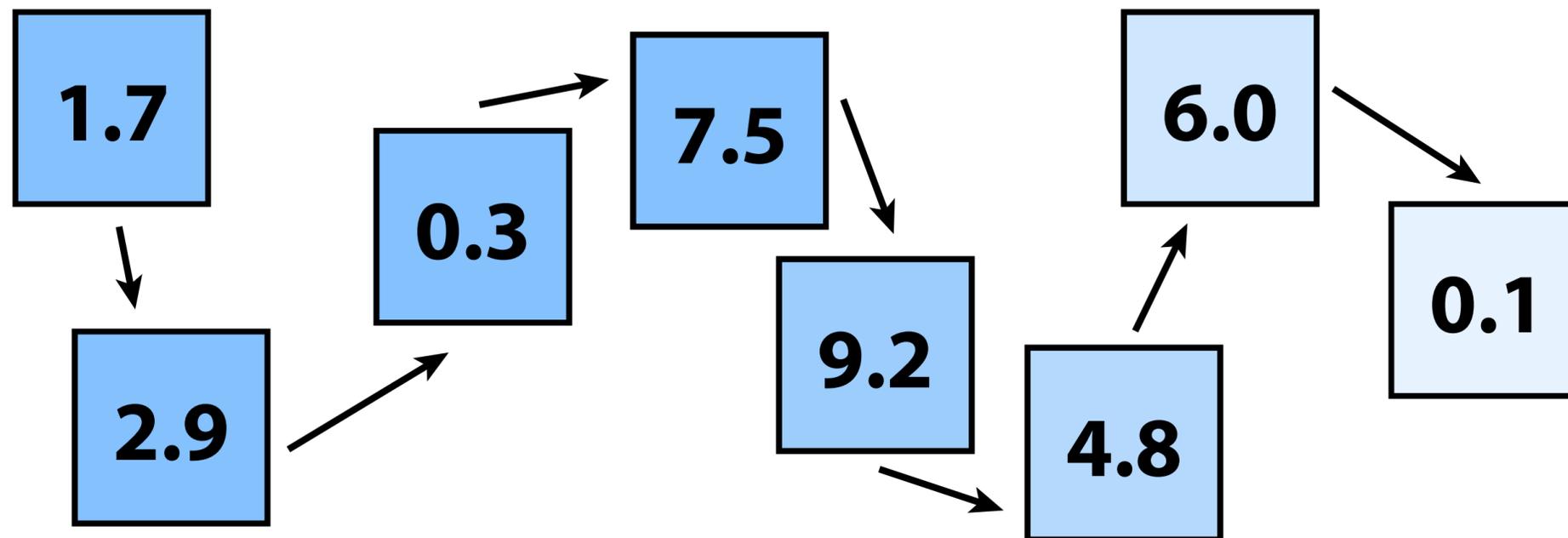
**How do we actually encode all this data?**

# Warm up: storing numbers

- Q: What data structures can we use to store a list of numbers?
- One idea: use an *array* (constant time lookup, coherent access)



- Alternative: use a linked list (linear lookup, incoherent access)



- Q: Why bother with the linked list?
- A: For one, we can easily insert numbers wherever we like...

# Polygon Soup

## ■ Most basic idea:

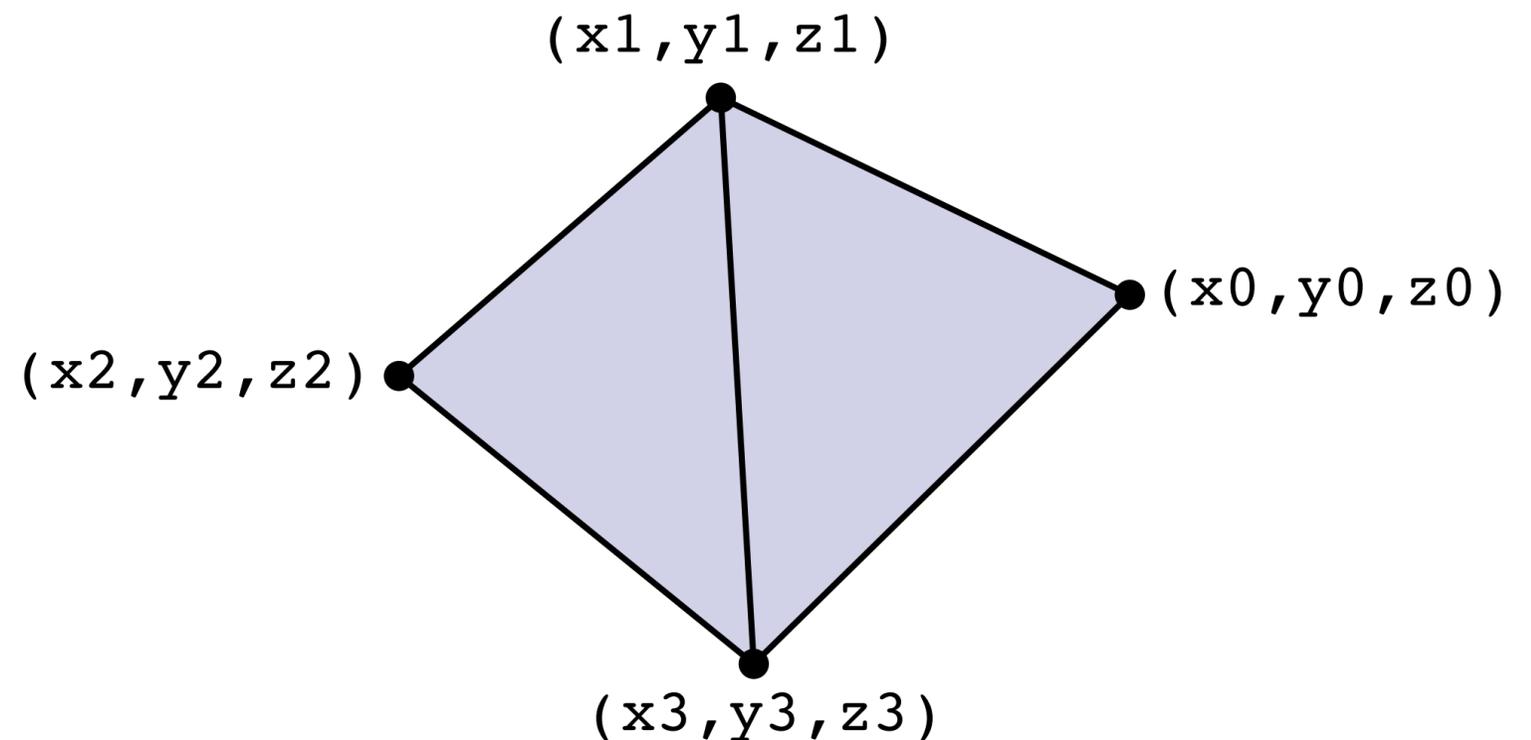
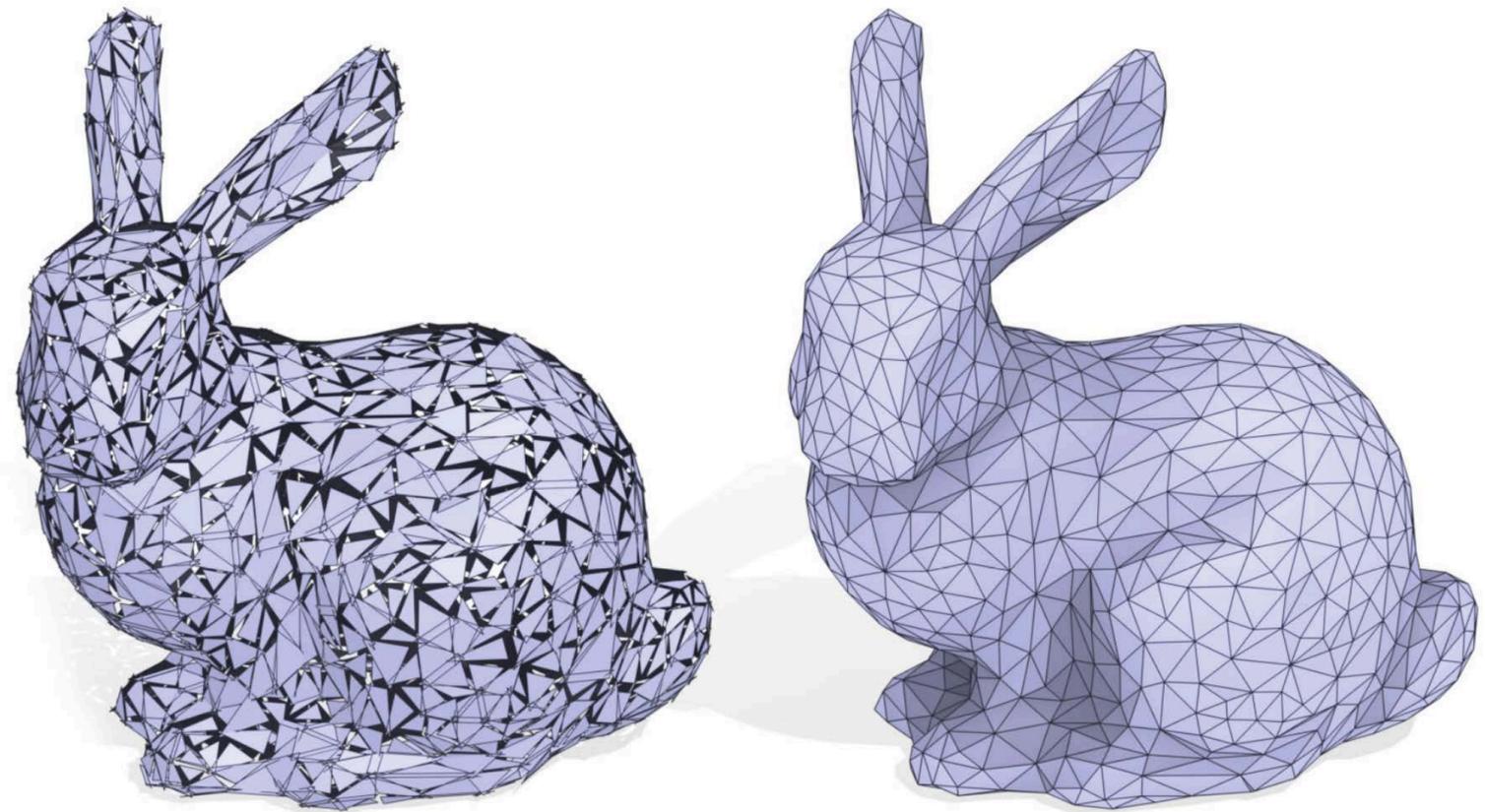
- For each triangle, just store three coordinates
- No other information about connectivity
- Not much different from point cloud! ("Triangle cloud?")

## ■ Pros:

- Really stupidly simple

## ■ Cons:

- Redundant storage
- Hard to do much beyond simply drawing the mesh on screen
- Need spatial data structures (later) to find neighbors



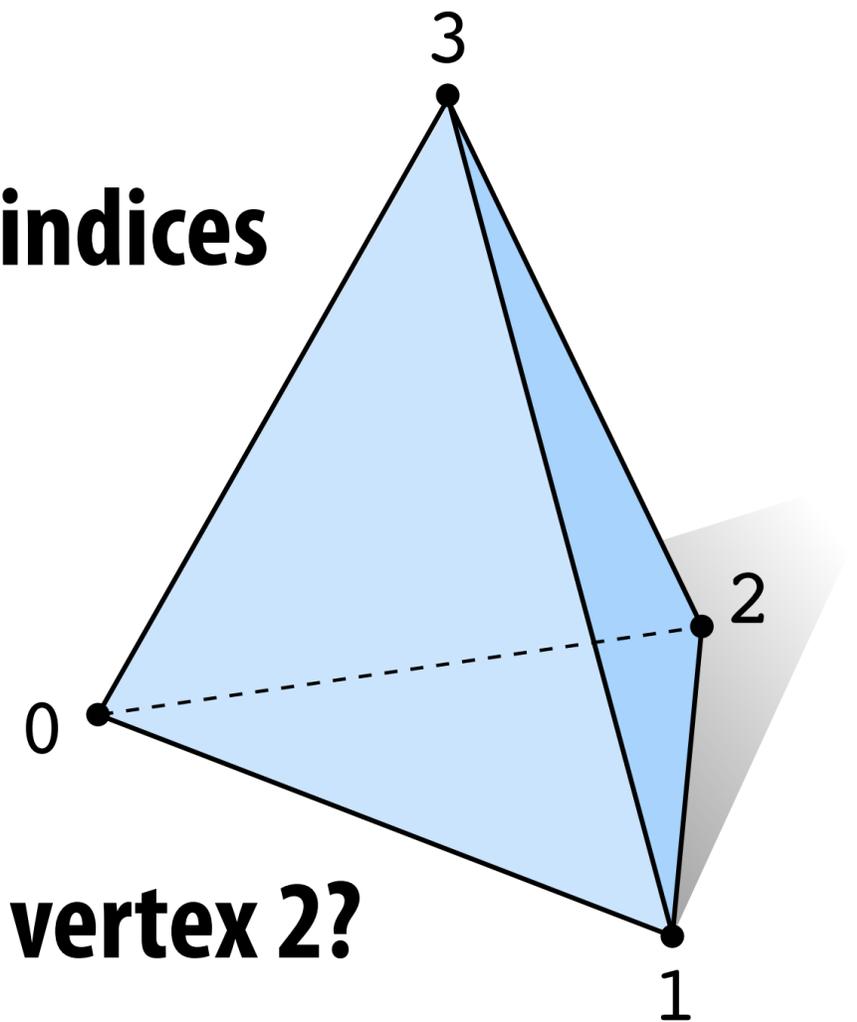
$x_0, y_0, z_0$	$x_1, y_1, z_1$	$x_3, y_3, z_3$
$x_1, y_1, z_1$	$x_2, y_2, z_2$	$x_3, y_3, z_3$

# Adjacency List (Array-like)

- Store triples of coordinates  $(x,y,z)$ , tuples of indices

- E.g., tetrahedron:

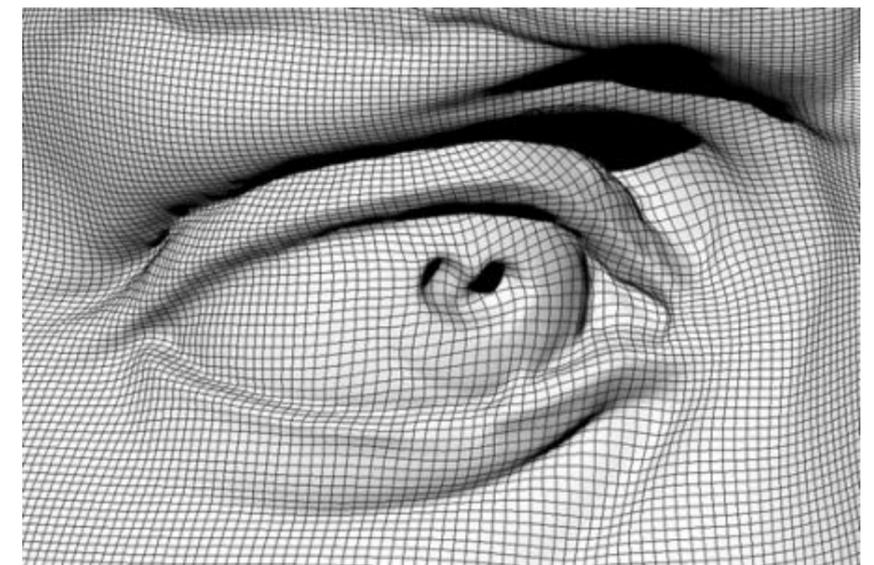
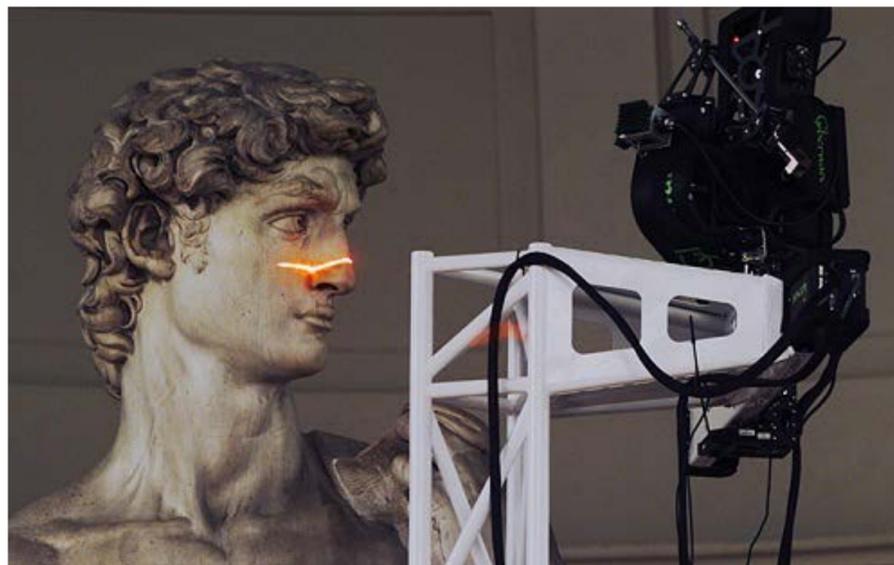
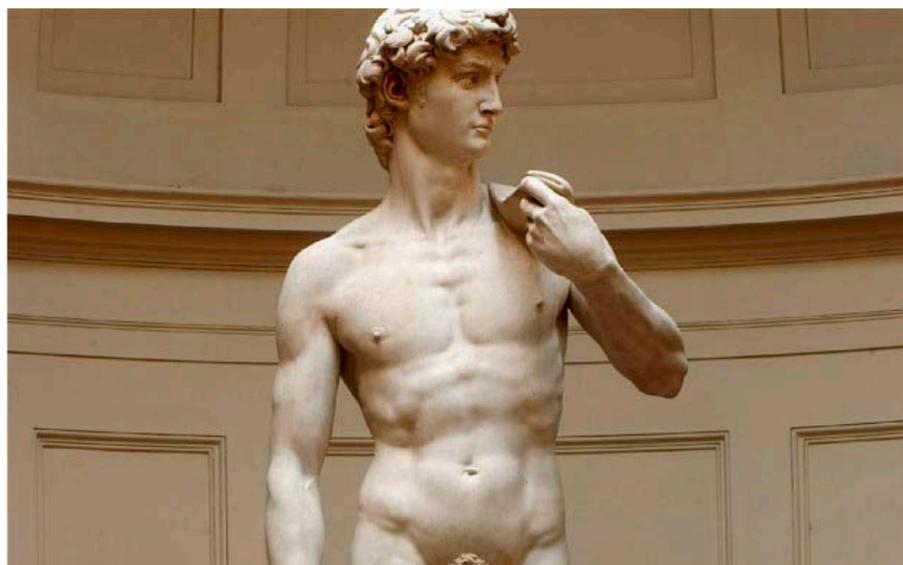
	VERTICES			POLYGONS		
	x	y	z	i	j	k
0:	-1	-1	-1	0	2	1
1:	1	-1	1	0	3	2
2:	1	1	-1	3	0	1
3:	-1	1	1	3	1	2



- Q: How do we find all the polygons touching vertex 2?

- Ok, now consider a more complicated mesh:

~1 billion polygons



**Very expensive to find the neighboring polygons! (What's the cost?)**

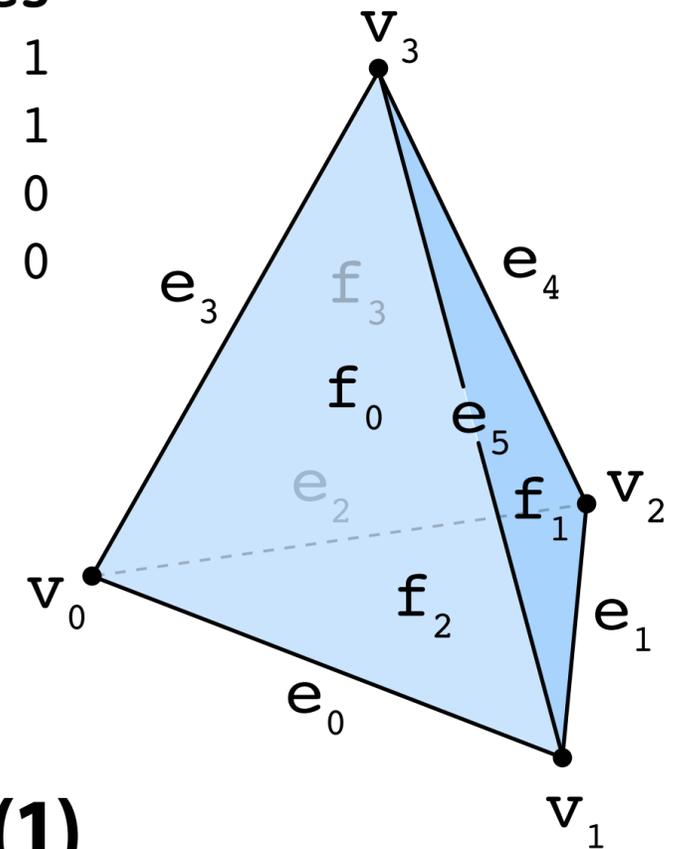
# Incidence Matrices

- If we want to know who our neighbors are, why not just store a list of neighbors?

- Can encode all neighbor information via *incidence matrices*

- E.g., tetrahedron:

	<u>VERTEX↔EDGE</u>				<u>EDGE↔FACE</u>						
	v0	v1	v2	v3	e0	e1	e2	e3	e4	e5	
e0	1	1	0	0	f0	1	0	0	1	0	1
e1	0	1	1	0	f1	0	1	0	0	1	1
e2	1	0	1	0	f2	1	1	1	0	0	0
e3	1	0	0	1	f3	0	0	1	1	1	0
e4	0	0	1	1							
e5	0	1	0	1							



- 1 means “touches”; 0 means “does not touch”

- Instead of storing lots of 0’s, use *sparse matrices*

- Still large storage cost, but finding neighbors is now  $O(1)$

- Hard to change connectivity, since we used fixed indices

- Bonus feature: mesh does not have to be manifold

# Aside: Sparse Matrix Data Structures

$$\begin{matrix} & 0 & 1 & 2 \\ 0 & \begin{bmatrix} 4 & 2 & 0 \end{bmatrix} \\ 1 & \begin{bmatrix} 0 & 0 & 3 \end{bmatrix} \\ 2 & \begin{bmatrix} 0 & 7 & 0 \end{bmatrix} \end{matrix}$$

- Ok, but how do we actually store a “sparse matrix”?

- Lots of possible data structures:

- Associative array from (row, column) to value

- easy to lookup/set entries, fast (e.g., hash table)

- harder to do matrix operations (e.g., multiplication)

(row, col) val

(0, 0)	->	4
(0, 1)	->	2
(1, 2)	->	3
(2, 1)	->	7

- Array of linked lists (one per row)

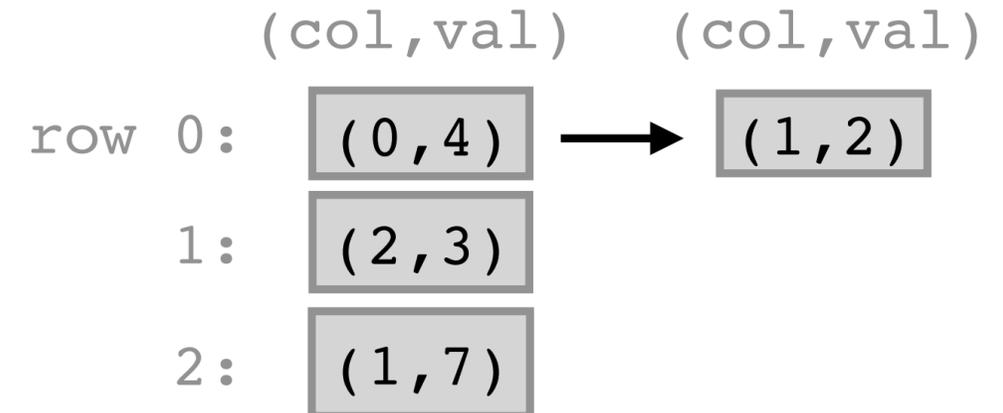
- conceptually simple

- slow access time, incoherent memory access

- Compressed column format—pack entries in list

- hard to add/modify entries

- fast for actual matrix operations



values 

4, 2, 7, 3
------------

row indices 

0, 0, 2, 1
------------

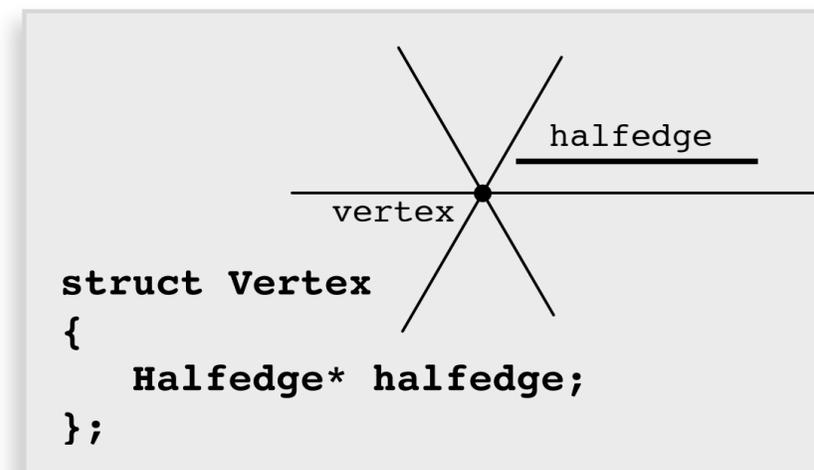
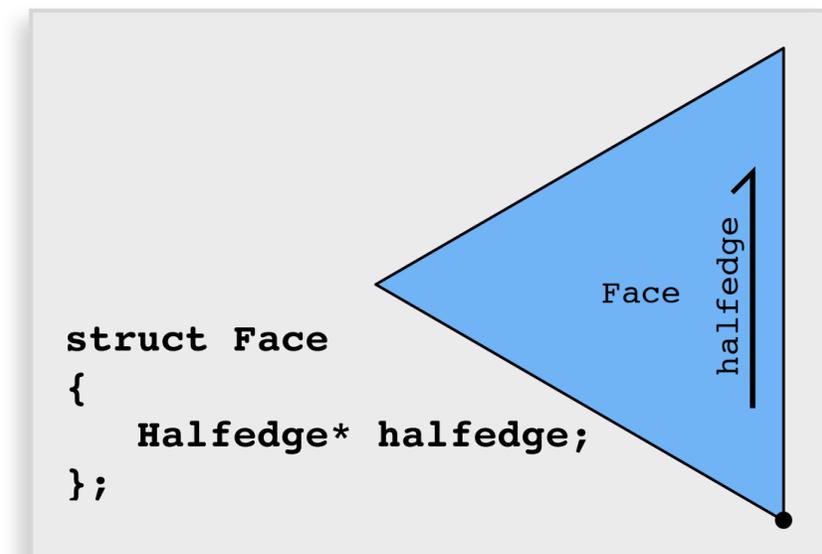
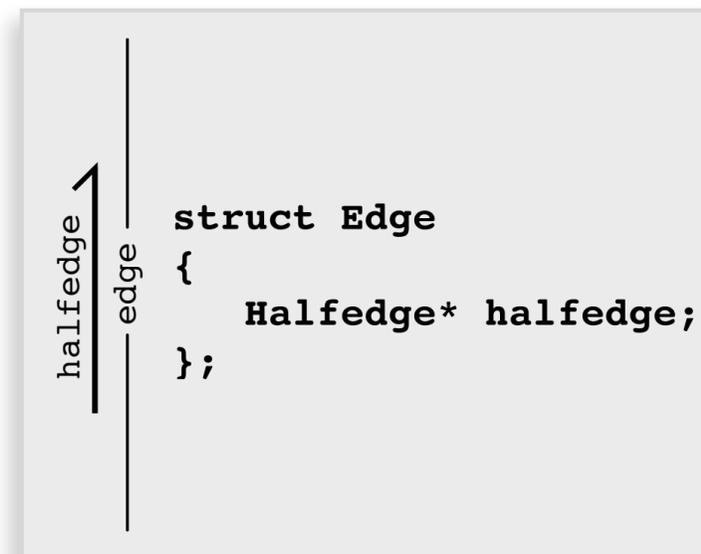
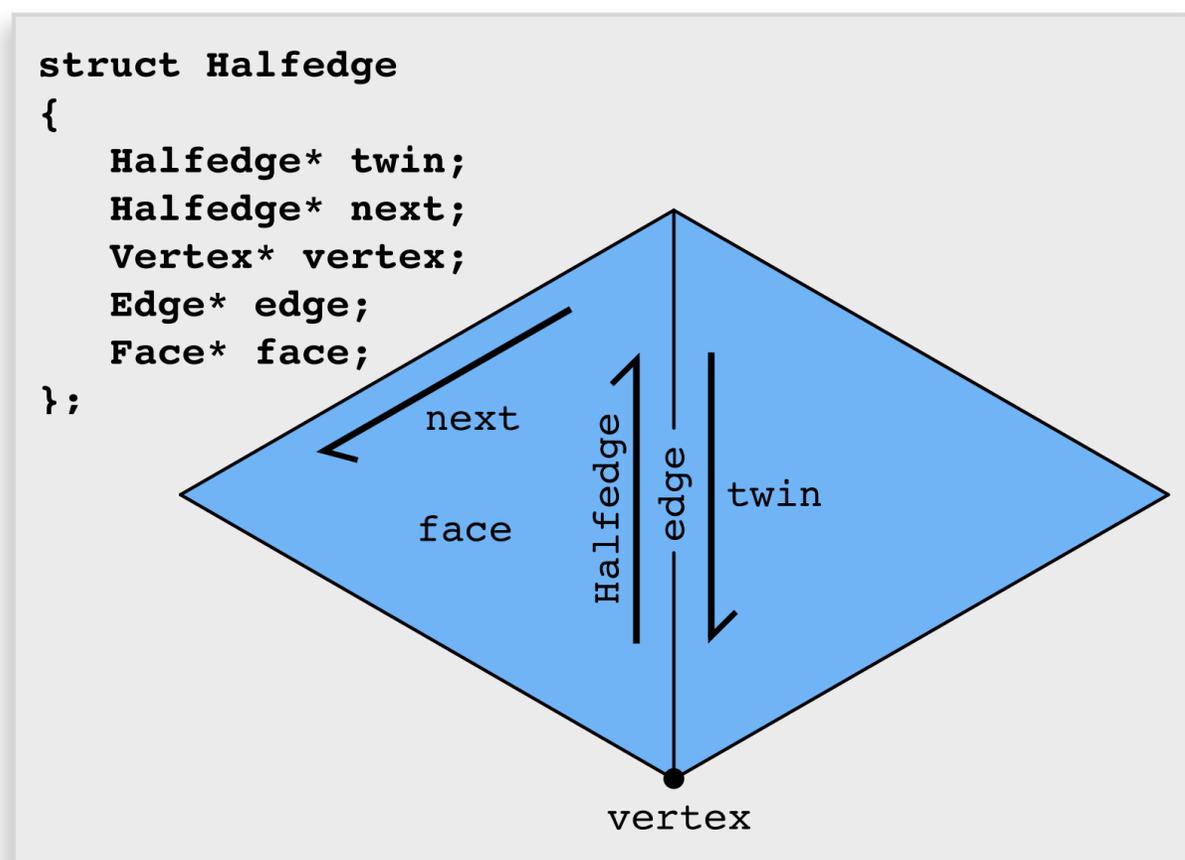
cumulative  
# entries  
by column 

1, 3, 4
---------

- In practice: often build up entries using an “easier” data structure, convert to compressed format for computation

# Halfedge Data Structure (Linked-list-like)

- Store *some* information about neighbors
- Don't need an exhaustive list; just a few key pointers
- Key idea: two *halfedges* act as "glue" between mesh elements:



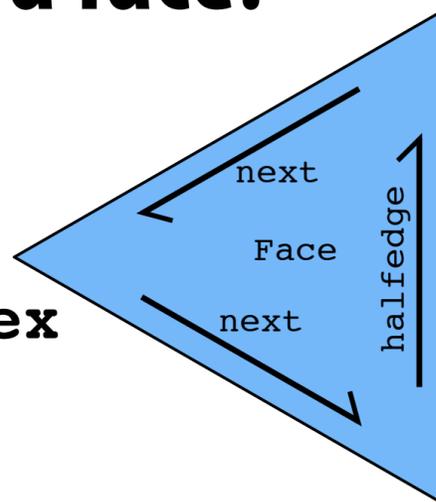
- Each vertex, edge face points to just *one* of its halfedges.

# Halfedge makes mesh traversal easy

- Use “twin” and “next” pointers to move around mesh
- Use “vertex”, “edge”, and “face” pointers to grab element

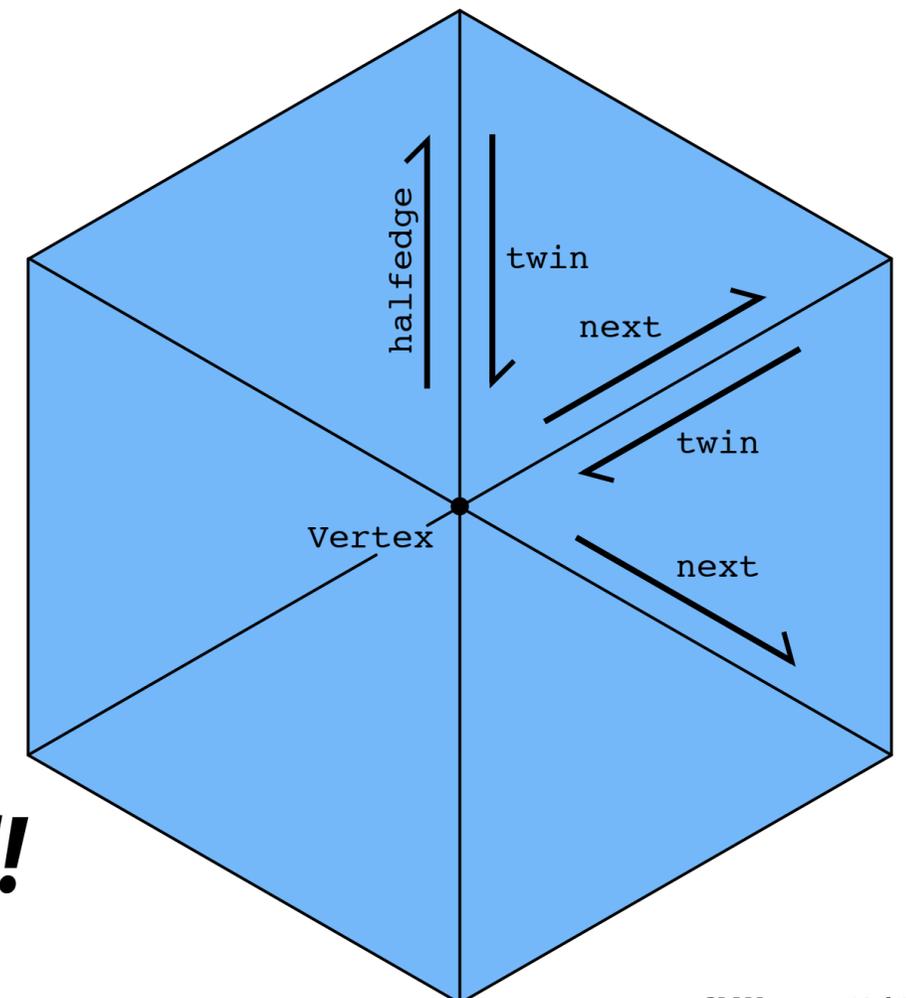
- Example: visit all vertices of a face:

```
Halfedge* h = f->halfedge;  
do {  
    h = h->next;  
    // do something w/ h->vertex  
}  
while( h != f->halfedge );
```



- Example: visit all neighbors of a vertex:

```
Halfedge* h = v->halfedge;  
do {  
    h = h->twin->next;  
}  
while( h != v->halfedge );
```



- Note: only makes sense if mesh is *manifold*!

# Halfedge connectivity is *always* manifold

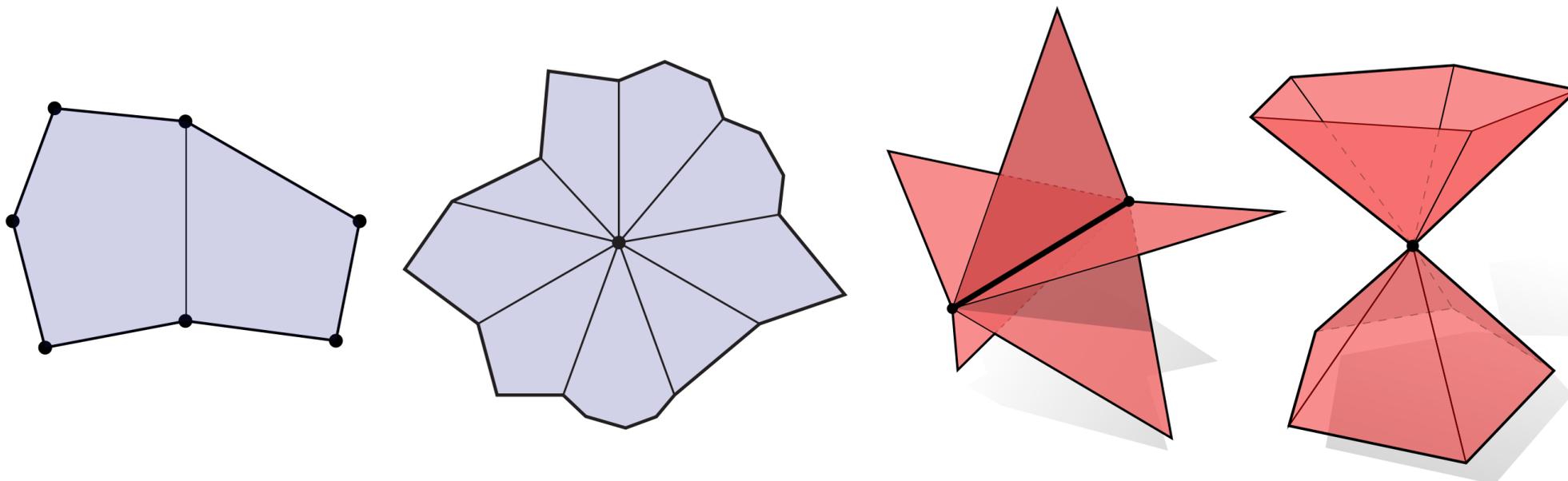
- Consider simplified halfedge data structure
- Require only “common-sense” conditions

```
struct Halfedge {  
    Halfedge *next, *twin;  
};
```

*(pointer to yourself!)*

```
twin->twin == this  
twin != this  
every he is someone's "next"
```

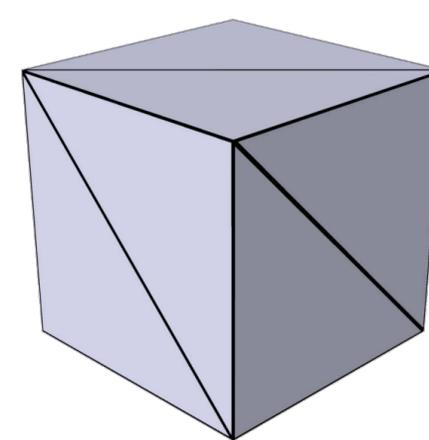
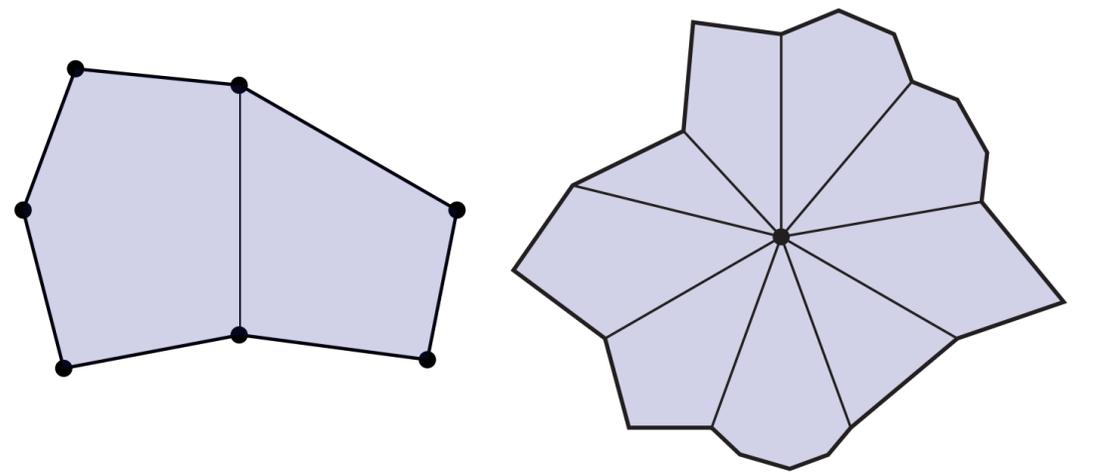
- Keep following `next`, and you'll get faces.
- Keep following `twin` and you'll get edges.
- Keep following `next->twin` and you'll get vertices.



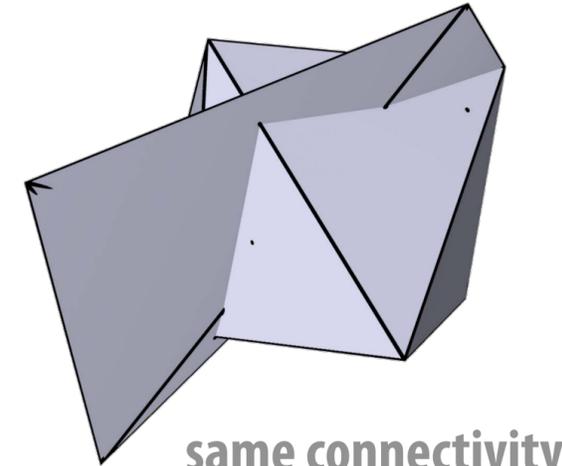
**Q: Why, therefore, is it impossible to encode the red figures?**

# Connectivity vs. Geometry

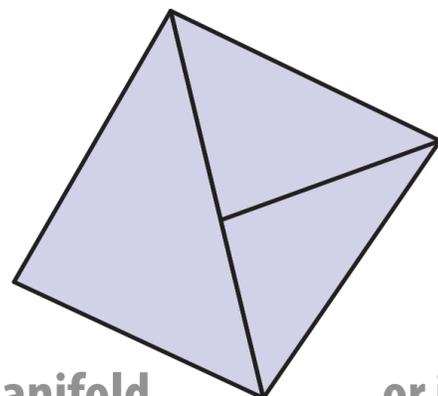
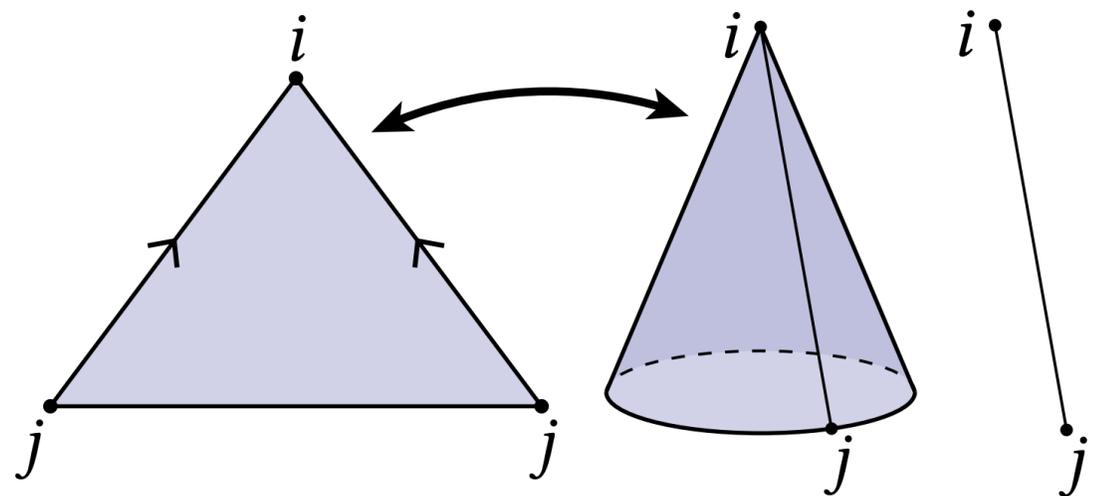
- Recall manifold conditions (fans not fins):
  - every edge contained in two faces
  - every vertex contained in one fan
- These conditions say nothing about vertex positions! Just connectivity
- Hence, can have perfectly good (manifold) connectivity, even if geometry is awful
- In fact, sometimes you can have perfectly good manifold connectivity for which any vertex positions give "bad" geometry!
- **Can lead to confusion when debugging: mesh looks "bad", even though connectivity is fine**



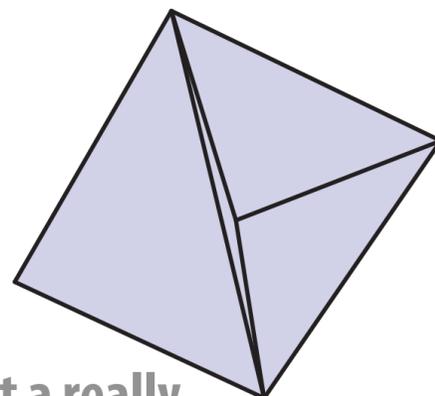
cube (manifold)



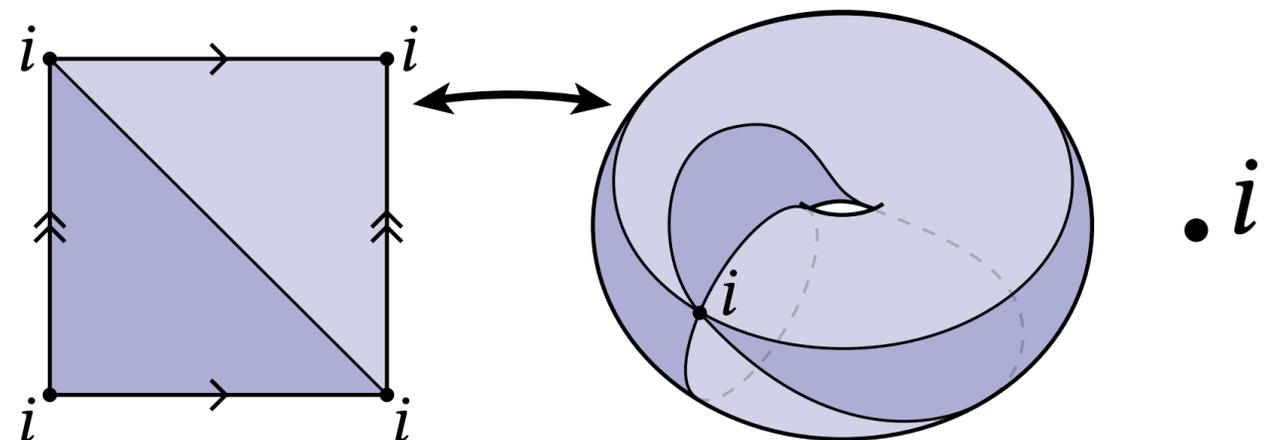
same connectivity, random vertex positions



non manifold connectivity?

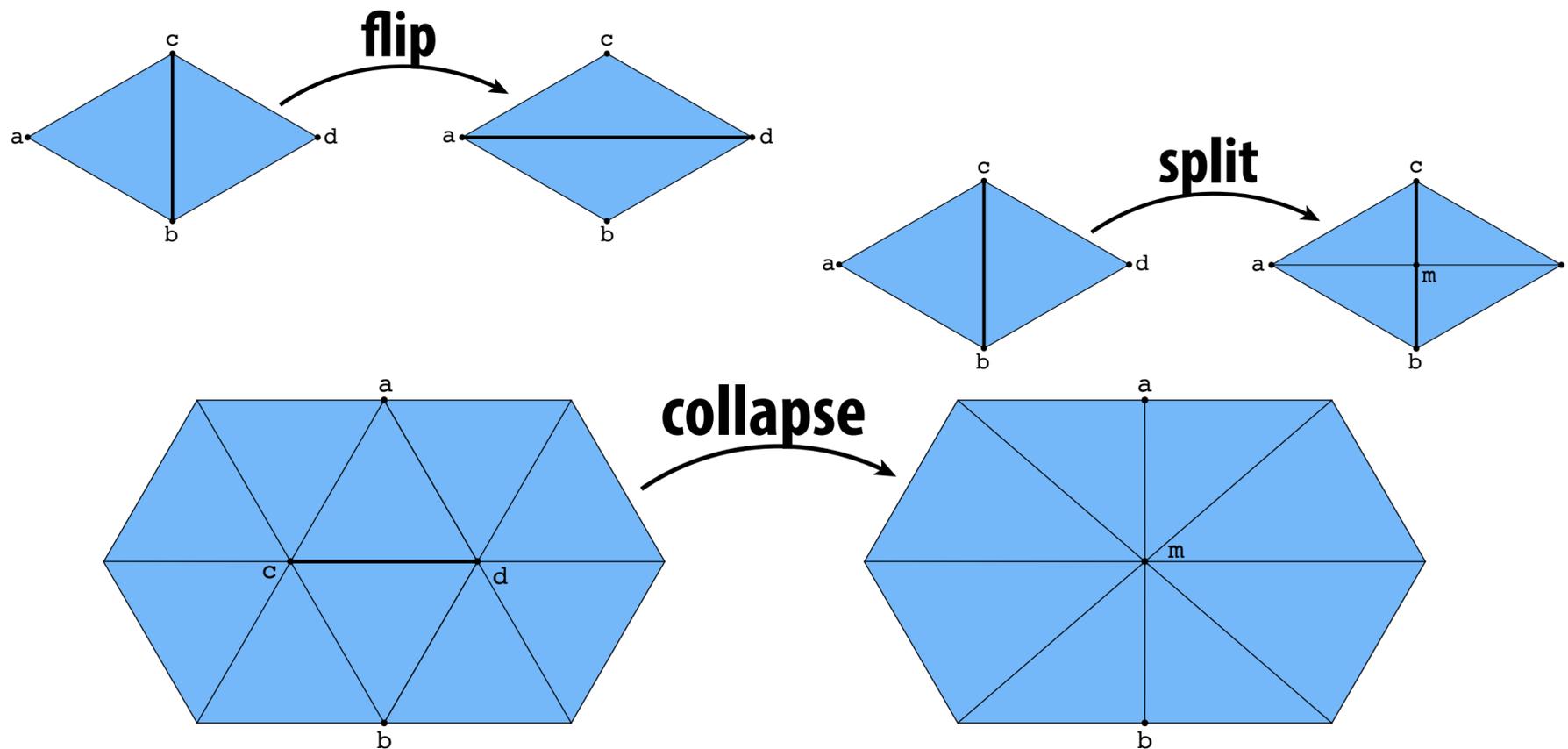


...or just a really skinny triangle?



# Halfedge meshes are easy to edit

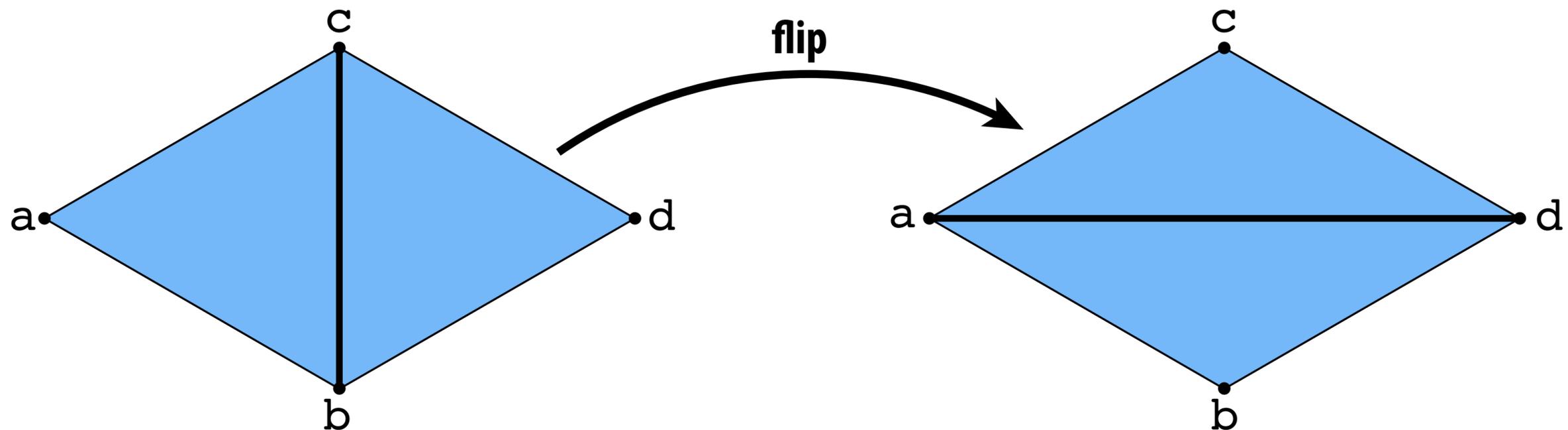
- Remember key feature of linked list: insert/delete elements
- Same story with halfedge mesh (“linked list on steroids”)
- E.g., for triangle meshes, several atomic operations:



- How? Allocate/delete elements; reassigning pointers.
- Must be careful to preserve manifoldness!

# Edge Flip (Triangles)

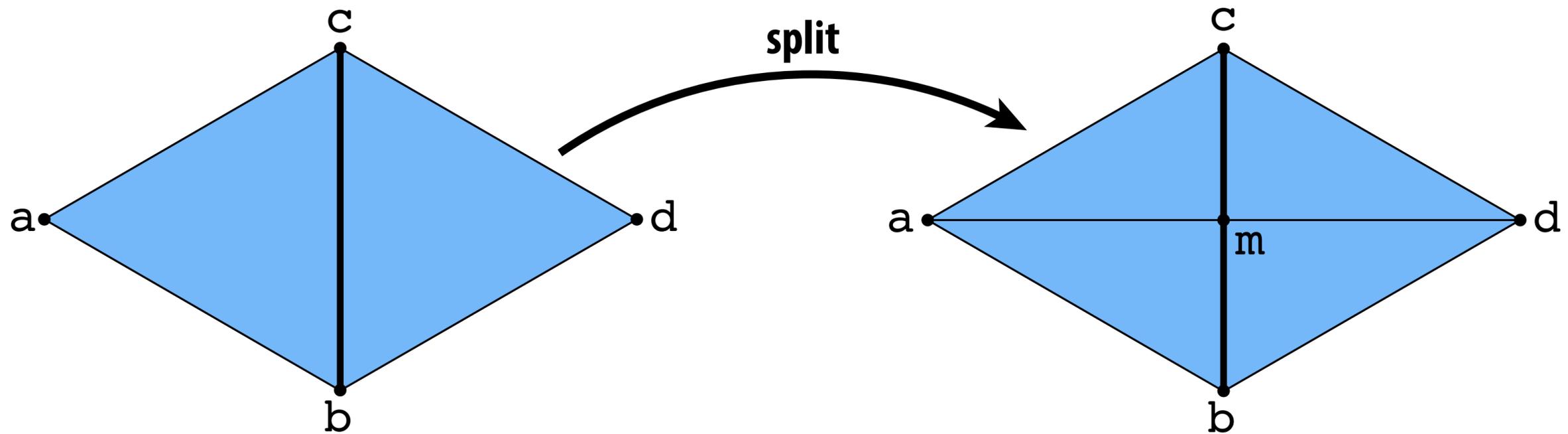
- Triangles  $(a,b,c)$ ,  $(b,d,c)$  become  $(a,d,c)$ ,  $(a,b,d)$ :



- Long list of pointer reassignments (`edge->halfedge = ...`)
- However, no elements created/destroyed.
- Q: What happens if we flip twice?
- Challenge: can you implement edge flip such that pointers are *unchanged* after two flips?

# Edge Split (Triangles)

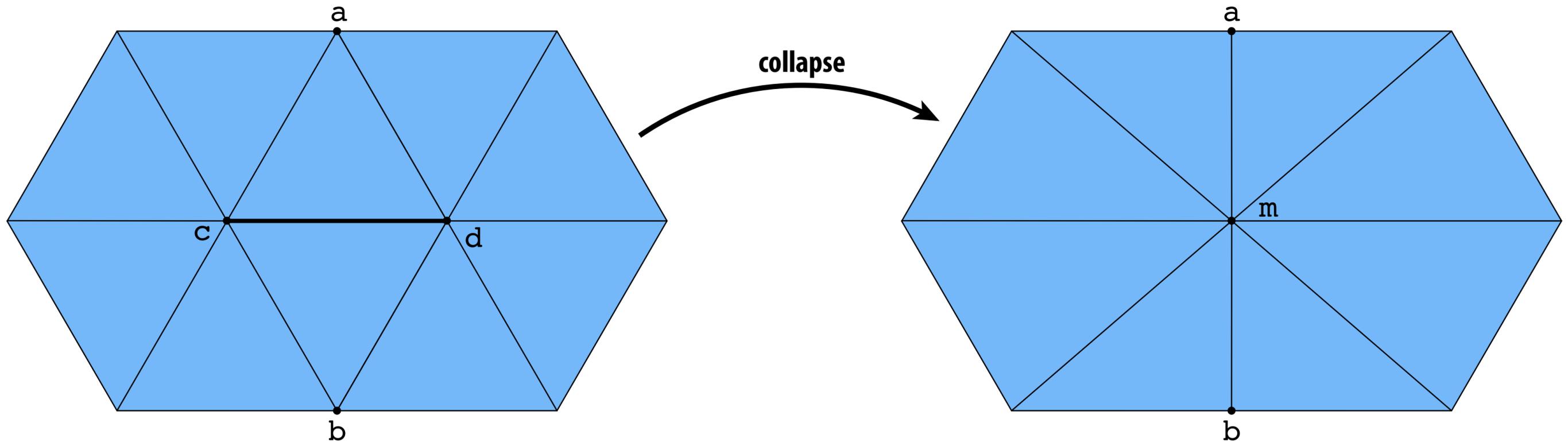
- Insert midpoint  $m$  of edge  $(c,b)$ , connect to get four triangles:



- This time, have to *add* new elements.
- Lots of pointer reassignments.
- Q: Can we “reverse” this operation?

# Edge Collapse (Triangles)

- Replace edge (b,c) with a single vertex m:



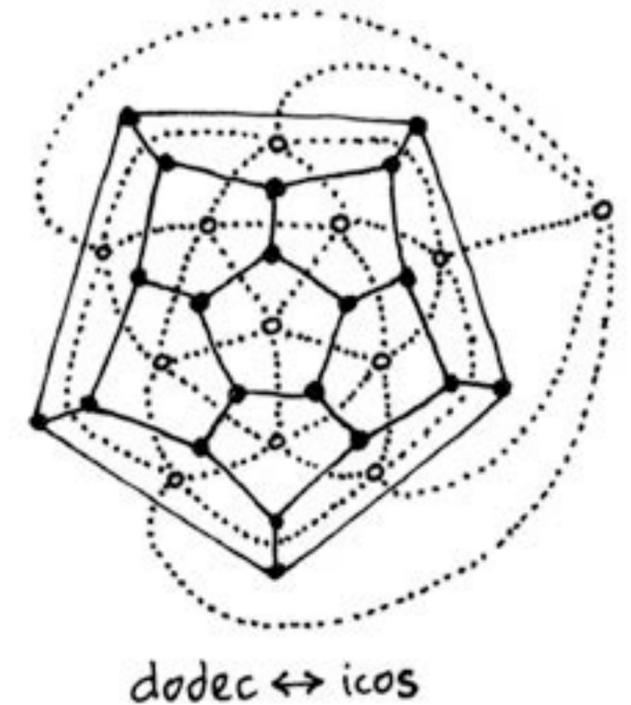
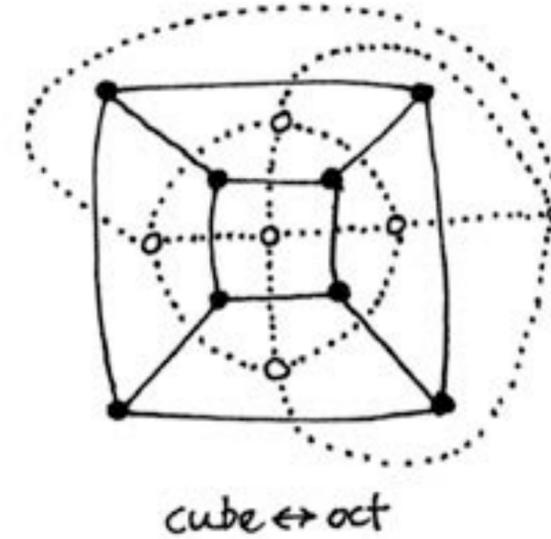
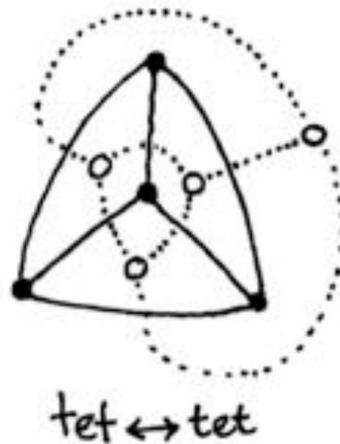
- Now have to *delete* elements.
- Still lots of pointer assignments!
- Q: How would we implement this with an adjacency list?
- Any other good way to do it? (E.g., different data structure?)

# Alternatives to Halfedge

Paul Heckbert (former CMU prof.)  
quadedge code - <http://bit.ly/1QZLHos>

## ■ Many very similar data structures:

- winged edge
- corner table
- quadedge
- ...



## ■ Each stores local neighborhood information

## ■ Similar tradeoffs relative to simple polygon list:

- **CONS:** additional storage, incoherent memory access
- **PROS:** better access time for individual elements, intuitive traversal of local neighborhoods

## ■ With some thought\*, can design halfedge-type data structures with coherent data storage, support for non manifold connectivity, etc.

\*see for instance <http://geometry-central.net/>

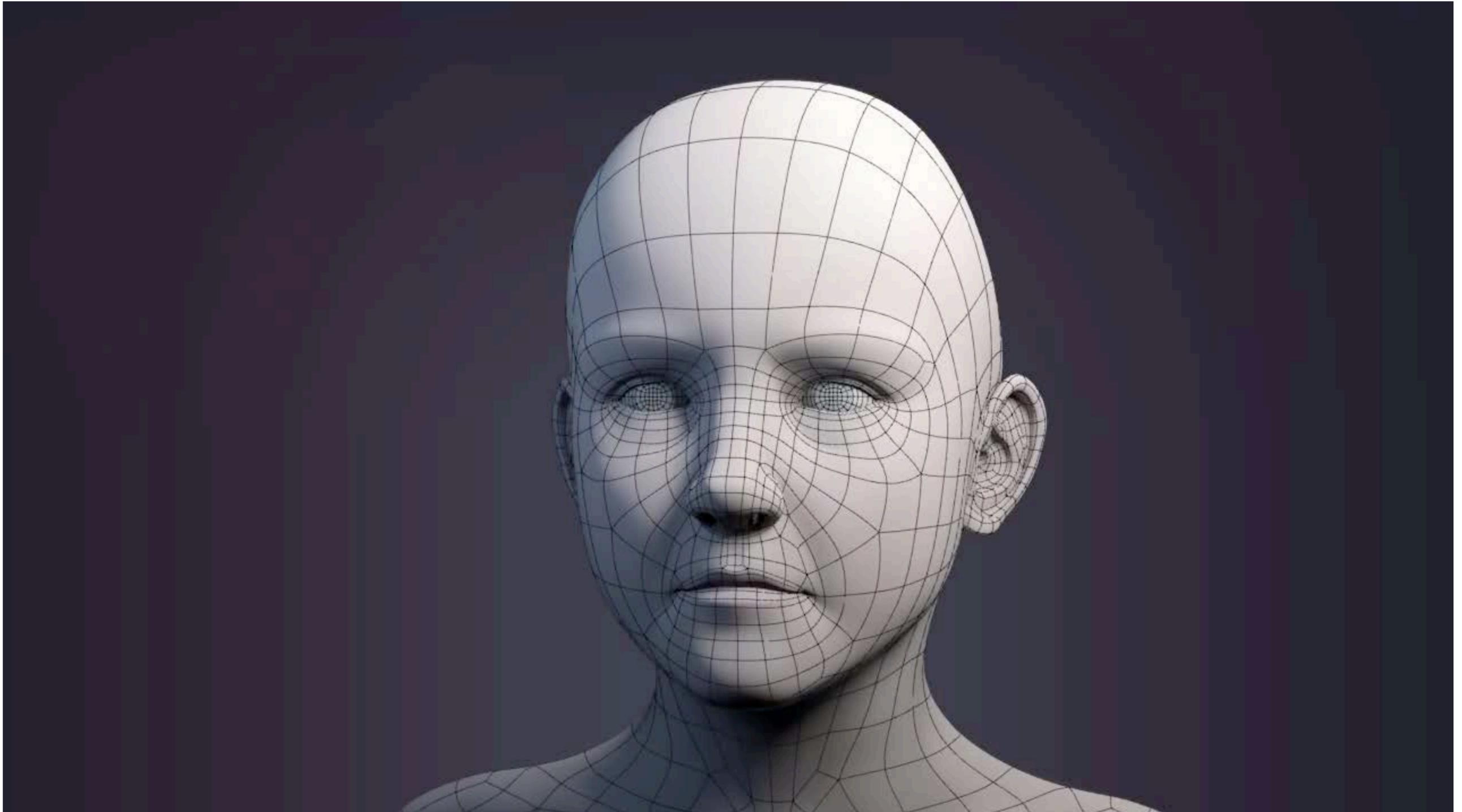
# Comparison of Polygon Mesh Data Structures

	Adjacency List	Incidence Matrices	Halfedge Mesh
constant-time neighborhood access?	NO	YES	YES
easy to add/remove mesh elements?	NO	NO	YES
nonmanifold geometry?	YES	YES	NO

**Conclusion: pick the right data structure for the job!**

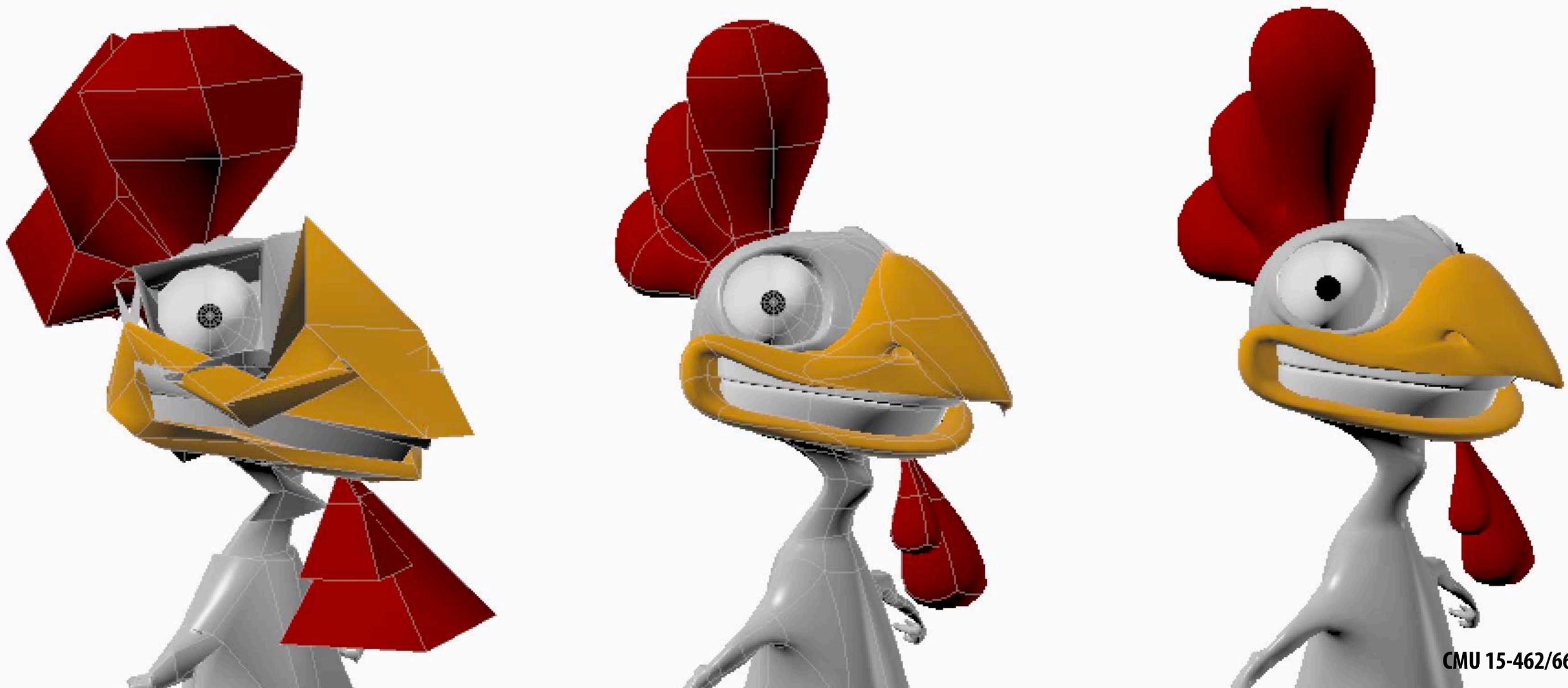
**Ok, but what can we actually *do* with our fancy new data structures?**

# Subdivision Modeling



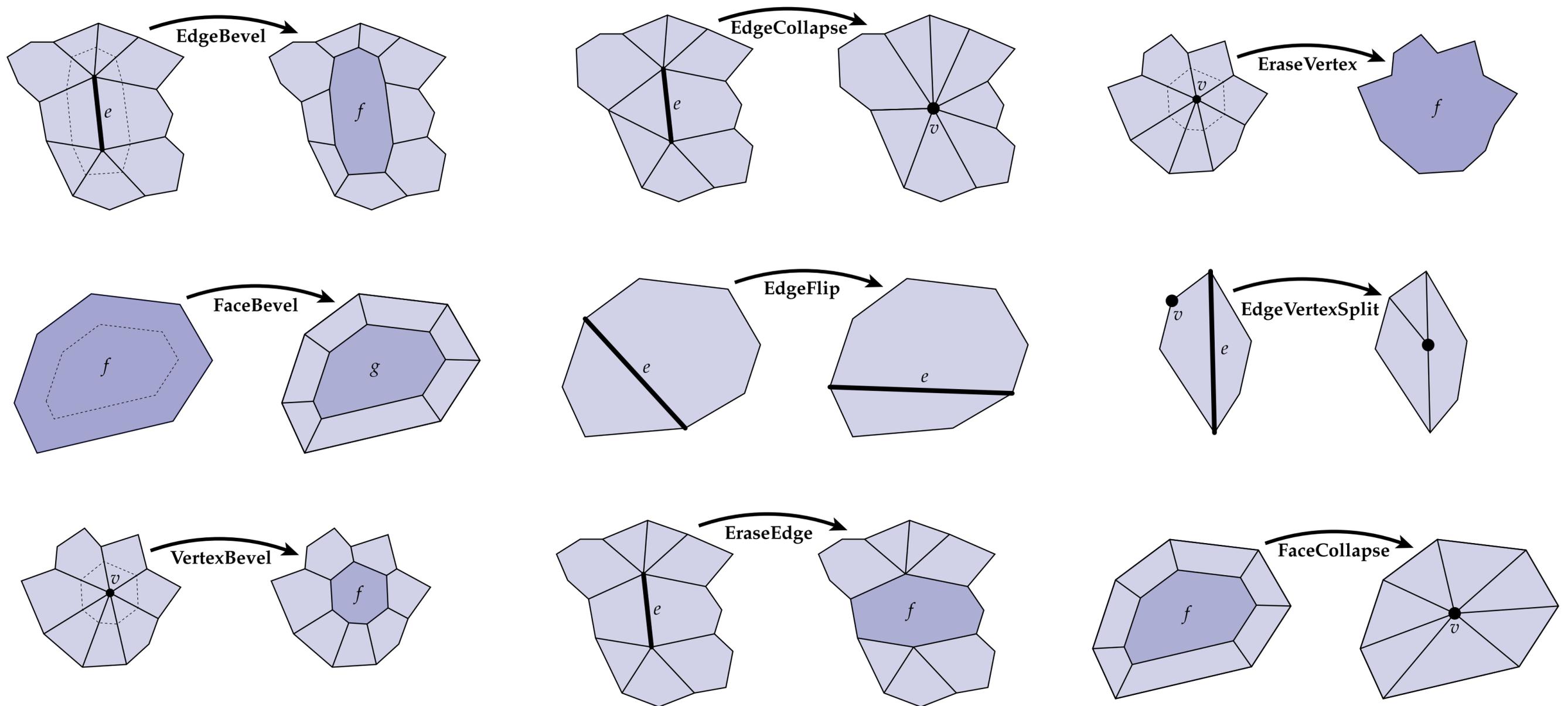
# Subdivision Modeling

- **Common modeling paradigm in modern 3D tools:**
  - **Coarse “control cage”**
  - **Perform local operations to control/edit shape**
  - **Global subdivision process determines final surface**



# Subdivision Modeling—Local Operations

- For general polygon meshes, we can dream up lots of local mesh operations that might be useful for modeling:



...and many, many more!

# Next Time: Digital Geometry Processing

- **Extend traditional digital signal processing (audio, video, etc.) to deal with *geometric* signals:**
  - **upsampling / downsampling / resampling / filtering ...**
  - **aliasing (reconstructed surface gives “false impression”)**

